

Anyone and everyone who routinely works with equations needs Eureka: The Solver

It solves the most complex equations in seconds. Whether you're a scientist, engineer, financial analyst, student, teacher, or some other professional, you need Eureka: The Solver!

Any problem that can be expressed as a linear or non-linear equation can be solved with Eureka. Algebra, Trigonometry and Calculus problems are a snap.

Eureka: The Solver also handles maximization and minimization problems, does plot functions, generates reports, and saves you an incredible amount of time.

X+exp(X) = 10 solved instantly instead of eventually!

Imagine you have to "solve for X," where $X + \exp(X) = 10$, and you don't have Eureka: The Solver. What you do have is a problem, because it's going to take a lot of time guessing at "X." Maybe your guesses get closer and closer to the right answer, but it's also getting closer and closer to midnight and you're doing it the hard way.

With Eureka: The Solver, there's no guessing, no dancing in the dark—you get the right answer, right now. (PS: X = 2.0705799, and Eureka solved that one in .4 of a second!)

How to use Eureka: The Solver

It's easy.

- 1. Enter your equation into the full-screen editor
- 2. Select the "Solve" command
- 3. Look at the answer
- 4. You're done

You can then tell Eureka to

- Evaluate your solution
- Plot a graph
- Generate a report, then send the output to your printer, disk file or screen
- Or all of the above

Eureka: The Solver includes

- ✓ Pull-down menus
- ▼ Context-sensitive Help
- ✓ On-screen calculator
- ✓ Automatic 8087 math co-processor chip support
- ▼ Powerful financial functions
- ☑ Built-in and user-defined math and financial functions
- Ability to generate reports complete with plots and lists

*Introductory price—good through July 1, 1987

- ✓ Polynomial finder
- Inequality solutions

Some of Eureka's key features

You can key in:

- A formula or formulas
- ✓ A series of equations—and solve for all variables
- ✓ A function to plot
- ✓ Unit conversions
- Maximization and minimization problems
- ✓ Interest Rate/Present Value calculations
- ✓ Variables we call "What happens?," like "What happens if I change this variable to 21 and that variable to 27?"

All this power for only \$99.95!

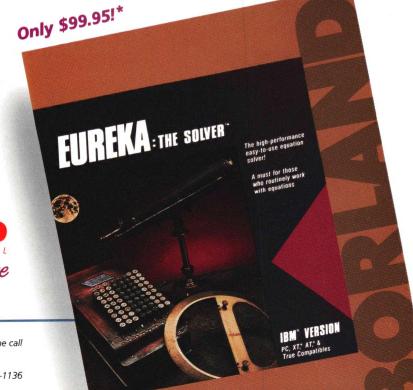
Equation-solving used to be a mainframe problem, but we've solved that problem.

Eureka: The Solver is all you need—and it's yours for only \$99.95!

That kind of savings you can calculate with your fingers!

System requirements

IBM PC, AT, XT, Portable, 3270 or true compatibles. PC-DOS (MS-DOS) 2.0 and later. 384K.





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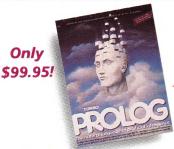
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Turbo Prolog

artificial intelligence, databases, expert systems, or new ways of thinking about programming, by all means plunk down your \$100 and buy a copy of Turbo Prolog.

Bruce Webster, BYTE ""



Turbo Prolog, the natural language of Artificial Intelligence, is the most popular AI package in the world with more than 100,000 users. It's the 5th-generation computer programming language that brings supercomputer power to your IBM PC and compatibles. You can join the AI revolution with Turbo Prolog for only \$99.95. Step-by-step tutorials, demo programs and source code included.

New! Turbo Prolog Toolbox

Our new Turbo Prolog Toolbox" enhances Turbo Prolog—with more than 80 tools and over 8,000 lines of source code that can easily be incorporated into your programs. It includes about 40 example programs

that show you how to use and incorporate your new tools.

New Turbo Prolog Toolbox features include:

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NEW

- ✓ Complete communications package ✓ File transfers from Reflex, dBASE III,
- 1-2-3, Symphony

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It's the complete developer's toolbox and a major addition to Turbo Prolog. You get a wide variety of menus—pull-down, pop-up, line, tree and box—so you can choose the one that suits your application best. You'll quickly and easily learn how to produce graphics; set up communications with remote devices; read information from Reflex,* dBASE III,* Lotus 1-2-3* and Symphony* files; generate parsers and design user interfaces. All of this for only \$99.95.

Only \$99.95!



System requirements

Turbo Prolog: IBM PC, XT, AT or true compatibles. PC-DOS (MS-DOS) 2.0 or later. 384K. Turbo Prolog Toolbox requires Turbo Prolog 1.10 or higher. Dual-floppy disk drive or hard disk. 512K.



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6 The language deal of the century. **Jeff Duntemann, PC Magazine**

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Turbo Pascal 3.0.

Includes 8087 & BCD features for 16-bit MS-DOS and CP/M-86 systems. CP/M-80 version minimum memory: 48K; 8087 and BCD features not available. 128K.

Turbo Basic

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The Critics' Choice

stretching the language without weighing us down with unnecessary details . . . Turbo Basic is the answer to my wish for a simple yet blindingly fast recreational utility language . . . The one language you can't forget how to use, Turbo Basic is a computer language for the missus, the masters, the masses, and me.

Steve Gibson, InfoWorld

Borland's Turbo Basic has advantages over the Microsoft product, including support of the highspeed 8087 math chip.

John C. Dvorak 77

Turbo Basic ends the basic confusion

There's now one standard: Turbo Basic.

It's fast, BASICA-compatible, and because Turbo Basic is a Borland product, the price is right, the quality is there, and the power is at your fingertips. You see, Turbo Basic's part of the fast-growing Borland family of programming languages—we call it the "Turbo Family." Hundreds of thousands of users are already using Borland's languages, so you can't go wrong. So join a whole new generation of smart IBM PC users—get your copy of Turbo Basic today. You get an easy-to-read 300+ page manual, two disks, and a free MicroCalc spreadsheet—and an instant start in the fast new world of Turbo Basic. All of this for only \$99.95—Order your copy of Turbo Basic today!

Free spreadsheet included, complete with source code!

Yes, we've included MicroCalc, our sample spreadsheet, complete with source code, so that you can get started right away with a "real program." You can compile and run it "as is," or modify it.

A technical look at Turbo Basic

- ▼ Full recursion supported
- ✓ Standard IEEE floating-point format
- ✓ Floating-point support, with full 8087 (math co-processor) integration. Software emulation if no 8087 present
- ✓ Program size limited only by available memory (no 64K limitation)
- Access to local, static, and global variables
- Full integration of the compiler, editor, and executable program, with separate windows for editing, messages, tracing, and execution
- ✓ Compile, run-time, and I/O errors place you in the source code where error occurred
- ✓ New long integer (32-bit) data type
- ▼ Full 80-bit precision
- ✓ Pull-down menus
- ✓ Full window management

System requirements

IBM PC, XT, AT and true compatibles, PC-DOS (MS-DOS) 2.0 or later. One floppy drive, 256K.





Turbo C: The fastest, most efficient and easy-to-use C compiler at any price

Compilation speed is more than 7000 lines a minute, which makes anything less than Turbo C an exercise in slow motion. Expect what only Borland delivers: Quality, Speed, Power and Price.

Turbo C: The C compiler for amateurs and professionals

If you're just beginning and you've "kinda wanted to learn C," now's your chance to do it the easy way. Like Turbo Pascal, Turbo C's got everything to get you going.

If you're already programming in C, switching to Turbo C will considerably increase your productivity and help make your programs both smaller and faster. Actually, writing in Turbo C is a highly productive and effective method—and we speak from experience. Eureka: The Solver and our new generation of software have been developed using Turbo C.

Turbo C: a complete interactive development environment

Free MicroCalc spreadsheet with source code

Like Turbo Pascal and Turbo Prolog, Turbo C comes

with an interactive editor that will show you syntax errors right in your source code. Developing, debugging, and running a Turbo C program is a snap.

Turbo C: The C compiler everybody's been waiting for. Everybody but the competition

Borland's "Quality, Speed, Power and Price" commitment isn't idle corporate chatter. The \$99.95 price tag on Turbo C isn't a "typo," it's real. So if you'd like to learn C in a hurry, pick up the phone. If you're already using C, switch to Turbo C and see the difference for yourself.

System requirements

IBM PC, XT, AT and true compatibles. PC-DOS (MS-DOS) 2.0 or later. One floppy drive. 320K.

Technical Specifications

- Compiler: One-pass compiler generating linkable object modules and inline assembler. Included is Borland's high performance "Turbo Linker." The object module is compatible with the PC-DOS linker. Supports tiny, small, compact, medium, large, and huge memory model libraries. Can mix models with near and far pointers. Includes floating point emulator (utilizes 8087/80287 if installed)
- ✓ Interactive Editor: The system includes a powerful, interactive full-screen text editor. If the compiler detects an error, the editor automatically positions the cursor appropriately in the source code.
- ✓ Development Environment: A powerful "Make" is included so that managing Turbo C program development is highly efficient. Also includes pull-down menus and windows
- ✓ Links with relocatable object modules created using Borland's Turbo Prolog into a single program.
- ANSI C compatible.
- ✓ Start-up routine source code included.
- ☑ Both command line and integrated environment versions included.

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Sieve benchmark (25 iterations)

	Turbo C	Microsoft® C	Lattice C
Compile time	3.89	16.37	13.90
Compile and link time	9.94	29.06	27.79
Execution time	5.77	9.51	13.79
Object code size	274	297	301
Price	\$99.95	\$450.00	\$500.00

Benchmark run on a 6 Mhz IBM AT using Turbo C version 1.0 and the Turbo Linker version 1.0; Microsoft C version 4.0 and the MS overlay linker version 3.51; Lattice C version 3.1 and the MS object linker version 3.05.

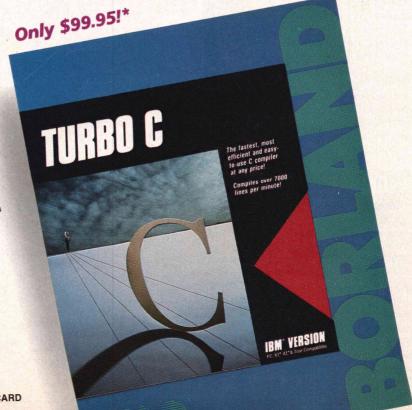
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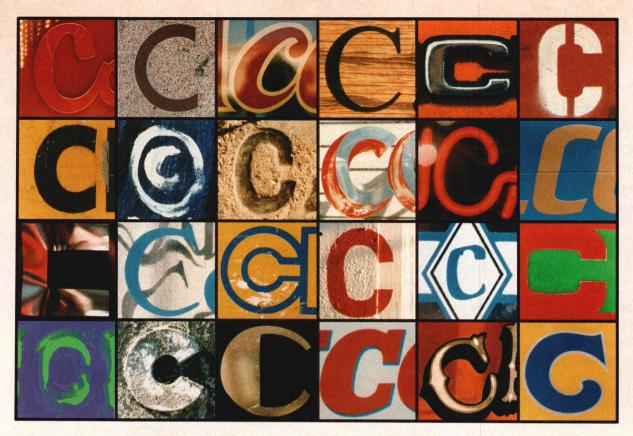
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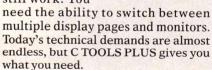




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also includes the "XMODEM" file-transfer protocol and support for Hayes-compatible modems. All source code is included for \$175. C TOOLS & C TOOLS 2—an indispensable combination still available at a low price of \$175, including all source code. See review in PC Tech Journal, 6/85.

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	ARTICLES
Efficient > queue control	ALGORITHMS: An Efficient Algorithm for Large Priority Queues by Robert Jay Brown By deferring decisions about lower-level priorities, this algorithm is able to assign the top priority more efficiently. COMMUNICATIONS: Two-Bit Analog-to-Digital Conversion
Mathematical basis for a space-saving trick	by John Musselman How to use hardware interrupts and a demonstration of techniques common in real-time programming ALGORITHMS: The XOR Chain
TSR serial driver	by David E. Cortesi The Swapped-Out Intern returns with a technique that exploits a curious property of doubly linked lists. COMMUNICATIONS: An Extended COM Port Driver by Thomas A. Zimniewicz Excom is a terminate-and-stay-resident program that gives the IBM PC interrupt-driven buffered input with flow-
Handling large Turbo Pascal programs Unix telecommunications	control selection and support for higher baud rates. LANGUAGES: Dynamic Memory Overlays for Turbo Pascal by Steve McMahon Steve shows how to get around the 64K limit on executable code imposed by Turbo Pascal without resorting to slow disk overlays. COMMUNICATIONS: A Unix BBS Using Shell Scripts by Jan L. Harrington Jan shows just how she wrote a bulletin board system
software	using Unix V Bourne shell scripts and XMODEM.
Queue control in C >	C CHEST by Allen Holub
For the programmer's bookshelf	Allen delves into the subject of priority queues and, in Flotsam and Jetsam, discusses standard #include files. 16-BIT SOFTWARE TOOLBOX by Ray Duncan Ray discusses programming books. He also offers a poor man's MAKE utility and an MS-DOS programming tip.
Object-oriented programming	ARTIFICIAL INTELLIGENCE by Ernest R. Tello Ernie presents some examples of object-oriented programming using SCOOPS, an extension of PC Scheme, "the Turbo Pascal of the PC LISP family."

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About the Cover

Handling the traffic flow of multitasking requires an efficient queueing algorithm such as the one that has just given the green light to a task in this month's cover illustration.

This Issue

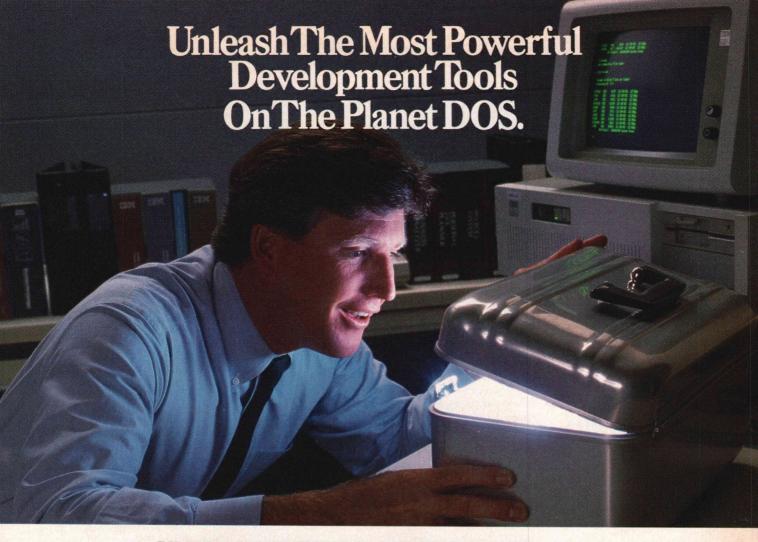
If this issue looks packed, it is. Three articles deal with communications techniques. Robert Jay Brown and Allen Holub discuss algorithms for priority queues. Dave Cortesi, a longtime DDJ columnist, contributes a rare bit on doubly linked lists. Steve McMahon gives aid and comfort to serious Turbo Pascal programmers; Ernie Tello and Allen Holub give tips to new PC Scheme and C programmers, respectively; and Forth, assemblylanguage, Modula-2, and BASIC programmers will all find something of interest in this issue.

Next Issue

So you really did it? You sold the Volvo and bought a 386 machine and now you're going to develop software for it? Only you don't want to wait for Microsoft to deliver OS/2 sometime next year? We hope you saved \$2.95 for our July issue, in which we will discuss development tools for the 386 that are available today.

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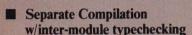
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CIRCLE 332 ON READER SERVICE CARD

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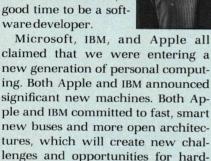
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CIRCLE 257 ON READER SERVICE CARD

FDITORIAL

The summer of 1987. Midyear. Time to take stock; time to assess what we've seen of 1987 and mull over what it all means What can we say? That it was the spring of hype, the summer of virtuality? I think it's a watershed summer, a good time to be a software developer.

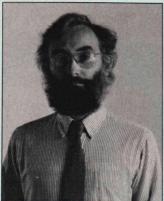


ware and software developers.

Microsoft revealed plans for getting parts of its OS/2 for 80286 and 80386 processors into developers' hands this year. Ever since Gary Kildall put an operating system on an 8080, the only arguably revolutionary advance in personal computer operating systems has been the Macintosh system software. It might be a stretch to suggest that OS/2 could be the next, but when it really arrives it should prove fertile ground for software development.

Both Unix and C took steps toward standardization in the first half of this year, and the crowded C compiler market ripened toward some sort of mitosis. Both Unix and C presented new tools and new environments for development. Version management and CASE became more real for personal computer-based development.

If you're reading this in Tokyo, you already know that several Japanese computer companies have taken a big step toward a viable Japaneseuser-oriented personal computer industry by standardizing on the TRON real-time operating nucleus. TRON should have a much greater impact on Japanese computer users than



the Fifth Generation Project, and it's also going to have a significant impact on software developers by providing an entirely new, coherent, Japanese-script-based programming platform.

These new environments and new tools speak of new opportunities for the software

developer. But they speak in the ambient din, and it's easy to miss their message.

If you're currently evaluating the software development opportunities presented by these recent advances. could I make some modest suggestions?

Don't listen too closely to the marketers. Marketing experts are magicians; when it's difficult to create value, marketers can create need. They can convince people that there is something terribly wrong with them that only your product can cure. But it's so easy to create real value with a piece of software that I'm just not sure that we need the digital deodorant. If the product is good, marketing it becomes a simpler matter.

And I wouldn't listen too closely to users, heretical as that statement may be. Listening to users doesn't generally contribute a lot to creating something new. User feedback is useful when you're tweaking existing products. Users may not know anything about software, but they know what they don't like.

Ultimately, I argue, the real breakthroughs come from listening to the technology. That should be particularly true at a technological watershed like this summer.

Hock the Volvo, buy that 68020 or 80386 machine, and start playing

Michael Swains Michael Swaine editor-in-chief

Dr. Dobb's Journal of ftware'

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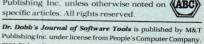
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RUNNING LIGHT

reetings, programmers. I'm Levi Thomas, assistant editor and host of DDJFORUM, our SIG on CompuServe. If my writing style seems conversational, it's be-

cause I'm accustomed to the SIG, on which my readers can talk back.

It's an interesting phenomenon, the on-line life. When I first began working for *DDJ* as electronic editor I felt like a DJ on the graveyard shift. Posting messages on the message board of the SIG, I got the eerie feeling that there was really nobody out there, that just maybe I was talking to myself. Luckily the FORUM members are a verbal bunch, and even the lurkers can be coaxed into responding if I ask the right questions or start a topic that gets their dander up.

Now I want to talk to you hardcopy readers. I wanna pull on yer coats about something. I hope you'll talk back to me, either on line or via the U.S. Snail.

<hopping onto soapbox>
Microcomputers have grown up.

We all know the story of the microcomputer's Wonder Bread Years. Born with features only a mother could love (front panel toggle switches, 2K of RAM, no permanent storage, and so on), they have matured into smart, fast machines.

The microcomputer community has also grown up. What started a few years ago as a hobby has become a business. The community has become an industry. For many of you, your fascination with computers has turned into a shot at The Dream Job: A chance to make a living at something you love to do. A chance to make work into play.

However, like many new adults, this industry seems determined to disown all of its childlike qualities as it seeks a grown-up identity. It's been giving up many of the values that fueled its growth. The Suits have moved in, and things are tightening up.

Examples: It's becoming next to impossible to get a job on pure ability; college degrees are the new prerequisite. The sense of community we once got from the computer magazines is disappearing. I see the term hacker changing-not evolving into the essence of what it originally meant but devolving into a derogatory term, a term feared by this new, "grown-up" industry. (Last year I was involved with the Hacker's Conference. Several computer businesses were nice enough to help fund the affair, but some of them asked that their names not be disclosed. They were nervous about having their company names associated with the word hacker.)

<returning soapbox to closet>

DDJ is still a reader-responsive magazine. If you would like to write for our October (Forth), November (graphics), or December (operating systems) issues, send us an outline. You can address your article proposals to our new editor, Tyler Sperry (you'll meet him on this page next month). If you just can't stand the thought of using the U.S. Sail, you can me on CompuServe (76703,4060), Usenet (well!levi), or arpanet (well!levi@lll-crg.arpa). We look forward to hearing from you.

Levi Thomas assistant editor

ARCHIVES

Inaugural Columns

"Welcome to the first edition of Programming Pastimes and Pleasures. Together we will explore the prosaic and fantastic possibilities of programming. Sometimes I will pose a problem to solve; sometimes I will discuss a programming technique or idea you may not have heard of; at other times I will tell you about a game you can play with your computer. At times I will be serious, and at other times I will be off-the-wall. Along the way, you will learn things you didn't know before, and you will find wonderful new ways to waste time with your computer. Most of all, I hope you will have fun." Charles Wetherell, DDJ, January 1979.

"Doctor Dobb's Clinic is a new venture. It is a place for the display of techniques and discoveries. We want to cover any sort of method or trick you've found in your exploration of (for example, but not limited to): CP/M or OASIS or FLEX or APPLE DOS or TRSDOS or NEWDOS or RSTS or UNIX or UCSD PASCAL or....

"During our rounds at the clinic we may examine a compiler or an interpreter, revealing its bugs or showing how it can be made to perform better. We'd like to tell of unobvious uses for standard utilities. We want to uncover errors in published documentation, to warn people away from pitfalls, and to show off those 'eureka!' moments that make systems work rewarding."—Dave Cortesi (Resident Intern), DDJ, May 1981.

'Just after watching Doc Dobb's nationally televised speech from the NCC in Texas, I was surprised to receive a collect call from the Old Man himself. 'The lack of public domain software for the 16-bit microcomputers is appalling!' remarked our Fearless Leader. 'No one ants to shell out that kind of money just so they can stare at the operating system's sign-on message. I want you to institute a regular column in DDJ that will address the needs of 16-bit system users and promote the discussion and interchange of software.' Luckily, after laving down that sweeping and rather alarming mandate, the good Doctor had to hang up to go out on a house call (someone had choked on an Apple) and I was left to my own devices."-Ray Duncan, DDJ, September 1982.



Microsoft Avoids Challenge

We challenged Microsoft to a C compiler duel-to-the-finish, comparing compile, link and execution times, and we offered to stop advertising for two months if they won...

by Roy Sherrill, President, Datalight

Microsoft purchased our C-compiler during February 1987 and we still haven't heard from them. OK, Microsoft, we are extending our challenge deadline from April 1, 1987 to May 15, 1987. After all, the Microsoft ad claims "the fastest C you've ever seen." Your reply, Microsoft!

Walter says Optimum-C is better

Walter Bright, the developer of Optimum C, says that Optimum C would win 7 out of 10 benchmarks as compared to Microsoft C, V.4.0. Walter explained to me that Optimum C includes a unique global optimizer that helps create compact code while increasing execution speed up to 30%. By the way, Borland, Walter is still waiting for his copy of Turbo C. V.1.0. Borland's ad claims "the fastest, most efficient and easy-to-use C compiler at any price.'

After reviewing Borland's benchmarks, Walter claims that Optimum C is faster. And, as for ease of use, all Datalight C compilers have been shipped with a free Learn C program for the last six months. Also, our new EZ Interactive Editor will show you each syntax error in your source code, then compile or "make" and run your program, all from within the editor. OK, so let the Microsoft challenge begin...

We only ask the following...

The benchmark suite will consist of the set of programs that Microsoft supplied to Computer Language for their February 1987 C compiler review issue. Microsoft will make available the programs to Datalight at least two weeks prior to the benchmarking. The benchmarking will be between Microsoft C 4.0 and Optimum-C. It will occur at a mutually agreed upon time and place. Interested individuals will be allowed to attend. The benchmarks will be compiled and run on a standard IBM PC-AT.

There will be two separate tests for each program: compile and link speed, and execution speed. For each test, a representative from each company will set up the compiler so that it performs at its best.

The benchmarks will be adjusted so that they take sufficiently long to run, that the tolerance involved in timing them is insignificant. The winner is determined by the compiler with the faster execution times for the majority of the benchmarks. We'd like an answer from Microsoft no later than May 15, 1987.

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DR. DOBBS, August 1986

"This is a sharp compiler!... what is impressive is that Datalight not only stole the compile time show completely, but had the fastest Fibonacci executable time and had excellent object file sizes to boot!"

COMPUTER LANGUAGE, February 1986

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- pr-Page printer
- pwd—Print working directory
- wc-Word count

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LETTERS



What's Right with High-Level Languages?

Dear DDJ,

Mike Suman's criticisms of Modula-2 in the February 1987 Viewpoint call for some comments.

First, although many Modula-2 implementations are now available, based on Wirth's *Programming in Modula-2* (either the second or third edition), both the language specifications and the preferred standard libraries are still being fine-tuned by such bodies as M2WG, the Modula-2 Working Group of the British Standards Institution. As I understand the situation, draft proposals from BSI

now await reaction and further input via the ISO. MODUS Quarterly (published by the Modula-2 Users' Association) has been covering this debate and provides a "democratic" forum for suggestions, including Wirth's own reactions. Mike should throw in his penn'orth before the stonecutters start chiseling. Contact George Symons, MODUS, P.O. Box 51778, Palo Alto, CA 94303; (415) 322-0547.

Second, I must defend Brian Anderson's use of the identifiers FIRST and LAST in X68000. Although in the final, published version these turn out to be nonce constants, during the development stage (and later when the MC68030 is incorporated!) who knows what values they might assume or where else they might occur?

Finally, you must really sympathize with Niklaus. You

want to keep a language "simple," "portable," and "small-core," yet hardly is the ink dry before committees emerge wanting to add their pet piece of FORTRAN or Ada. Where do you draw the line? Certainly array and record constants would be useful (as in Turbo but not standard Pascal) and may find their way into the canon. Modula-2 does have the useful, non-Pascal *Module Body* construct (a piece of code invoked just once when the module is first executed).

Mike's preferred tabular layout for setting up Brian Anderson's *Table68K* will be possible under my own forthcoming language called Modula 1-2-3.

Stan Kelly-Bootle 25 Parkwood Ave. Mill Valley, CA 94063

Stan Kelly-Bootle is the author of The Modula-2 Primer. —eds.

Dear DDJ,

In his February 1987 Viewpoint, Mike Suman raises a couple of objections that he suggests show something wrong with Modula-2 and other high-level languages. I don't think his objections are cogent, and I don't think they have anything to do with high-level languages.

His first complaint concerns the use of constants to replace Arabic integers. This is not peculiar to highlevel languages but is a feature of most assembly languages, too, so whatever is wrong, it can't be a problem of Modula-2 or of high-level languages in general.

His particular complaint concerns the use of *FIRST* and *LAST* to replace 1 and 118, respectively; in the program on which he is commenting, they appear in the declaration:

Table68K : ARRAY[FIRST..LAST] OF TableRecord:

This does look odd, but the usual use of constants in this situation would omit the constant *FIRST* and replace the constant *LAST* by something more informative, such as *Number-OfOpcodes*, so the declaration would look like this:

Table68K: ARRAY[1..NumberOfOpcodes] OF TableRecord;

> I can't see why he would object to that. He does say, "no one can maintain that it is really easier, or safer, to change values in an early definition than it is to change them in the only place in which they are used later, when it is obvious what the effects of the change are going to be." But six months and five revisions later, how will you know that the number of opcodes is used only once in the program, and how easy will it be to find that one place?

His larger complaint is that Modula-2 contains no compact and readable way to initialize a complex table. But then, what language does?

There are really two points mixed up here. One is that Modula-2 provides no way to initialize variables other than by assignment statements; there are no facilities for de-



Architecture of Wendin's Concurrent I/O System

Wendin's Operating System Toolbox makes it possible to write a concurrent driver for any I/O device.

Wendin's Operating System Toolbox kernel is composed of three major components: the scheduler, the memory manager, and the I/O subsystem. This month we will show how the I/O subsystem works, at the QIO (Queued Input/Output) device driver level.

The I/O system supports concurrent and waited I/O operations from multiple processes simultaneously. Its architecture was derived from the VAX/VMS QIO system, which was developed by Digital Equipment Corporation with these ideas in mind. The whole QIO mechanism takes advantage of many of the kernel's capabilities, including ASTs (Asynchronous System Traps), and event flags (which are used to signal I/O completion).

Before a program can perform I/O operations on a device, it must assign a channel to the device. A channel is simply a number, similar to an MS-DOS file handle. The operating system uses an assigned channel number as an index into an array of channel control blocks (CCBs) to get information about the channel and the corresponding device.

After assigning a channel with the SYS__ASSIGN system service, the user program can make calls to SYS _QIO or SYS__QIOW (for concurrent or waited I/O) through the INT FF software interrupt. The kernel calls the routine EXEQIO, which validates the caller's parameters and I/O privileges. If the call parameters are valid and the channel number corresponds to an assigned CCB, EXE-QIO will deliver a KERNEL mode AST to the routine ASTQIO in the same process context. This allows the driver to run on its own stack, but in the context of the calling process. ASTQIO receives all information about the I/O operation in an IRP (I/O Request Packet), created by EXEQIO from the user program's parameters.

```
void ast_qio (packet)
   struct irp *packet;
{
   struct ccb *c;
   struct ucb *u;
   void (*driver)(struct irp *);

   c = &io_chantbl [packet->chan];
   u = &io_unittbl [c->unit];
   driver = u->driver;
   (*driver)(packet);
} /* ast_qio */
```

Listing 1. How ASTQIO dispatches the device driver given a channel number.

ASTQIO uses the channel number to find the corresponding CCB, and obtain a pointer to a data structure describing the device assigned to, in the format of a unit control block (UCB). The unit number is used as an index into an array of UCBs, much as the channel number is used as an index into an array of CCBs. The UCB contains information about a specific device, including what general type of device is being referenced (for example, a terminal, mailbox, or disk), and a pointer to the device driver code. This double indirection (the channel obtains a CCB, which in turn references a UCB, which finally references the device driver) allows several channels to be assigned to the same device at the same time, and for any number of UCBs to share the same device driver code (for example, the disk driver defines a UCB for each disk drive, but there is only one device driver for all of them).

Once ASTQIO has obtained the UCB for the requested device, it calls the device driver directly. The device driver is responsible for determining what to do based on the I/O function given in the IRP. The device driver will typically call the system BIOS, or interface directly to the hardware, to perform an I/O operation.

Finally, the kernel must signal the user process that the I/O is done. The SYS__QIO services allow two

methods for this: an event flag can be set, or an AST can be executed when the operation finishes. The kernel routine STOPIO examines the IRP associated with an operation, and determines which event flag to set, and whether to queue an AST to the user process. Finally, the IRP is removed from the system I/O queues, and returned to the memory pool. Control is then passed back up the chain and returns to the user program.

A device driver can be written to talk to any device, based on standard models provided by the Operating System Toolbox kernel. Examples include mailboxes (functionally equivalent to UNIX pipes), terminals, disk drives, optical disks, or even expanded or extended memory. Basically, the minimum I/O functions that the driver needs to support are the read and write block functions. For some devices (like disks) there will be a distinction between virtual, logical, and physical I/O. For others (like mailboxes), all of these functions will perform the same operation. You can decide how you want a device driver to work, and it's easy to integrate a driver into the toolbox kernel.

Next month, we'll continue with the internals of the toolbox. If you'd like to learn more about operating system design and architecture, pick up a copy of Operating System Toolbox from Wendin and follow along.

Operating System Toolbox: \$99 Wendin, Inc. P.O. Box 3888 Spokane, WA 99220 (509) 624-8088 fining initial values at the point of declaration as in some other languages (although the compiler I use provides this as an extension). Unless the compiler is very clever, this may mean some unnecessary work at execution time, but this does not seem like a major problem and certainly not one inherent in high-level languages.

The second point is that it is tedious to initialize a large table (in the example Suman discusses, it consists of 118 records of 4 fields each) with assignment statements. But there are other ways. For instance, I like to write a table in a readable form in an ASCII disk file and read it at run time to initialize the table. It's hard to beat this for readability, the table can be edited without recompiling the program, and the code to read the table is usually not difficult to write. There is a problem with reading in enumerated types from an ASCII file: because it can be a nuisance to translate the ASCII version of the identifiers for the values, you can compromise by using Arabic numbers for their ordinal values instead. This can be endured because the table file can easily contain a key for the values.

Suman offers an assembly-language version of the table in which sets are initialized by 16-digit binary numbers. I suppose it's a matter of taste, but I don't find 118 entries like that very readable, and I'm sure I'd make plenty of errors typing them into my program. And what would he do if he had a table with real numbers in it?

I suppose it would be nice if we could write our tables in readable tabular form in our code and have the compiler do the work, but there are lots of tasks we might like our compilers to perform, and we can't build them all in. I would guess that initializing complex tables is one of those tasks that is just as well left out of general-purpose languages.

John G. Bennett 301 Roslyn St. Rochester, NY 14619-1813

Dear DDJ,

Though the Viewpoint by Mike Suman is interesting and raises some valid points, I find it to be misleading

in three specific ways.

First, the point about the constants FIRST and LAST being used only once is not quite right. Consider the index variable i. The initialization and manipulation of i should always be in the range FIRST . . LAST, and further, it should exactly cover this range. The program is poorly written in that the index should have been initialized to FIRST, not 1 explicitly. Next, the value of the index at the end of the initialization should have been checked to see that it was equal to LAST, so that you know all the elements of the table have been initialized. Far from pointing out a place in which Modula-2 is too strict, this points out a place in which the programmer just didn't use the language properly.

Second, the fact that Modula-2 doesn't allow the specification of constant arrays is not an indication that Modula-2's type checking is too strict or that its model of programming is inadequate. It means that Modula-2 doesn't have this feature. It is clear that allowing constant aggregates to be uttered in the language is an independent issue from whether the language is inconveniently strict or not. Strictness and features provided are two completely different issues.

Third, asking if imagined (or even real) problems with the Modula-2 language mean there is "something wrong with the direction in which we are being led" is to miss the point altogether. The point of "higher-level" languages that have a strict type structure is not to prevent programmers from full expression or to make it harder to write legal programs. The point is to help programmers by making it easier for them to tell when they have written a meaningless or illegal program. Here we see that a programmer misused the constant definition facility of the language and that the language lacks the ability to utter constant-valued aggregates. But neither of these faults is a fault of strict type checking, nor of modularity, nor of the high-level nature of the language.

So, if we are being led in the direction of modularity and static type safety, Mike hasn't even begun to present evidence that the direction is wrong. He has shown that small languages often lack features that are awkward to do without, and he has shown that people often forget that index variables ought to be initialized and manipulated using the same bounds that were used to declare the array they index into. These are indeed things to be wary of, and they show that Modula-2 is neither omnipowerful nor errorproof. But this is no reason to throw away the baby of type checking and modularity with the bathwater of a possibly too small language and human error.

Wayne Throop 86 Fearrington Pittsboro, NC 27312

Happy Ducks

Dear DDJ.

Thank you for the very funny February cover revealing the true nature of us WordStar-imprinted programmers. But no thanks for the text editors article. For a *DDJ* feature article, it was amazingly uninformative—simply an excuse for a two-barrel discharge against us happy ducks.

Especially unfair was the statement that WordStar-style editors "usually use 'weird' file formats that can't be read by any other editor without some sort of conversion." WordStar itself works in straight ASCII when in its program-editing (non-document) mode. But that is indeed not the most usual WordStar-style editor. The MUWSSE is, of course, the Turbo Pascal editor, which many waterfowl use with other compilers as well. And that is nothing but straight ASCII. So, just what were you talking about?

Another gross misstatement was the wish list. Every programmer I know, even strictly dry land, will tell you that wish number 1 is speed, 2 is speed, and 3 is speed. Other wishes appear only after pausing for breath. No mention of that in your list. Of course, again, the MUWSSE—the Turbo editor—is the fastest thing under ten fingers. And for the price of just about any dedicated programming editor, you can get Turbo Pascal and SuperKey, and then you have the

(continued on page 122)

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BASIC

Index

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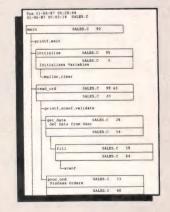
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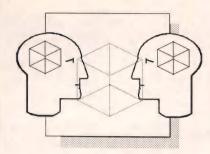
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VIEWPOINT



What's Right with High-Level Languages?

In "What's Wrong with High-Level Languages" (Viewpoint, February 1987), Mike Suman brings up some interesting points about the limitations inherent in any computer language but fails to mention any of the strengths of Modula-2 or of modern high-level languages in general.

As a college instructor who has spent several years developing and teaching a course in assembly-language programming, I can well appreciate the advantages of low-level programming. Assembly language has some shortcomings that can be much more serious than the limitations of high-level languages, however.

Assembly language is much harder to learn than are most high-level languages, and you must substantially relearn it if you move from one computer to another. I grant that the second assembly language is much easier to learn than the first. It would, however, take a considerable effort to adapt from the 68000 to the

by Brian R. Anderson

8086, for instance, especially when you consider the peculiarities of Microsoft's MASM.

Assembly language is hard to write and even harder to debug (compared to Modula-2), partially because often the only control structures available

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are the conditional and unconditional jump and the call to subroutine. To test several conditions for a loop, you must write a cascade of comparisons and jumps. To produce a simple repetition, you must split the control instructions between the outside of the loop, the top of the loop, and the bottom of the loop. Consider Examples 1 and 2, below, in which some elements of a character array are cleared using 68000 assembly language and Modula-2, respectively. In both examples, the range of items to be cleared has been calculated and left in variables named TOP and BOT. In the assembly-language example, the loop constraints are developed in five different lines of code, whereas

in the high-level language example, all loop constraints are developed on a single line.

Modern high-level languages allow programmers to express data in terms that naturally match a wide range of typical problems. Modula-2 handles the mathematical concepts of real and integral numeric types and sets, the organizational concepts of records and files, and the universal concepts of arrays and characters. In assembly language, the only data type is the computer word. Programmers must impose a structure and then provide algorithms to perform even the simplest task. Consider how much easier it is (for exam-(continued on page 124)

```
FOR loop in 68000
                                 First Array Element to Process: TOP
                                 Last Array Element to Process: BOT
                                 The array: DATA
                                 Used to Index into array: DO
                                 Used as Pointer to array: AO
000000
         00
                             DATA
                                              500
                                                              ;array of bytes
0001F4
         0000
                             TOP
                                      DC
                                              0
                                                              ; first to process
0001F6
         0000
                             BOT
                                      DC
                                              0
                                                              ; last to process
                                                              ;other data
                                                              ; and code
                                Assumes TOP & BOT previously calculated
0001FB
         41F900000000
                                     TEA
                                              DATA, AO
                                                              ;set pointer
0001FE
         3039000001F4
                                     MOVE.
                                              TOP, DO
                                                              ;init index
000204
         B079000001F6
                            LOOP
                                     CMP
                                              BOT, DO
                                                              ; check bounds
00020A
         6206
                                     BHI.S
                                              LABEL
                                                              ;end loop
0020C
         42300000
                                     CLR.B
                                              0 (AO, DO)
                                                              :loop body
000210
         60F2
                                     BRA.S
                                              LOOP
                                                              ;repeat
                            LABEL
000212
                                     END
```

Example 1: Clearing elements of a character array in 68000 assembly language

```
VAR
DATA: ARRAY [1..500] OF CHAR;
TOP, BOT, i: CARDINAL;
BEGIN
(* Assumes TOP & BOT have been previously calculated *)

FOR i:= TOP TO BOT DO
DATA[i]:= OC;
END;

END TRIAL.
```

Example 2: Modula-2 version of the code in Example 1

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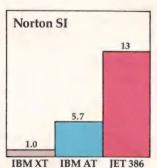
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CIRCLE 130 ON READER SERVICE CARD

An Efficient Algorithm for Large Priority Queues

by Robert Jay Brown



n many real-time programs, it is necessary to share resources that are in short supply. To help manage these situations, a priority system is often established. When several contenders are competing for the same resource, the one

The central concept is that inserting and removing from the queue can be viewed as merging two separate queues.

with the highest priority gets the resource. When the resource again becomes available, the next-priority contender gets it. A simple first-in/first-out queue can be thought of as a priority queue in which the time a contender is enqueued is the priority and lower numbers have higher priority.

In practice, the resource being waited on may be a physical device, such as a printer; a logical device, such as a file or a record of a file; or the CPU itself. Alternatively, a timer scheduler may be implemented by using the desired dispatch time as the priority and having the dequeueing operation wait until the time of day equals the dispatch time at the head of the queue.

Basic Queue Operations

A priority queueing scheme in a real-time system must be

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able to perform the following operations on the elements, or nodes, of the queue: determine the highest-priority node, add a new node, and remove a node. Removing the highest-priority node is a special case of the more general operation

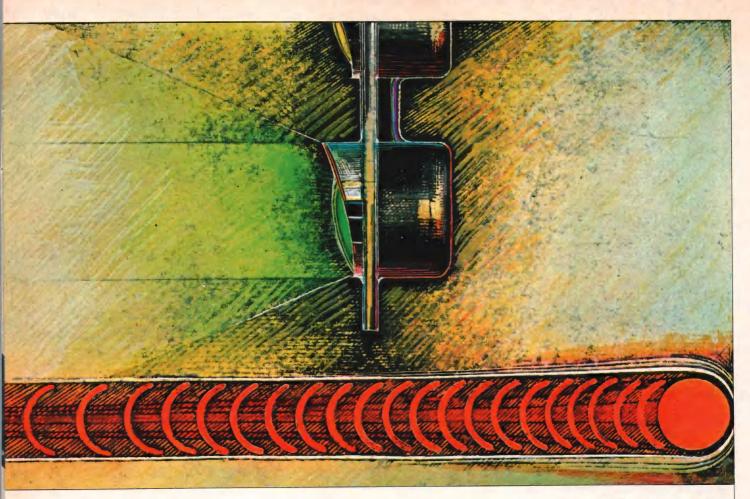
of removing a node anywhere in the queue. This is called preempting.

The most simple implementation for priority queues results in the time to perform at least one of the above operations being directly proportional to the size of the queue. For extremely large queues, this becomes unworkable. The problem is similar to sorting, and sorting can be done in a time proportional to the logarithm of the number of elements, so you should be able to do as well for a priority queue.

In fact, you can do a bit better than this. Although the priority queue algorithm I have implemented (see Example 1, page 18) performs in logarithmic time, it does not keep the queue completely sorted. It defers determining the second-priority element until after the top-priority element is removed. This does not change the logarithmic component, but it makes each iteration go faster.

A Binary Tree

The representation for the queue is a binary tree with left, rite, and father connections. Each element also con-



tains a sort *key*, which can be interpreted as the priority. In addition, a *dist* cell is used to indicate the minimum distance from that node to a leaf, which is a connection to no element, or a terminator. The *father* connection is not used to maintain the ordering of the queue but is used to allow rapid removal of any node in the queue (see the example, lines 28–37). (Note: my structure declaring words [lines 3–26], together with examples of their use, are available on the East Coast Forth Board as the file STRUC.ARC, and so I will not describe them further here.)

The word and data cells (lines 36–37) are used to contain dispatching information. Word is the address of a Forth word to perform when the node is dispatched, and data is a word of data, typically a pointer, that is pushed on the stack before performing the word.

The following priority queueing algorithm was first described by Clark Allen Crane in 1971. My implementation is an extension of a revised version of Crane's algorithm that is described by Donald Knuth. You can refer to Knuth's *The Art of Computer Programming: Volume 3, Sorting and Searching* (Reading, Mass.: Addison-Wesley, 1973) for a complete description of the algorithm and an analysis of its performance.

Merging Queues

The central concept behind Crane's algorithm is that the operations of inserting and removing an element from the queue can be viewed as merging two separate queues. Crane's algorithm keeps the highest-priority node at the root of the tree, and the subtrees follow the

same pattern. Thus, to dispatch the head element of the queue (lines 103-106), the left and right subtrees are pruned from the root and merged back together, leaving the root out of the result. To insert an element on the queue (lines 98-101), that element is viewed as a little priority queue, in its own right, of one element. This queue is merged with the original queue, and the element is thereby inserted into the original queue. To remove a node from somewhere in the middle of the queue-that is, to preempt it (lines 108-111), the node's left and rite subtrees are cut off and merged, forming an intermediate result. Next, the preempted node is removed by using its father pointer to cut it off from the element that points to it. Finally, the original queue, less the preempted node and its two subtrees, is merged with the intermediate tree described above.

The implementation in the example has been tested on a 10-MHz 80286 with one wait state on memory access and 3.2 percent DRAM refresh interference running under LMI PC/FORTH+ 3.1. For 10,000 iterations of an *unque* on a randomly selected element of the queue, followed by an *enque* using a randomly chosen priority, it produced the results shown in Table 1, page 18.

Availability

The complete source code, as a standard Forth screen file, including the benchmark test, is available on Compuserve (GODDJ), and on the East Coast Forth Board ([703] 442-8695). Also, all the source code for articles in this issue is available on a single disk. To order, send \$14.95 to *Dr*.

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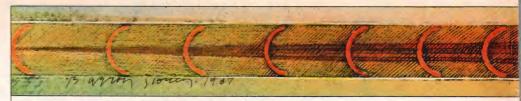


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QUEUEING ALGORITHM (continued from page 17)

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#nodes	msecs
10	7.4
20	9.3
50	11.9
100	13.8
200	15.5
500	17.9
1000	19.4
2000	20.7
5000	21.9
10000	23.1

Table 1: Benchmark test results

```
( Timer queue implementation
                                               rjb 16:06 07/2186 )
   ( ----- structure declaring words -----)
   : <struc> CREATE , DOES> @ + ; ( 2nd generation defining word )
 7
   ( 'struc' is used to create the defining word for a structure )
   : struc CREATE ,
                      ( org struc <struc> ; creates a structure )
 9
           DOES>
                   ( size <struc> <element> ; creates an element )
10
           DUP @ DUP >R ROT + SWAP !
                                     ( update location counter )
11
           R> <struc> ;
                                         ( make word for element )
12
13 ( 'array' is used to create an array of structures )
14 : array CREATE ( #slots slot-size array <array> ; makes array )
15
           DUP , * ALLOT
                             ( save slot-size and allocate array )
16
           DOES>
                          ( subscript <array> -- &<array>element )
17
           DUP @ ROT * + WSIZE + ; ( return pointer to the slot )
18
19 ( 'org' is useful for doing the 4th version of a C 'union')
20 : org
           ( n org <strucname> ; re-initializes location counter )
21
           BL WORD FIND NOT ABORT" Undefined! " >BODY !;
22
23
  ( 'sizeof' is used for declaring structures of structures )
  : sizeof ( sizeof <strucname> -- n ; gets the size of a struc )
           BL WORD FIND NOT ABORT" Undefined " >BODY @
           STATE @ IF [COMPILE] LITERAL THEN ; IMMEDIATE
26
27
28
             ----- timer queue entry -----)
29
30
           struc
                 tq
                                     ( a node in the timer queue )
31
           4
                  tq
                                   ( the sort key: dispatch time )
32
           4
                   tq
                          dist
                                      ( distance to nearest leaf )
33
                   tq
                          left.
                                      ( pointer to left sub-tree )
34
           4
                  ta
                          rite
                                     ( pointer to right sub-tree )
35
           4
                  tq
                          father
                                     ( pointer to father of node )
36
           4
                   ta
                          word
                                   ( word to execute at dispatch )
37
                  tq
                          data
                                      ( pointer to data for word )
38
39
       ----- timer queue implementation -----)
40
41
  ( left! & rite! set the left and rite subtrees, respectively
42
    of a node, called the father. The father pointer of the son
43
    node is updated to point to the pointer, left or rite, that
44
    points to the son, so that it may be cleared on an unque. )
45
46 : left! DUP 0<> IF OVER left OVER father !
                                                ( Father Son -- )
```

Example 1: Timer queue implementation



```
THEN SWAP left ! ;
47
48
49 : rite! DUP O<> IF OVER rite OVER father!
                                                  ( Father Son -- )
           THEN SWAP rite ! ;
50
52 : go-rt DUP rite @ >R DUP ROT rite! R> ;
                                                ( go do rite side )
53
54 : dist@ DUP IF dist @ THEN ;
                                        ( P -- P=nil ? 0 : dist @ )
56 (This is step M2: from Knuth p 619, sol'n to prob. 32, p 159)
                        (PQR -- PQRD; merge priques P & Q, )
57 : list-merge
        BEGIN ( R is Roving pointer, D is Distance to near leaf )
58
           OVER 0= IF
                                                      (0 = nil ?)
59
                2 PICK dist@ EXIT THEN ( yes, D = P->dist; done! )
60
                                                      (P = nil ?)
61
           2 PICK 0= IF
               ROT DROP OVER SWAP
                                                     ( yes, P = Q )
62
                2 PICK dist@ EXIT THEN ( yes, D = P->dist; done!
63
           2 PICK key @ 2 PICK key @
                                              ( P->key < Q->key ?
64
           < IF ROT go-rt ROT ROT ELSE
                                             ( yes, P moves right )
65
                SWAP go-rt SWAP THEN
                                             ( no, Q moves right )
66
                      ( loop until one of the trees is eliminated )
67
       AGAIN :
68
69 (This is steps M3: and M4: from Knuth p 619)
   : fix-dist ( P Q R D -- P ; fixes up distance to nearest leaf )
70
      REGIN
71
                                          (R = nil ? yes, done!)
         OVER 0= IF 3DROP EXIT THEN
72
73
         ROT DROP OVER rite @ ROT ROT
                                                    (Q = R - > rite)
74
         OVER left @ dist@ OVER < IF
                                            (R->left->dist < D?)
                                          (D = R \rightarrow left \rightarrow dist + 1)
75
            DROP DUP left @ dist@ 1+
                                              ( R->rite = R->left )
            OVER DUP left @ rite!
76
            OVER 4 PICK left!
                                  ELSE
                                                    (R->left = P)
77
                                               ( D++; R->rite = P )
78
            1+ OVER 4 PICK rite! THEN
                                                    (R->dist = D)
         DUP 2 PICK dist !
79
         >R ROT DROP SWAP DUP R>
                                                   (P = R; R = Q)
80
                                   ( spin down the right sub-tree )
81
      AGAIN ;
82
83 (tq-merge is Knuth's Algorithm M from p 619)
84 : tg-merge 0 list-merge fix-dist; (PQ -- P; merge 2 tq's)
85
                                               VARIABLE TO 0 TO !
86 (the timer queue root pointer)
87
88 ( TQ! updates the father pointer in the first node, & sets TQ )
                               ( tq -- ; sets TQ and father of tq )
89 : TQ!
                                       ( set the timer queue head )
 90
                                              ( empty? yes, done! )
            DUP 0= IF DROP EXIT THEN
 91
                                     ( no, set father of top node )
 92
            father TQ SWAP ! ;
93
                                                      ( tq -- flag )
 94 : ?waiting DUP father @ @ = ;
95
 96 ( ----- timer queue package entry points -----)
97
                                        ( &tq -- ; enques to timer )
 98 : tq-enque
                                         ( nullify both sub-trees )
 99
           O OVER left ! O OVER rite !
                                        ( distance to a leaf is 1 )
           1 OVER dist !
100
                                   ( merge new node with old queue )
           TQ @ tq-merge TQ! ;
101
102
                                      ( -- &tq ; deques from timer )
103 : tq-deque
           TQ @ DUP 0= IF EXIT THEN ( returns nil on empty queue )
104
                                           ( save head for answer )
105
           DUP
           DUP left @ SWAP rite @ tq-merg TQ! ;
                                                  ( merge remains )
106
107
                                   &tq -- &tq ; removes from timer )
108 : tq-unque
           DUP ?waiting NOT IF EXIT THEN 0 OVER father @ !
                                                            ( cut
109
           DUP left @ OVER rite @ TQ @ ( both subtrees of tq & TQ
110
                                          ( paste it back together )
           tq-merge tq-merge TQ!;
111
```

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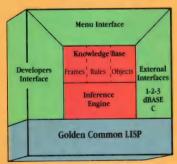
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GOLDHILL

Two-Bit Analog-to-Digital Conversion

by John Musselman

his article describes a simple technique for measuring an analog voltage that is not only useful but also demonstrates some programming techniques commonly used in real-time applications. The key to such programming is the use of hardware interrupts. A lot of programmers might be wary of hardware interrupts, but once you use them you will find they are a powerful tool. I work with embedded controllers a lot, and so I use them all the time.

Using Hardware Interrupts

The simplest interrupt scheme utilizes a hardware timer to generate an interrupt at regular intervals. How this is done depends on the hardware in the system, so it is not possible to go into much detail. In general, the idea is to set a timer to its free-running mode. In this mode, the timer is set to some value that decrements on each system clock pulse. When the timer reaches zero, it generates an interrupt, is automatically reloaded, and starts to decrement again. Thus, the interrupt routine is entered at regular intervals. This creates a time base for any routines running in the interrupt and guarantees that each routine in the interrupt will be executed within a certain interval. I almost always have one such interrupt running in any system I design. Many routines that might otherwise qualify for their own interrupt can often

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An example of putting interrupts to work and how software can directly interact with hardware

run in this ''master'' interrupt, which helps simplify things.

The driver for the analog-to-digital (A-to-D) circuit appears in this kind of interrupt. The frequency of the interrupt is not critical, although the faster it is, the more quickly or more accurately the conversion can be made. Once every millisecond (1,000 times a second) is a convenient figure.

Figure 1, below, is a simplified schematic of the A-to-D converter circuit. Just two bits—one output and one input—interface the circuit to the computer. To understand the circuit, first consider the part of it represented by the output bit, the resistor,

and the capacitor. This is actually a simple D-to-A converter. The output bit is set high or low during each interrupt. The resistor and capacitor values are relatively large, so the voltage on the capacitor is an average of the output voltage over time. More 1s and fewer 0s produce a higher voltage; fewer 1s and more 0s give a lower voltage. If the ratio of 1s to 0s is held constant over time, the voltage will be essentially constant and proportional to the percentage of 1 bits being output.

A simple way, then, to program a D-to-A with one output bit, a resistor, and a capacitor would be to repeatedly output m 0s followed by n 1s, where m+n is a constant. This could be programmed with a software counter that decrements on every interrupt. Comparing the counter against some variable determines whether to output a high or a low during that particular interrupt. When the counter reaches 0, it is reset to m+n. The variable thus would control the output voltage as the variable ranges from 0 to m+n.

The A-to-D converter does not use

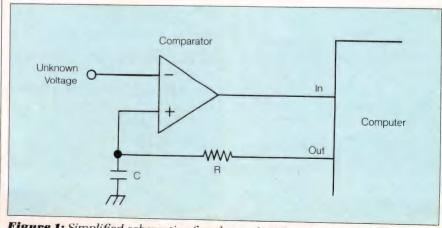


Figure 1: Simplified schematic of analog-to-digital converter circuit

exactly this method to create its output voltage, but the concept is the same. A counter is used to control the conversion cycle, and the ratio of 1 bits to 0 bits output during each cycle determines the output voltage. The difference is that the output is not all 0 bits followed by all 1 bits; the 1s and 0s can appear in any order.

In Figure 1, notice that the unknown analog voltage is fed into one input of the comparator. The other input of the comparator is the voltage generated by the computer, as explained earlier. The output of the comparator is a digital signal that indicates whether the unknown voltage or the generated voltage is higher.

The Software Driver

The A-to-D software driver samples the comparator output at each interrupt. If the generated voltage is lower than the unknown voltage, a 1 is output, charging the capacitor to a slightly higher voltage by the next interrupt. If the generated voltage is higher than the unknown voltage, a 0 is output, discharging the capacitor to a slightly lower voltage. In this way, the generated voltage seeks the level of the unknown voltage and then hovers about it. With large resistor and capacitor values, the searching effect is small and the two voltages are essentially equal. You can determine what voltage you are outputting by counting the percentage of 1 bits; this value is the answer sought.

Example 1, right, shows how you do this. The code is for a TMS7000 series processor. It is fairly easy to read, even if you are not familar with this device, but let me point out a few things. First, the MOVD instruction is a 16-bit move. The second byte of the variable is specified as the operand for this instruction, as in MOVD X+1,Y+1. Because the processor does not have a 16-bit increment instruction, I have used a 16-bit decrement to count in one's complement. The count is complemented before being saved as the result.

The conversion cycle takes 1,000 1-millisecond interrupts, or 1 second. A 16-bit counter, *COUNT*, keeps track of this time period. Another 16-bit counter, *HIGH*, keeps track of the number of 1 bits output during the cycle. At the end of the cycle, *HIGH* is

```
INBIT.
            EOU
                                                :INPUT BIT POSITION
                        2
                                                COUTPUT BIT POSITION
OUTBIT:
            EOU
                                                . OF INTERRUPTS IN A CONVERSION
                        1000
PERIOD:
            FOU
                                                COUNT OF HIGH BITS
HIGH:
            DS
                        2
                                                CONVERSION CYCLE COUNTER
COUNT:
                                                RESULT OF CONVERSION
RESULT:
            DS
                        2
            : INITIALIZATION ...
              MUST BE INCLUDED IN THE SYSTEM INITIALIZATION BEFORE INTERRUPTS
                        #PERIOD-1, COUNT+1
            MOVD
             INTERRUPT ROUTINES...
              THE FOLLOWING ROUTINES MUST APPEAR
              IN AN INTERRUPT WHICH OCCURS
              AT REGULAR INTERVALS
              1/0...
              SEE IF INPUT BIT IS HIGH OR LOW ...
                         APORT. B
 ATOD:
             MOVP
                         #INBIT, B, HI
             BTJO
               IF IT'S LOW, GENERATED VOLTAGE IS BELOW THE UNKNOWN VOLTAGE. OUTPUT A HIGH. COUNT ONE MORE HIGH BIT...
             ORP
                         #OUTBIT, APORT
 LO:
                                                 : NOTE ONES COMPLEMENT COUNT
             DECD
                         HIGH+1
                         TODONE
             JMP
               IF IT'S HIGH, GENERATED VOLTAGE IS
               ABOVE THE UNKNOWN VOLTAGE. OUTPUT A
               TOW ...
                         $255-OUTBIT, APORT
             ANDP
 HI:
 TODONE :
                CONVERSION CYCLE ...
              ; SEE IF CONVERSION CYCLE DONE ...
                          COUNT+1
              DECD
                          ATODDN
              JC
                IF CONVERSION DONE, SAVE RESULT
                AND RESET COUNTERS . . .
                                                  RESULT IS ONES COMPLEMENT
                          HIGH+1,B
              MOVD
                                                              OF COUNT
               COM
               COM
                          B, RESULT+1
               MOVD
                                                   ONES COMPLEMENT OF ZERO
                           #-1, HIGH+1
               MOVD
                           *PERIOD-1, COUNT+1
               MOVD
   ATODDN:
```

Example 1: Analog-to-digital converter driver

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TWO BIT A TO D (continued from page 23)

dumped to RESULT, and both counters are reinitialized. RESULT can be read by the main routine at any time it is desired to know the input analog voltage. It updates once a second. All this is transparent to the main routine.

The number generated by the A-to-D routine, RESULT, can take any value from 0 to 1,000. This means that there are 1,001 possible values-1 more than the number of interrupts in the conversion cycle. The actual voltage that RESULT represents is (RESULT/ 1,000) times the voltage swing of the output bit. If you actually build this circuit, be careful to measure this voltage. A TTL output may not provide the full 5 volts that an NMOS or CMOS output will.

Also, there is nothing magic about the number 1,000. The conversion cycle may consist of any number of interrupts. The more interrupts used in a conversion, the greater the resolution but the longer it takes to get the result. You can choose the number of interrupts to suit the application.

This circuit is a good example of how to put interrupts to work and how software can interact with hardware in a direct manner. Using software in a feedback loop, as in this example, is one of the best ways to convert from the analog to the digital world and vice versa.

Availability

All the source code for articles in this issue is available on a single disk. To order, send \$14.95 to Dr. Dobb's Journal, 501 Galveston Dr., Redwood City, CA 94063 or call (415) 366-3600 ext. 216. Please specify the issue number and format (MS-DOS, Macintosh, Kaypro).

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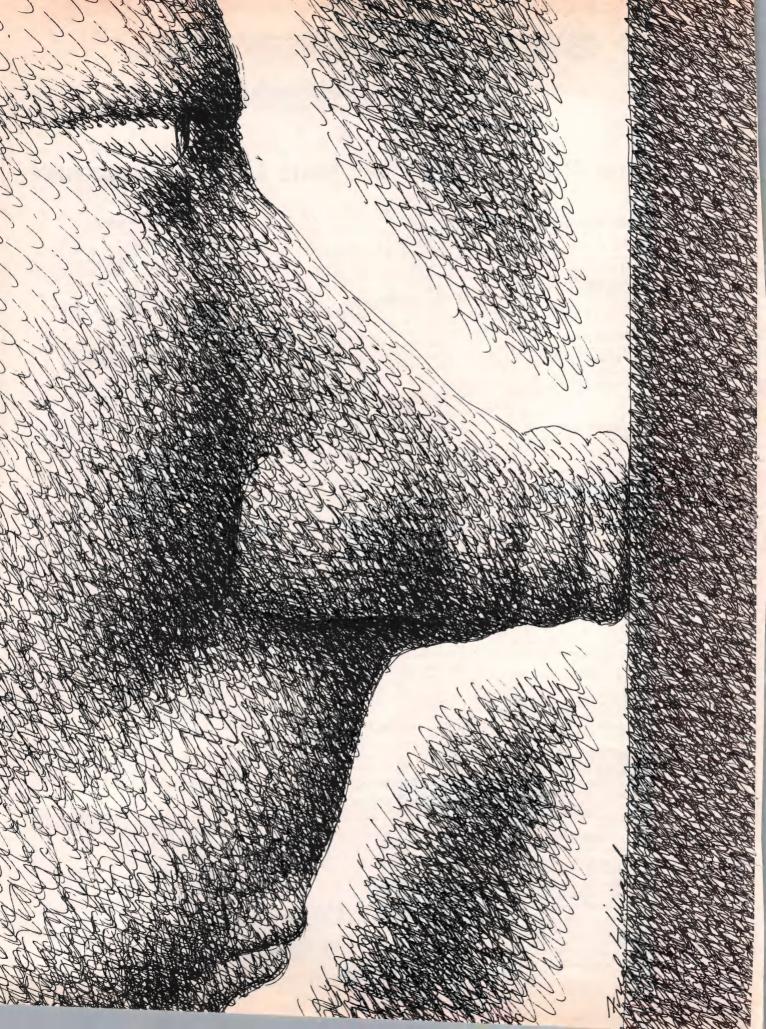
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The XOR Chain

by David E. Cortesi

n a system in which unsigned binary words and storage addresses are compatible types, a peculiarity of the exclusive-OR function allows both links of a doubly linked list item to be stored in a single word. This space-saving gimmick has been part of the folklore of computing for some time. In this article I explore the mathematical basis of the trick and develop a package of C functions to implement it in a modular style.

The Exclusive-OR Function

Consider the exclusive-OR function (XOR). You can think of XOR as an arithmetic function similar to addition and subtraction; it obeys similar metarules. Like them, for instance, it possesses 0 as an identity element. The identity element of a function is the value whose use preserves an identity over the function. For addition:

$$A + 0 = A$$

That is, adding 0 is an identity operation; it leaves A unchanged. In subtraction:

$$A - 0 = A$$
$$A - A = 0$$

Here, not only does subtraction of 0 leave an identity unchanged, but subtracting identical values yields the identity element. Just so:

$$A XOR 0 = A$$

$$A XOR A = 0$$
(1)

Relation (2) is the basis of the common

David E. Cortesi, 415 Cambridge St., #18, Palo Alto, CA 94306. Dave is a former DDJ columnist. You can construct a doubly linked list with both links in the same word.

assembly-language trick of clearing a register by XORing it with itself.

The addition operation is commutative—that is:

$$A + B = B + A$$

This is not true of subtraction, but it is true of XOR:

$$A XOR B = B XOR A$$
 (3)

Addition is also associative—that is:

$$(A + B) + C = A + (B + C)$$

Again, this is not true of subtraction but is true of XOR:

$$A XOR (B XOR C) = (A XOR B) XOR C (4)$$

The XOR operation differs from addition and subtraction in that all the relations (1) to (4) hold for it, whereas only (1) and (2) are true of subtraction and only (1), (3), and (4) hold for addition. Because all four relations hold for XOR, so does this peculiar rule:

$$(A \text{ XOR B XOR C}) \text{ XOR B XOR C} = A$$
 (5)

This may not be immediately obvious. Come at it in steps: apply relations (3) and (4) several times to rearrange the formula so it reads:

By relation (2), this is equivalent to:

A XOR (0) XOR (0) = A

but by (1), 0 XOR 0 is just 0, leaving you with:

 $A \times OR 0 = A$

which merely restates (1).

Three Links in a Word

Mathematically, (5) is a tautology, a gassy expansion of (1), but it has a practical use for data storage. From it follows this: if you have a binary word W that contains the result of:

W = (A XOR B XOR C)

and if you know any two of these values independently, you can recover the third value from W. By relation (5):

 $W \times C = A$

Similarly, you can recover B knowing A and C, or C knowing A and B. You can use this fact to construct a doubly linked list using but a single binary word to store both links in each list item.

How do you do this? In a doubly linked list, each list item contains the addresses of its predecessor and its successor. Call these A and C, and use the item's own address as B. Set its single link word W to:

W = (A XOR B XOR C)

Now, if a program is examining a list item, it obviously must know that item's address. Provided it arrived at the item by tracing the list, it must have come down to it from its predecessor or back from its successor, and hence it must know one of the addresses A or C. Therefore, by (5), it can recover the unknown address from word W and use it to proceed onward in the list.

That's the basic idea. A complete implementation, however, must handle the cases of the empty list and lists of one and two items, should allow for inserting and deleting elements, and ought to be packaged as functional subroutines for simplicity.

Here I'll develop such a package in C. In the following code, I assume that the unsigned binary integer is exactly compatible with a C pointer (as may not be the case in, for instance, the Intel 8086 with a large-memory model), and I've left out the casts that the more picky C compilers may want inserted to make that equivalence manifest. The code favors clarity over absolute efficiency, so you may wish to insert register declarations and other changes for speed.

Defining an Item

In C, the XOR function is denoted by an up arrow and inversion of a Boolean value by comparison to 0. For readability, assume that the following *defines* hold in the example code:

```
#define Xor ↑
#define Not 0==
#define Nil 0
```

Now you can define the structure of a list item:

```
struct Item {
    unsigned link;
    /* other contents of item */
    };
```

The link word contains the XORed linking addresses. What the "other contents" might be depends on the application, and so I assume that functions to create and destroy *Items* are defined elsewhere:

```
extern struct
Item *MakeItem( );
extern void
DropItem(i) struct Item *i;
```

Defining a List

A list consists of a chain of zero or more *Items*, but there has to be an anchor for the chain: a single fixed place where the head and tail of the list can be found. I define its form as *Anchor*:

```
struct Anchor {
    struct Item *head, *tail;
    };
```

For any list there will be just one Anchor. Furthermore, I legislate that if a list is empty:

You can think of XOR as an arithmetic function.

```
(head == Nil) && (tail == Nil)
```

Otherwise, *head* points to the first item of the list and *tail* points to the last item. If the list contains but one item:

```
(head = = tail) && (head != Nil)
```

You can express these rules in Boolean functions that test an *Anchor*:

```
int ZeroItems(a) struct Anchor *a;
{return(
 (Nil = = a-> head)&&
 (Nil = = a-> tail)
 );}
```

```
int OneItem(a) struct Anchor *a;
{return(
(a->head = =a->tail)
&&(a->head!= Nil)
);}
```

Scanning a List

The Anchor of a list gives you access only to the first and last items. If you are to reach intermediate items, you must scan over the list, either forward from the head or backward from the tail.

While scanning the list, your position is always maintained in a pair of pointers that contain the addresses of two items that are logically adjacent in the list. The item that is nearer the head of the list I call the prior item; the one closer to the tail I call the next item. When I have these items set up with valid addresses, I refer to such a pair of pointers as a Scan, because I can use the pair to scan over the list in either direction. Because a pair of pointers is always required, I put them in a record. Because I am always scanning a particular list, I include the address of the list's Anchor.

```
struct Scan {
    struct Item *prior, *next;
    struct Anchor *base;
    };
```

Although a list may have only one *Anchor*, it may have any number of *Scans* associated with it. No *Scan* is valid until it has been associated with some list, however. Here's a procedure to do that:

```
void Associate(s,a)
struct Scan *s;
struct Anchor *a;
{ s->base = a; }
```

This procedure is very simple, but in some applications, it might have to do more. It might be desirable, for example, to keep track of the number of scans that are active on a list, and the *Associate* function could do that by incrementing a count in *Anchor*.

A program may scan a list by moving a *Scan* structure from the current position to an adjacent one. But a scan has to start somewhere. This procedure sets a scan to the head of the list:

```
void ToHead(s) struct Scan *s;
{
    s->next = s->a->head;
    s->prior = Nil;
}
```

When a scan is at the head, the *prior* pointer contains *Nil* (nothing precedes the head of a list) and the *next* item is the head item of the list. A simple function can test a *Scan* for this condition:

```
int AtHead(s) struct Scan *s;
{ return( s-> prior == Nil ); }
```

XOR CHAIN

(continued from page 29)

A procedure can set a scan to the tail of its list:

```
void ToTail(s) struct Scan *s;
{
    s->prior = s->a->tail;
    s->next = Nil;
}
```

When a scan is at the tail of a list, the *next* pointer contains *Nil* (nothing follows the tail of a list) and the *prior* item is the tail item of the list. A Bool-

ean function can test for this condition:

```
int AtTail(s) struct Scan *s;
{ return( s->next == Nil ); }
```

Recall that a list is empty when the head and tail pointers in its *Anchor* are *Nil*. Verify for yourself that, for a *Scan* associated with an empty list, applying either *ToHead* or *ToTail* makes both *AtHead* and *AtTail* report true. Earlier I defined a function *Zeroltems*, which tests an *Anchor* to see if it's empty. Now I can define a function that tests not an *Anchor* but a

Scan:

int EmptyList(s) struct Scan *s; { return(AtHead(s) && AtTail(s)); }

Consider the implications for loop control. Either the sequence:

ToHead(z); while (Not AtTail(z)) {step forward in list}

or the sequence:

ToTail(z) while (Not AtHead(z)) {step backward in list}

can work safely—that is, do nothing—on an empty list.

Stepping a Scan

Having set a *Scan* to the head of a list, you want to be able to step it forward in the list. You do that using the identities of XOR. Consider an *Item* that is neither at the head nor the tail of a list. If it is located at address B and if its predecessor is at address A and its successor at address C, then its link word contains A XOR B XOR C. The link word plus the contents of the *Scan* let you step it forward:

void GoFwd(s) struct Scan *s;
{ struct Item *i;
 i = s->prior Xor
 s->next Xor
 s->next->link;
 s->prior = s->next;
 s->next = i;
}

Variable *i* receives the link word of the next *ltem* minus the contributions made by its own and its predecessor's address: in short, its successor's address.

In this way you can process all the items of a list:

Associate(z,a);
ToHead(z);
while(Not AtTail(z))
{
 Process(z->next);
 GoFwd(z);
}

If you are accustomed to C, you may prefer to see this written using a for

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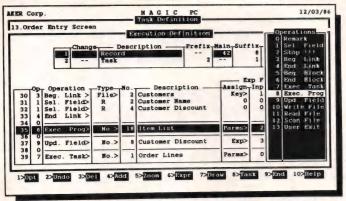
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Order No: 999 Order Date: 99/99/99 Item Type Description Quantity Unit Price Total Price 9,994 499 499 999,999 99 Description Sales Tax 99.994 Order Total 1>Opt 2>Undo 3>Del 4>Add 5>Zoom 6>Expr 7>Draw 1>Task 9>End 100 Help

program in the Program Menu, you tell Magic PC to Zoom into the Item List program through the window shown on the right screen. The window will automatically scroll the Item List data horizontally and vertically, and allow query, "cut and paste" copy or even creation of new Items.

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XOR CHAIN

(continued from page 30)

loop:

Associate(z,a); for(ToHead(z):

Not AtTail(z); GoFwd(z))

Process(z->next);

which amounts to the same thing.

When you process a list from head to tail, the *Items* to be processed are those that turn up in the next posi-

tion. When you go backward, from tail to head, you take them from the prior slot. Here's a going-backward procedure:

```
void GoBak(s) struct Scan *s;
{ struct Item *i;
    i =s->prior Xor
        s->next Xor
        s->prior->link;
s->next = s->prior;
s->prior = i;
}
```

Here you take the link word of the

predecessor *Item* minus the contributions of its own address and of its successor: in short, the address of its predecessor. With this procedure and *ToTail*, you can process a whole list:

```
Associate(z,a);
ToTail(z);
while( Not AtHead(z) )
{
    Process(z->prior);
    GoBak(z);
}
```

It may not be clear that these procedures work for all items of all lists. Before you can be sure they do, you need to define what the link words of the head and tail items contain. The simplest rule works: the link word contains xOR o for irrelevant addresses. Table 1, below, shows a complete four-item list whose anchor is *Test*. Presume that the following loop is to be executed using the list in this table:

Associate(z,Test);
ToHead(z);
while(Not AtTail(z)) GoFwd(z);

Trace the contents of *prior* and *next* at each step and assure yourself that all the pointers work out correctly. (Surely you don't take articles such as this one on faith, without at least desk-checking the code?) Then trace them through this loop to verify the rest of the code given so far:

for(ToTail(z,Test); Not AtHead(z); GoBak(z));

What would happen if a program applied *GoFwd* just once too often? Or *GoBak*? Clearly you ought to account for these end effects. What shall be done with the head item (which has no predecessor) and the tail item (no successor)? You can either wrap around or stick. Here is a safe *StepFwd* procedure:

Address	Contents	
Test.head	A	
A->link	A xor B	
B->link	A xor B xor C	
C->link	B xor C xor D	
D->link	C xor D	
Test.tail	D	

Table 1: A four-item list whose anchor is Test

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void StepFwd(s) struct Scan *s;
{ if (Not AtTail(s)) GoFwd(s); }

that simply sticks when it reaches the tail. Here is a *StepBak* procedure that wraps around:

```
void StepBak(s) struct Scan *s;
{
    if (AtHead(s)) ToTail(s);
    else GoBak(s);
}
```

Inserting an Item

You construct a list by starting with *Zeroltems* true and inserting new items. Example 1, below, shows one way to insert a new item between the prior and next items of a scan. In the statement:

s->prior->link Xor=
(i Xor s->next)

```
void Insert(i,s)
struct Item *i;
struct Scan *s;

{
  if (AtHead(s))
    s->a->head = i;
  else
    s->prior->link Xor=
        (i Xor s->next);

  if (AtTail(s))
    s->a->tail = i;
  else
    s->next->link Xor=
        (i Xor s->prior);

  i->link =
        s->prior Xor s->next Xor i;
  s->prior = i;
}
```

Example 1: A way to insert a new item between the prior and next items of a scan.

Address	Contents
Test.head	Nil
Test.tail	Nil
A.link	?
B.link	?
C.link	?
D.link	?

Table 2 An empty list and four prepared items

you take the predecessor's link, remove from it the contribution of the successor's address (by XORing with *s*->*next*); and install the value of its new successor—the item with address *i*.

A similar statement fixes the link word in the next item, except that it's the predecessor's address that is removed in favor of *i*. Finally, the new item's link is formed from its own address and that of its new predecessor and successor. It becomes the prior item, so a series of insertions will build up in head-to-tail

order.

To verify this, commence with an empty list and four prepared items, as shown in Table 2, below, and execute the following sequence of operations by hand:

Associate(z,Test); Insert(A,z); Insert(D,z); GoBak(s); Insert(B,z); Insert(C,z);

keeping track of the new contents of

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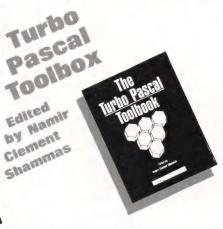


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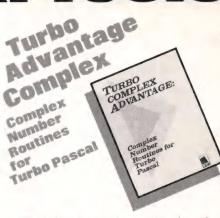


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(continued from page 35)

Anchor, the Item links, and prior and next.

Deleting an Item

You need a way to delete an item of a list. Deleting comes down to fixing the things that point to an item so they point to other items instead. Once nothing points to it anymore, an item can be discarded. The deletion procedure in Example 2, page 36, deletes the next *Item* of the scan. If there is no next *Item*, it does nothing (ergo, it correctly does nothing to an empty list).

The item to be deleted is *s->next*. The first step is to get the address of its successor, which will take its place. It is obtained from its link field in the manner of *GoFwd*.

Next you repair the predecessor of the deleted item. If it's the head of the list, the anchor must be repaired; otherwise, the predecessor's link field is adjusted to point to the deleted item's successor.

The successor of the deleted item is then repaired. If there isn't one, the deleted item was the tail of the list, and its predecessor becomes the new tail.

You could empty a list with this loop:

Associate(z,Test);
ToHead(z);
while (Not AtTail(z)) Delete(z);

Although this procedure won't delete when the scan is *AtTail*, the tail item can nevertheless be deleted using:

ToTail(z); StepBak(z); Delete(z);

Earlier I said that any number of *Scan* structures could be associated with a list. Now that I've defined *Insert* and *Delete*, I should perhaps modify that claim. What might the effect be on one *Scan* if *Insert* were performed via a different *Scan* of the same list? (Answer: not much—the new item might be missed by the other *Scan*.)

What about Delete? (Answer: catastrophe-the other Scan might step using the link word of a deleted, now-invalid Item.)

Conclusion

The C functions shown here form the basis of a package for handling doubly linked lists. These lists have several advantages: because they are doubly linked, they can be traversed in either direction with equal ease, and insertion and deletion are simple and speedy operations. By packaging a list position as a single *Scan* structure, you make it convenient to track several list positions at once, as long as no deletions are allowed.

There are disadvantages, of course. The greatest is that the only practical access to the list is from its ends; it isn't possible to enter it at the middle and then traverse forward or backward. (Traversal requires the addresses of two logically adjacent items, which can't be maintained outside the list itself in the presence of *Insert* or *Delete*.) Thus this list form can't be indexed for quick entry at intermediate points.

Readers of an imaginative turn of mind might explore some other possibilities. For instance, the structure treated here as a linear list could be wrapped around to make an endless ring. In a double-linked ring, *Anchor* might be eliminated in favor of a mandatory single *Item*.

DDJ

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void Delete(s) struct Scan *s; } struct Item *i; if (AtTail(s)) return: $i = s-\rangle next-\rangle link$ Xor s-)prior Xor s-)next: if (AtHead(s)) $s-\rangle a-\rangle head = i$: PISP $s-\rangle prior-\rangle link Xor = (s-\rangle next$ Xori); if (i = = Nil) s->a->tail = s->prior; PISP i->link Xor = (s->next Xor s->prior); DropItem(s-)next); s-)next = i;

Example 2 A procedure to delete an item in a list

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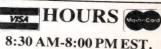
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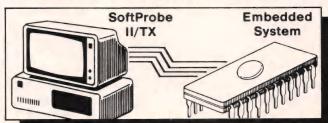
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An Extended IBM PC COM Port Driver

by Thomas A. Zimniewicz

he IBM PC and its compatibles support two serial ports-COM1 and COM2—in a wide variety of configurations. Four basic functions are provided—data in, data out, set configuration, and read status. What the IBM PC does not do is walk and chew gum at the same time. If your disk is seeking and more than one character arrives, it is lost. This article describes what I think a COM port driver should do, discusses some design rules and implementation pitfalls of device drivers, and presents a device driver modification with the following features:

- buffered input—no lost data
- input flow control—XON/XOFF or control lines
- output flow control—control lines
- · works with existing software
- · removal—without a reboot
- 19,200- and 38,400-baud operation
- type-ahead with CRTs
- · upward compatible

What's Wrong?

During the first part of my career as a software engineer, I was spoiled by minicomputer operating systems that were good at capturing serial data. More recently, I had the pleasure of working with Unix—a pleasure, that is, in all aspects except capturing serial data. My introduction to Unix was on a PDP-11/44 that could miss characters from serial data input at 300 baud on a busy day. Every time my job required monitoring serial data, it was a bad day. Then one

Thomas A. Zimniewicz, 2695 Pond Rd., Lima, NY 14485. Thomas is a software consultant. What the IBM PC does not do is walk and chew gum at the same time.

day I noticed an idle IBM PC that had nothing to do (no other users to get in my way) except perhaps collect my data. It would be simple—just getchar from COM1 and putchar to disk, and with no other work, the IBM PC could certainly handle at least 9,600 baud. A while later it was working, but (and there always is one) it seemed that there were gaps in the data. (I discovered later that the gaps coincided with the flashing of the disk light.) I gave up. Much later I found myself using an IBM PC a little. I didn't like the keyboard and monitor, so I hooked up my CRT to a COM port and then found there was no type-ahead. I got mad enough to fix it.

So, what were my goals? The main thing was buffered interrupt-driven input to allow type-ahead and highspeed serial data capture. I wanted to use these features with existing programs, too. This meant the solution had to be fully compatible and somehow get between programs and the hardware. I had to modify PC-DOS and/or the BIOS in a way that was transparent to an application-level program. In this sea of bad news, there was one very bright light-everything that PC-DOS and the BIOS lacked in capability, it made up for in flexibility. The hooks were there, device drivers could be installed, interrupts could be stolen. If the new COM port software were installed on an IBM PC, then all well-behaved programs (those that used PC-DOS or the BIOS instead of directly accessing the hardware) would automatically gain the new features. Several methods of input/output flow control were possible and the hardware could handle higher baud rates . . . I sure do get carried away easily.

Can It Be Fixed?

The IBM PC actually has two levels of software to handle the COM ports. The high-level code is part of PC-DOS and is loaded into RAM during the boot-up of PC-DOS. This code handles those functions that do not depend on the details of the hardware and provides an interface to the COM port. This level is called a device driver and is accessed by high-level PC-DOS operations such as those that open and read serial data system calls. The low-level code is hardware specific and is contained in ROM. This code is specific to the operation of the COM port hardware. This level is called the BIOS and is called by the device driver to perform operations specific to a COM port, such as setting the baud rate or getting a character from the port hardware.

When I started to investigate the job at hand, I suspected that a whole new device driver would be required. The good news was that PCDOS faithfully used the BIOS for all COM port operations. This meant that all I needed to do was to write a replacement for the parts of the BIOS that handle the COM port. For a while the job looked really easy. I considered stealing the software interrupt

(int 14h) that is used to access the BIOS COM port code, going to my code for operations that needed changing, and otherwise just passing the job off to the BIOS. Unfortunately, I had to abandon this easy way out because of negative side effects caused by almost every BIOS function. For example, the BIOS "send a character" function sets the control signals in ways that are incompatible with some of the required flow control modes.

To illustrate some of the technical issues of handling a COM port, let's take a simplified look at the job of reading serial data. When the BIOS call is made to read a character, it loops until a character arrives at the port or a counter expires. Either the character or a time-out status is returned. Most programs and PC-DOS ignore the time-out and try again. The COM port hardware has only a single character buffer. If the buffer in the COM port hardware overflows while the processor is doing anything else, it is lost. The solution to these lost characters is to have the COM port generate an interrupt when a character arrives and have the interrupt handler save the character in a buffer until the program gets around to needing it. The code might look like this:

COM port interrupt handler read character if buffer count < buffer size put character in the buffer increment buffer count else set overrun error

read character routine
loop
if buffer count > 0
decrement buffer count
take character and return it
if timer expired
return time-out status
endloop

If this algorithm were implemented, it would probably pass all its initial tests. Then, when you least expected it, your input would be garbled and no errors would be reported. What you have here is your classic race condition. Imagine the buffer is exactly full when the read routine is called and the count is decremented to indicate room in the

buffer. Immediately after the count is decremented but before the character is removed, an interrupt occurs. The handler puts another character in the buffer that is really still full. Things go downhill from here. This type of error is an easy trap for beginners (and sometimes a big embarrassment for pros). One possible fix is to disable interrupts during critical portions of the read character routine. (Note that COM port interrupts are automatically disabled by the hardware during the processing of a COM port interrupt; there is no need to worry about reentering the handler itself). This works, but it has

I wanted
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serious drawbacks. If interrupts are off for too long, serial data (and other real-time events) can be lost. The best solution is to order the events in the read character routine so that an interrupt cannot hurt it. If interrupts must be disabled, try to limit the duration. To fix the bug above, just take the character from the buffer before decrementing the buffer count.

Code Walk-Through

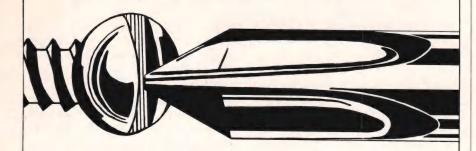
Excom is functionally equivalent to the BIOS int 14h COM port handler with three significant enhancements-it provides interrupt-driven buffered input and has an extended set of configuration parameters, including flow-control selection and higher baud rates. Excom is made up of three major sections-the COM port data input interrupt handler, the replacement for the BIOS int 14h handler, and the code to install and initialize excom. A problem exists in trying to describe excom, unlike a well-written application in which functions can be isolated to a single routine or group of routines, because its structure dictates that many of its capabilities are distributed through-

out the code. I will first take you on a quick trip through the listing (Listing One, page 64) to show its structure and then later describe several of excom's more interesting functions.

Excom is a .COM program that is run once before the COM ports are needed and that uses the terminateand-stay-resident feature of PC-DOS to remain active as long as is needed. If required it can be removed, restoring the interrupt vectors to their previous values and freeing memory. Two requirements for a .COM program are to have all segment registers pointing to the start of the program's memory space and to have an executable origin of 100h to make room for the program header. Execution starts at location 100h, which is a jump to the initialization code. A few constants are then defined, the first three of which have to do with the buffer size and can be modified to tune the excom buffers to your needs. The Microsoft assembler supports a concept similar to structures in C. A structure called pcb (port control block) is defined to describe the dynamic state of a port. It is defined as a structure so that one piece of code can handle two ports just by setting a register to point to the appropriate structure. Storage for two port control blocks is then allocated. Space is reserved to store the old interrupt vectors required for the program to remove itself. Next the baud rate table contains "magic" numbers required to set the UART's baud rate clocks. Note that the last two entries are used for the excom extended baud rates.

The first executable code is the COM port hardware interrupt handler. The label intOB is the entry point for int 0Bh, COM2, whereas int0C is for int 0Ch, COM1. Each entry sets ds:si to point to the pcb associated with its port and enters common code. The rest of the routine reads the character, puts it in the buffer, takes any action required for flow control, and exits. There is no analogy to this code in the BIOS. The second code segment is the int 14h handler, which performs five distinct functions corresponding to the value in the register ah. The functions are initialize port, transmit a character, receive a character, get port status, and extend initialization. At the

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COM PORT DRIVER (continued from page 43)

end of excom is the initialization. This is the code that runs only when excom is being installed. Initialization starts by releasing the memory for excom's copy of the environment, which is not used. The body of the initialization uses PC-DOS calls to get the old interrupt vectors and then install the three new interrupt handlers. The interrupts that are thus stolen are OBh (COM1 hardware). 0Ch (COM2 hardware), and 14h (BIOS COM port calls). Then the excom COM port data structures are initialized. This initialization code is at the end so that it is not kept as part of the resident code. Initialization ends with the PC-DOS call to make everything before the initialization code stay resident.

Key Functions

The BIOS does COM port input by waiting for a character to arrive. Even though it is not used, the IBM PC has all the hardware required to generate an interrupt when a character arrives. This capability is activated by changing the interrupt vector (a RAM location) to point to the appropriate handler (see lines 555-592 of Listing One). Now the vectors are set up, but because the BIOS doesn't use COM port interrupts, the hardware has not been initialized to generate interrupts. Two separate pieces of hardware have to be initializedfirst the COM port hardware (an 8250 UART) has to be programmed to generate interrupts and then the interrupt controller (an 8259) has to be set to allow the interrupt to be sent on to the CPU. I expected this to be a simple flip through the data sheets and a dozen lines of code. It didn't workno interrupts were generated. After much pain, I discovered that a spare output called OUT2 on the UART was used to gate the interrupt signal from the UART. I have no idea why. I then asserted OUT2 on the UART and voilà-interrupts (see lines 290-295). When an interrupt does occur, the CPU stops what it's doing and jumps through the appropriate vector-in this case to int0B or int0C (lines 73-158). The interrupt handler then reads the character, which resets the interrupt status in the UART and

sends an interrupt complete command to the interrupt controller before returning.

The BIOS port status request simply reads the two UART status registers and returns them to the caller. One register is called the line status and contains bits for receive errors, receive data ready, time-out, and transmit ready. The other is called modem status and has the UART's control lines. When characters are transmitted or received, the ah register returns the line status. When a port is initialized or status is requested. ah is returned as above and al returns the modem status. In excom input characters are stored in a circular buffer as they are read from the UART and are then removed from the buffer when a character read takes place. This presents a dilemma for excom because some status bits are associated with the received character whereas others are really HART status. The solution is to have the interrupt handler read the status associated with receive characters and buffer both character and status together. Then when a status is needed and there are characters in the buffer, the status associated with input characters is taken from the next character in the buffer. The rest of the status is taken directly from the UART (see lines 429-457). This approach is not always perfect, but excom has been used successfully with many port handling packages, including COMMAND.COM.

One of the major goals for excom was to provide configurable flow control for input data. The objective of flow control is to signal the source of the data to stop sending when the buffer is getting full (see lines 96-120). Then, when the buffer is emptied a bit, the sender is told to resume sending (see lines 394-409). The exact number of characters in the buffer when input is stopped and started is configurable (lines 16-18). Having the stop and restart values be different by a few characters tends to minimize the number of times the data is stopped and restarted. The actual mechanisms used to tell the sender to stop are any combination of DTR (data transfer register), RTS (request to send), and XON/XOFF (Control-S/ Control-Q). DTR and RTS are turned off as the buffer fills and turned back

on when room becomes available. Control-S is transmitted as the buffer fills, and Control-Q is transmitted when room becomes available. I always thought that RTS was used as described by its name, but I've seen others use it for input flow control. It was easy to add and it didn't seem to do any harm, so I used it for input flow control, too. Add to this arsenal the ability to cross a few wires in a cable, and almost any situation is covered-except, of course, the sender that ignores your attempts. This is where excom shines—interrupts and big buffers provide great power. Data can be copied from a COM port to a floppy on a lowly old 4.77-MHz IBM PC at 9,600 baud without missing a character. To complete the picture, I also added output flow control. The BIOS requires that both DSR and CTS be set before a character can be sent. Excom can be configured to ignore either or both of these signals.

In order to be useful, excom had to be upward compatible with the BIOS *int 14h*. The area of extended initialization options presented some special problems. There are no unused

bits in the BIOS set configuration command. Therefore I had to add a new command-int 14h with ah having the values from 0 to 3 were already used, so ah = 4 was the obvious choice (lines 460-499). The problem was that the old initialization—ah =0-included the baud rate and the new initialization could also specify baud rates not included in the BIOS. The obvious solution was to require that the BIOS initialization occur before the extended initialization. If excom was to be used with existing software, though, this restriction was impractical. The actual solution was to have excom ignore the baud rate in the BIOS initialization if the extended initialization had been used to set a nonstandard baud rate (lines 304-306). Thus, you set the baud rate to 19,200, start your existing software, and a BIOS initialization call is made, but excom ignores the baud rate. Another concession to upward compatibility is that the extended options are relative to the BIOS defaults. Excom allows DTR input flow control to be enabled, whereas output CTS flow control can be disabled.



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COM PORT DRIVER (continued from page 45)

Not All Is Gold

Along the way from the problem to the solution, there were many little struggles, one of which may be of interest. The rest are too embarrassing to put into print.

While first using excom, I ran into a strange problem. I have a homemade shell that reads the keyboard (CRT hooked to COM port) by calling int 14h directly. I ran a program that did CRT output via PC-DOS but did not read any keyboard characters. While it was running, I typed a command ahead. When the program stopped, my type-ahead was echoed minus the first character. This was early in my use of excom, so I assumed a bug in excom. Ignoring the problem, I continued what I was doing. Some time later I ran an interactive program that used PC-DOS for keyboard inputs, and my lost character from 20 minutes ago appeared on the screen as the first character echoed. Wow! At first I passed it off as a bug in excom's input buffer and a coincidence. It happened again! I considered a career as a crossing guard.

You may have noticed that I described the exact way in which the above programs read their inputthat was the key to understanding. PC-DOS makes a feeble attempt to do some input buffering, and when PC-DOS calls int 14h to transmit a character to the port, a status is returned. If a data ready is indicated, PC-DOS does a data read and saves the character for later. PC-DOS only stores one character in this way as far as I can tell. So my type-ahead was eaten by PC-DOS, to be provided to the next call to PC-DOS for a character. Beware of mixing PC-DOS and BIOS character reads!

Anything Left to Do?

Of course there is! The two questions I've heard most often as a software engineer are "When will the code be complete?" and "How many more bugs are there?" The answers have always been the same: "About a week or two" and "Three." The answers are the same here.

If you put a file on a floppy and enter the command *type file*, you will notice a flip-flop of data scrolling by and disk activity. This is because

the BIOS waits for both the disk and the display. Excom has input buffering to prevent loss of data. Output buffering won't fix any problems, but it will speed up data output in the case when the source of the data is a slow device. If excom had output buffering, it would take the characters faster than they were actually being transmitted. Thus PC-DOS would think a block of data had been sent and go to the source for more while excom was still sending data out of its buffer. The flow would be limited by making PC-DOS wait when excom's output buffers were full.

Earlier I mentioned that excom

An
interesting
possibility
with excom
is nesting
several
invocations.

can work with any well-behaved program. What about the naughty ones? At work, I use an old version of CrossTalk that steals the COM port interrupt and doesn't put it back. I don't know if this has been fixed. Are you listening CrossTalk folks? Upon return from CrossTalk, excom has lost the COM port interrupt. One possible solution would be to run Cross-Talk from a batch file that removed excom, ran CrossTalk, and then reinstalled excom. The problem with this is that you must identify and fix all offenders. It may be possible for excom to examine its environment for damage and perform repairs. I'm sure this would not cover all such problems, but I think all of them that I am currently aware of could be fixed. Self-repairing programsmaybe I should do an article on "Core Wars."

When I use my shell on a COM port with or without excom installed, XON/XOFF output flow control seems to work. There must be problems. I have heard of requests to provide XON/XOFF in relation to printer driv-

ers and so on. In any case, because PC-DOS may not see characters until some time after excom does, any PC-DOS attempt at XON/XOFF output flow control is hampered by excom. This implies that XON/XOFF output flow control should be added to excom.

Because the IBM PC does not have memory management hardware, having a memory-resident program occupy the wrong location can cause problems. When a shell runs a program, it is loaded just after the shell. If the program remains resident upon its termination, the shell is in a squeeze. The shell cannot get more memory to hold shell variables, environment, temporary buffers, and so on. A desirable extension to excom would be to have it relocate itself to the end of RAM and remain resident there.

Excom has a bug! Software is no fun if its perfect. Control-S and Control-Q are sent without any regard to the UART being ready or the control lines being correct. This sounds pretty bad, but in reality it's only a problem when there is a heavy flow of data in both directions, which is pretty rare. The reason I allow this bug to exist is that the proper solution requires interrupt-driven output. *DDJ* has only so many pages for my giant listings.

The version of excom on the listings disk is an updated version with interrupt-driven output, corrected sending of Control-S/Control-Q, the ability to repair the damage done by ill-behaved programs, the ability to do output XON/XOFF flow control, and self-relocation to the end of RAM. (See the "Availability" section at the end of this article.)

Use Excom on Your IBM PC

To make excom.com, use the followings commands:

masm excom.asm; link excom.obj; exe2bin excom.exe excom.com del excom.exe

Normally, you should install excome arly in the boot process—before a print /d—because the print spooler steals int 14h. The real-time clock initialization program on my clone messes up the int 14h vector for an unknown reason; therefore I in-

"How to protect your software by letting people copy it."

By Dick Erett, President of Software Security



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protection of intellectual property.

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COM PORT DRIVER (continued from page 46)

stall excom after running the clock initialization. A bit of experimentation may be required. Excom is selfinstalling; just run excom to install it. If all you need is buffered input, you've finished. Excom will be properly set up by programs using the COM ports, and the existing mode command will configure excom. To tap the power of excom, you may need to access some of the extended features. If you use excom with your own programs, you can just make an int 14h call with ah = 4 to set any extended options you may need. If you use excom with existing programs, you need a way to set the extended options from your keyboard or a batch file.

Listing Two, page 75, is a simple little C program called exmode that sets extended configuration parameters. Exmode works with the Microsoft 4.0 C compiler. If you have a different one, you may need to rewrite the function *int14*—it just sets registers and performs a software *int 14h*. You

can also use exmode to install and remove excom, usually useful only while testing.

Exmode is self-documenting; just run exmode and it will tell you what it can do. Note that running exmode does not add to existing extended options but sets the complete set—that is, exmode com1 nodsr followed by exmode com1 nocts is not the same as exmode com1 nocts is left set; in the second, both nocts and nodsr are set. COM1's settings can be cleared by running exmode com1. The previous examples will not change any of COM2's settings.

An interesting possibility with excom is nesting several invocations of excom. Say you have excom set up the way you want it for a shell running on your CRT and want to temporarily run a communications program. Just install excom again, set it up for communications, run the communication program, and then use exmode to remove excom. The most recent copy of excom is removed, restoring your old excom untouched.

Availability

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(Listings begin on page 64.)

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Dynamic Memory Overlays for Turbo Pascal

by Steve McMahon

bo Pascal has a significant limitation that hinders its usefulness for developing large programs. Because the compiler can generate code for only a single code segment, it imposes a 64K limit on executable code space. Turbo Pascal's overlay system, designed to circumvent this restriction, potentially allows for much larger amounts of code but at the cost of dramatically slowing program execution as overlay code is loaded from disk files.

This article describes a way to use Turbo Pascal's overlay facility in conjunction with dynamic memory to build memory overlays that can be loaded and executed far more rapidly than can disk overlays. These memory overlays are easy to set up and can optionally be left on disk at run time if sufficient memory is lacking. This technique lets you have the best of both worlds: a program that can execute in limited memory with disk overlays but that can take advantage of extra memory to load overlays into RAM.

Turbo Pascal's Overlay Scheme

Borland's solution to the large-program problem is an overlay system. Programmers confronted with a program that would otherwise overrun the 64K code segment limitation can specify that the code of a group of procedures and functions should

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A fast and versatile way to handle large programs in Turbo Pascal

share the same execution space in the code segment. You make this declaration by grouping the routines contiguously in the source code and preceding the declaration of each overlay routine with the Turbo reserved word overlay. On compilation of the program, all the routines of an overlay group are compiled to a separate overlay file, where the executing program can find them as necessary. You can specify as many overlay groups as you wish within a program. The only coding limitation is that routines within an overlay group cannot call themselves (they cannot be directly recursive) or other routines within the same overlay group. (The recursion limitation can be circumvented by calling the overlay routine from a recursive, nonoverlay function or procedure.)

The advantages of this scheme are that it is easy to use and that it allows you to construct programs with code not only larger than 64K but also potentially larger than available RAM. The disadvantage is a significant performance degradation: disk activity and DOS file handling overhead are now required every time an overlay routine is executed. Careful selection and grouping of overlay routines can minimize such performance degradation, but it is still possible for a program to spend far more time reading

overlays from disk than doing anything else.

Overlay performance can be significantly enhanced by storing the overlay file on a virtual memory, or RAM, disk. Turbo Pascal's OvrPath procedure makes this possible by allowing run-time specification of the location of overlay files. This solution to the performance problem. though, has requirements-that a RAM disk be available, overlay files be located there, and the program be correctly informed of that locationthat are unacceptable for application programs intended for use in normal DOS environments by unsophisticated operators.

A solution to the overlay performance problem that is faster, more versatile, and completely invisible to an application program's users is presented here. But first, let's take a look at how Turbo Pascal's overlay system works.

How Overlays Work

When Turbo Pascal compiles a program that has overlay procedures, it reserves an overlay area in the code segment (and in the .COM file) large enough to accommodate the largest of the procedures or functions in a given overlay group. All the procedures and functions in that overlay set are then compiled for execution at the same position within the reserved overlay area. Those procedures and functions are stored all together in a separate overlay file that bears the same name as the main program but has the number of the overlay group (.000, .001, and so on) as a file-name extension.

When some other portion of the program makes a call to an over-

layed function or procedure, the calling code passes in registers the offset of the required overlay function or procedure inside the overlay file and the length of the code fragment to be loaded. The call is made to a fragment of code positioned at the top of the reserved overlay area. This fragment passes control to the overlay handler in Turbo Pascal's run-time package, having placed the address of the reserved overlay area on the stack. The overlay handler responds by checking to see if the code required is already in place in the overlay area (the overlay file offset of whatever code is currently in the overlay area is stored in a data area at the top of the reserved overlay space along with the name of the overlay file). If the code required is already in place, the overlay handler executes it. If not, the code is read from the overlay file: the overlay file is opened, a seek is made to the reguired offset, and the specified quantity of code is read into the overlay area. Finally, the file is closed, the overlay file offset of the now-current code fragment is saved in memory for future reference, and the overlay function or procedure is executed.

This overlay handling method is ripe for alteration. If the contents of the overlay file could be positioned in some otherwise unused portion of memory, then a replacement overlay handler could snatch code fragments from that area of memory rather than from a disk file. The time savings could be impressive: overlay calls would require no disk activity. Such a scheme should even provide significant speed advantages over using Borland's overlay handler in conjunction with a RAM disk (to contain the overlay file) because DOS file handling overhead would be eliminated. What a memory overlay scheme couldn't do, though, is allow for programs that are truly larger than memory because all the overlay code would need to fit in memory available to the program.

Such a scheme can be implemented with no changes to the compiler and with surprisingly little extra code. Best of all, it can be done with little modification of existing programs that use overlays. All that's necessary is to include the procedures contained in MemOvrly.Inc

(Listing One, page 78) in a program that uses overlays and to add some initialization and deinitialization code to the program. (Some limitations on the scheme are discussed later.)

Using MemOvrly.Inc

You can include the three procedures that make up the memory overlay handler in an existing program by adding the line:

{\$I MEMOVRLY.INC}

at a point in the procedure and function declaration part of a program that uses overlays (do not nest it inside a procedure or function). Then, you must add a procedure call, probably inside the main body of the pro-

gram, to initialize the substitute overlay handler for each overlay group that you wish to use as a memory overlay group. You can add code to dispose of the dynamic memory used by the memory overlay group or groups at the end of the program. This cleanup code is particularly important if the program chains to another Turbo Pascal program because Turbo Pascal preserves the dynamic memory heap on chaining.

To set up a particular overlay group as a memory overlay group, all that's necessary is to run the *Init-Overlay* procedure, passing it the address of some procedure or function in the overlay group. If, for example, an overlay group contained, among other routines, the function or procedure *One*, the entire overlay group

```
PROGRAM OverlayTest;
    (* Memory Overlay Demonstration Program.
    {$I MEMOVRLY.INC}
  VAR
   c : Char;
  OVERLAY PROCEDURE One;
   BEGIN
      WriteLn ('This is Overlay Procedure One.');
   END;
  OVERLAY PROCEDURE Two;
    BEGIN
      WriteLn('This is Overlay Procedure Two.');
    END:
  BEGIN
    {Install the new overlay handler by passing it the address
    offset of ONE procedure or function from the overlay group.
    Multiple invocations for multiple overlay groups should be
    no problem. }
    InitOverlay(Ofs(One));
      Write('Hit any key to run the overlays (^Z to stop): ');
      Read (Kbd, c);
      WriteLn;
      IF c <> ^Z THEN
        BEGIN
          One;
           Two:
        END:
      WriteLn;
    UNTIL c = ^Z;
     (Free up the heap space used by the replacement overlay
    handler by passing the same offset as above to the
    DisposeOverlayStorage Routine -- VITAL if you're chaining
    to another program. The heap is preserved in a chain operation.}
    DisposeOverlayStorage (Ofs (One));
```

Example 1: Short program demonstrating memory overlays

TURBO PASCAL OVERLAYS (continued from page 51)

could be initialized as a memory overlay group with the statement:

InitOverlay(Ofs(One));

Ofs is a built-in Turbo Pascal function that returns the offset within a segment of a variable or procedure. Because vou know that the reserved overlay area associated with the routine One is in the code segment, passing the offset of the routine is sufficient to establish where the reserved overlay area for this group is located. This instruction allocates dynamic memory sufficient to contain all the overlay code, reads that code into memory, and installs the procedure NewOverlayHandler as overlay handler for that group. Any overlay groups not set up in this fashion are handled by Turbo Pascal's normal overlay handler. Also, if not enough dynamic storage is available to accept the overlay code, the group is left as a normal, nonmemory overlay.

Likewise, the memory used to contain all the overlay code for a given overlay group can be reclaimed for other uses with the instruction:

DisposeOverlayStorage(Ofs(One));

It's crucial that you do not try to call any routine in the overlay group after issuing this instruction. The DisposeOverlayStorage procedure does not restore the old, nonmemory overlay handler for the specified overlay group. It merely frees up the memory previously occupied by overlay code for other uses. So, the DisposeOverlayStorage command is usually used only (if at all) at the end of a program. If the program does not use Turbo Pascal's Chain facility to chain to another Turbo Pascal program, you probably don't even need to use the DisposeOverlavStorage procedure. If you do need it, it should be executed once for each memory overlay group in use.

How It Works

When asked to initialize a memory overlay group, the *InitOverlay* procedure pulls the name of the overlay code file from the reserved overlay area and opens the file as a Turbo Pascal untyped file. The file is sized and a check is made to make sure that enough dynamic memory is available to hold the entire file's contents with sufficient heap space left over to satisfy the program's other needs. (The amount of dynamic memory otherwise required by the program is determined by the constant *RequiredHeap*, which can be changed to reflect the needs of a particular program.)

If not enough dynamic memory is available, things go no further and the overlay group is left as a normal, nonmemory overlay group. Otherwise, heap space is allocated and the overlay file is read into memory. Then, the initialization procedure substitutes codes comprising a call instruction for the memory overlay handler for code that would normally call Turbo Pascal's native overlay handler. A data area that formerly contained the overlay file name is subsequently filled with information more pertinent to the memory overlay handling routine: the location and size of the heap space containing all the overlay code.

When a routine in the memory overlay group is needed, the call is directed to the new overlay handler. This handler uses the overlay offset and size location it receives to find the required code on the heap. The code is then moved into the reserved overlay area in the code segment with a string move and is executed.

The procedure that disposes of overlay storage on the heap is pretty straightforward. The only special thing it does is check to make sure that the overlay group it has been pointed to is actually a memory overlay group. If there wasn't adequate dynamic memory when memory overlay initialization was attempted, then the group might have been left as a normal overlay group and there would be no dynamic memory to free.

Limitations and Cautions

The memory overlay scheme described here does not work with overlay groups for which the overlay file is 64K or larger.

Nested overlays, created by declaring overlay groups within already overlayed functions or procedures, cannot be made into memory overlays. This would require modification of code inside overlay files or on the heap and seems beyond the bounds of what could be reliably, simply implemented.

Turbo Pascal's *OvrPath* procedure won't work in conjunction with memory overlay procedures. This is one limitation that could be overcome without too much difficulty by anyone who understands the code in MemOvrly.Inc—but it's probably not necessary because the primary utility of *OvrPath* is the placement of normal overlays on RAM disks.

MemOvrly.Inc is highly version dependent. It definitely won't work with any version of Turbo Pascal prior to 3.0, and if Borland changes its overlay handling technique in later versions, MemOvrly.Inc may have to be modified to take account of the changes.

Finally, as with overlays in general, memory overlays should be used only for functions and procedures that have been thoroughly debugged. Debugging functions and procedures placed in overlay groups can be exceptionally difficult.

Customizing MemOvrly.Inc

You may wish to customize the memory overlay handling routines in two ways. First, the initialization routine contains no input/output error checking code.

Second, when a program needs significant amounts of dynamic memory, the *RequiredHeap* constant at the top of MemOvrly. Inc should be customized to reflect that fact (as should the "minimum free dynamic memory" setting on the Turbo Pascal memory usage menu).

Availability

All the source code for articles in this issue is available on a single disk. To order, send \$14.95 to *Dr. Dobb's Journal*, 501 Galveston Dr., Redwood City, CA 94063 or call (415) 366-3600 ext. 216. Please specify the issue number and format (MS-DOS, Macintosh, Kaypro).

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(Listing begins on page 78.)

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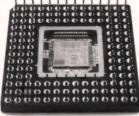
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A Unix BBS Using Shell Scripts

by Jan L. Harrington

wo years ago I assumed the responsibility of developing a BBS (now known as Scholastech Telecommunications) for Scholastech, a nonprofit organization made up of college professors and other industry professionals whose goal is to facilitate and enhance academic computing. The BBS is designed to provide a distribution mechanism for Scholastech's publicdomain/shareware software collection and to futher the exchange of information between educators at the college and university level.

Scholastech Telecommunications' software was originally developed and installed on an AT&T 7300 with 1megabyte RAM and a 20-megabyte hard disk (by the time this article is published, the system will have been transfered to an AT&T 3B2 with a 72megabyte drive). It has been in operation since October 1985, and to the best of AT&T's knowledge, it is the only BBS that has ever run on a 7300. Though the decision to use that particular machine was dictated only by circumstances (no other Scholastech hardware had the necessary configuration), the Unix operating system has provided an ideal environment in which to develop and operate a bulletin-board system. Not only is Unix multiuser and multitasking but it also contains significant support for telecommunications.

This article briefly discusses the Unix telecommunications environment and then looks in depth at the software written to implement Scho-

Jan L. Harrington, 4002 Stearns Hill Rd., Waltham, MA 02154. Jan is an assistant professor in the Computer Science Department at Bentley College. She is the author of Macintosh Assembly Language: An Introduction. Unix
has made it
remarkably simple
to develop a
bulletin-board
svstem.

lastech Telecommunications. Although it may seem like exposing the trade secret of the century, it is nonetheless true that the Unix operating system has made the development remarkably simple; the entire BBS software consists of a set of Unix V Bourne shell scripts, some supporting text files, and a public-domain XMODEM file transfer program.

The Unix Telecommunications Environment

Most implementations of Unix support both electronic mail (mail) and intersystem file transfers (uucp-Unix to Unix copy). When these functions are present, the system is able to answer and place phone calls. Unix also supports cu (call up) operations, which turn the system issuing the call into a dumb terminal. If the system called happens to be another Unix machine, cu supports file transfer as well. In terms of developing a BBS, cu is of little use because it is unreasonable to assume that all incoming calls will be from another Unix system. Cu's file transfer functions therefore cannot be used as part of BBS software.

On the other hand, the presence of facilities to manage electronic mail and *uucp* transfers means that the system can handle incoming calls from any type of system (it is nontheless true that the full power of Unix

telecommunications can only be realized when the communication is with another Unix system). The steps for preparing a Unix system for telecommunications do not vary greatly from one implementation to another, though the details of how and where the activities occur may be different. Generally, setting up a Unix system to receive incoming calls involves the following two actions:

1. Configure communications port(s): The AT&T 7300 has two internal modems; port configurations are handled by an applications shell known as The Office (because The Office performs some undocumented actions when it configures ports, users of the 7300 are wise not to attempt to do the configuration from the Bourne shell). More commonly, however, the configuration is handled from the shell by making appropriate entries in device configuration files, though the actual names of the files will vary between implementations.

2. Verify that a *getty* process is available for the port(s): *Getty* is a Unix system process that polls devices for input; it is also the process that actually logs a user onto the system. Instructions as to how the system should respond to device activity are contained in the file inittab. For example, the following inittab is used by the 7300:

is:2:initdefault

rc:bootwait:/etc/rc > /dev/window </dev/w1 2>&1

vid:2:respawn:/etc/getty window 9600

:ph0:2:respawn:/etc/getty ph0 1200 ph1:2:respawn:/etc/getty ph1 1200 000:2:respawn:/etc/getty tty000 9600 The first line in inittab puts the system in multiuser mode when the system is started; the second invokes rc, the shell script that handles system boot. The remaining lines are for actual devices. The colon in front of ph0 indicates that the first modem line is inactive and should not be polled; the line in use is ph1. The word respawn, seen in the four device lines, indicates that if the getty process is not active when the device is polled, it should be started.

For some implementations of Unix, the same port cannot be used for both incoming and outgoing calls; the getty process must be inactive for outgoing lines and active for incoming lines. Configuring such systems for outgoing calls means that a colon must be placed in inittab at the beginning of the entry for each port that will be used to dial out, as was done above for the inactive pho line. The 7300, however, kills the getty on a port automatically whenever an outgoing call is placed, making it possible to use the same line for both incoming and outgoing traffic.

Once one or more ports have been configured (assuming a modem is attached) and a getty is polling those ports, Unix answers incoming calls and takes users through the log-in process without intervention from an application program. Unix also intercepts the system log-out command, Ctrl-D, and logs users off. It is important to realize that this is not enough to configure the system for outgoing calls; outgoing calls require at least an entry in the file L-devices which describes the type of modem in use. Uucp and mail also require entries in the file L.sys, which identifies the names and phone numbers of other Unix systems that the host can call. Additional configurations may be necessary under specific Unix implementations.

What does this mean for a BBS that is to run on a Unix machine? The BBS software does not have to manage telecommunications: it does not have to poll a serial port to detect incoming calls, it does not have to log users onto the system, and it does not have to log users off the system. A simple Unix BBS needs to be concerned only with the actual functions remote users will perform after they are actu-

ally logged onto the system.

Overview of the System

The Scholastech BBS software provides the following functions:

- · file upload and download
- public message exchange
- private mail
- user help
- · change of password

Each user is given a separate Unix account under what is known as the

The steps
for preparing a
Unix
telecommunications
system
do not vary
greatly.

"restricted shell." The restricted shell prevents users from changing directories, from accessing directories not in their default *PATH* (set in a .profile, an executable shell script that is run automatically when a user logs in), and from using any commands not contained in their rbin directory. The rbin directory is also made part of the user's default path. The restricted shell therefore effectively locks BBS users into a small segment of the system, preventing them from even listing the contents of directories that are not part of the BBS.

BBS commands are each implemented as a separate shell script. There are at least four advantages to this approach:

- 1. The short shell scripts are easier to debug than a long, single program, regardless of whether the program is a shell script or compiled C.
- 2. New commands can be tested and added at any time without disturbing BBS operation.
- 3. Individual shell scripts simplify the logging of users' activities and the monitoring of their activities while they are on the system.
- 4. Shell script programming doesn't require the nearly 10 megabytes of

disk space needed by the C development system on the 7300.

There are, of course, some major disadvantages to working with separate shell scripts:

- 1. Each user must have his or her own account (as long as the number of users is small, this is not an enormous problem, but as the number of users rises, account management begins to consume significant amounts of time).
- 2. Because they are interpreted rather than compiled, shell scripts run much more slowly than would, for example, programs in compiled C.
- 3. The Bourne shell itself, although containing powerful flow of control statements, is weak in terms of arithmetic operations. Even simple addition must be prefaced with the command *expr* and enclosed in single quotes before the sum can be assigned to a variable.

Scholastech Telecommunications has two special accounts without passwords, new and info. The new account is designed only for on-line account requests; the info account displays information about Scholastech activities and can accept on-line sign-up for Scholastech workshops. Code for these functions is contained in their .profiles.

BBS Directory Structure

Scholastech Telecommunications files are maintained within the directories /u/bbs and /u/bbs/rbin and their subdirectories. The contents of the directory in /u/bbs are:

- rbin directory (contains all commands available to restricted shell users)
- MS-files directory (contains MS-DOS software)
- Unix-files directory (contains Unix software)
- Mac-files directory (contains Macintosh software)
- Uploads directory (destination for all uploaded software)
- new (log-in for on-line sign-up for new accounts)
- info (log-in for Scholastech information and on-line workshop sign-up)
- msg directory (contains support files for the public message

UNIX BBS

(continued from page 55)

subsystem)

- log.file (a text file that records BBS command usage)
- msg1 (a text file whose contents are displayed by a user's .profile)
- · directories for each BBS user

The directory /u/bbs/rbin contains:

- dwnld (shell script for downloading software)
- *help* (shell script for displaying BBS instructions)
- *list* (shell script for listing the contents of a data library)
- pmail (shell script for sending and reading private mail)
- readmsg (shell script for reading public messages)
- scan (shell script for viewing the headers of public messages)
- send (shell script for sending public messages)
- xmodem (public-domain file transfer program; compiled C)
- supporting text files:

- a. MS.list (listing of MS-DOS data library)
- b. Mac.list (listing of Macintosh data library)
- c. Unix.list (listing of Unix data library)
- d. help.file (text file with BBS instructions)
- e. user.file (listing of system users)
- all Unix commands used by the BBS shell scripts (either copied or linked into this directory)

The Shell Scripts

BBS User's .profile

The environment for each BBS user is configured by a short .profile:

PATH = /u/bbs/rbin:/u/bbs/MSfiles:/u/bbs/Mac-files:/u/bbs/Unixfiles:/u/bbs/Uploads

export PATH PS1='>' cat /u/bbs/msg1 echo

The .profile establishes the rbin directory and the program library di-

rectories as the user's default *PATH* (although technically there is only one restricted shell, /bin/rsh, there can be many rbin directories within a single Unix file system, each of which has a different path name and contains a different set of commands). The default prompt is also changed from the standard Bourne shell \$ to a > . Finally, the .profile displays the contents of a short welcome message (stored in msg1) and then returns control to the operating system.

Users logging in for the first time or after a long absence usually use the *help* command to display a screen full of instructions. The *help* command shell script is only two lines long:

echo 'who am i | cut -f1 -d" " ' 'date | cut -c1-c16' "help" >>/u/bbs/log.file

cat /u/bbs/rbin/help.file

The first line writes a record to the command use log file, /u/bbs/log .file; the second simply displays the

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contents of a text file containing instructions.

Users who wish to change their passwords can issue the command passwd. This Unix command is not contained in a shell script but is made available to BBS users at their > prompt. Passwd prompts for the new password and then asks that it be entered again for verification. On the 7300, passwd also enforces format constraints—passwords must be two words separated by a special character (only passwords established by the superuser can violate those constraints).

The Program Libraries

Three commands support user interaction with the program libraries—list, which displays the contents of a program library; dwnld, which downloads a file; and upld, which uploads a file.

The command list (Listing One, page 80) uses the Unix command more to display the contents of a text file that contains a listing of files available in a specific program library (lines 6, 11, and 16). Although the more widely used cat command also displays the contents of a text file, more displays the text one screen at a time (users advance by pressing the space bar or Return key) and is therefore better suited for screen displays. The Unix command to display a file directory, ls, could also be used to show the contents of the program libraries directly; this would eliminate the need to maintain the separate text files. In terms of security, however, Is is a "dangerous" command. If ls is present in the rbin directory, users who know Unix can use it from the > prompt to view the contents of that directory. They can then be aware that potentially more dangerous commands such as my, which moves a file from one name or place to another, are available under the restricted shell. The best situation is to avoid the use of dangerous commands entirely, but they are essential to more than one of the BBS shell scripts.

The listings for each of the three program libraries are kept in separate files. Users must specify which program library as an argument to the command. For example, *list MS* displays the contents of /u/bbs/

rbin/MS.list. The shell script traps the argument at lines 2, 7, and 12. If no argument is present (line 17), the shell script prints an error message (line 18) and exits. *List*, like all the other BBS commands, also writes a record to the command use log file (lines 4, 9, and 14).

Dwnld (Listing Two, page 80) transfers files from the program libraries to a remote user. The command requires two arguments—a designator for the program library and the name of the file. For example, dwnld MS PCFILE.ARC initiates transfer of the file /u/bbs/MS-files/PCFILE.ARC. The first 12 lines of the dwnld script identify the program library. If no program library is included, the script prints an error message (line 11) and does not execute a transfer. Assuming that a valid program library has been entered, dwnld then verifies that a file name is present (line 13); the error message for a missing file name appears in lines 37 and 38. The presence of a file name on the command line, however, does not guarantee that the file exists. A check for a valid file occurs on line 15 (the error message for an invalid file name appears in lines 32–34). After verifying that the requested file exists, dwnld writes a log record (line 17) and then prompts the user for the file transfer method (line 18). Files can be transferred as ASCII text (lines 24–29); the script simply cats the file to the user, who must have instructed his or her terminal emulator program to capture incoming data to disk.

Binary file transfers are performed using the XMODEM file transfer protocol (line 22). Xmodem, the program used to implement the transfers, is a stand-alone file upload and download program that can be executed from the Unix command level as well as from within a shell script. Two versions exist in the public domain, umodem and uc, both of which are available as C source code. Xmodem was obtained by downloading the umodem source code from an existing Unix BBS and then compiling it on the 7300 (this was done prior to loading the program libraries,

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when there was still enough space of the disk for the entire development system). The xmodem command line is straightforward: the argument -sb indicates that a file is to be sent and that its format is binary; the second argument is the path name of the file to be transmitted.

Upld (Listing Three, page 81) handles file transfers in precisely the opposite manner to dwnld. ASCII uploads echo incoming characters to a text file (lines 14-19); to upload binary files, xmodem is invoked with the arguments -rb, receive binary (line 12). Upld also contains code to verify that a name for the file to be uploaded has been entered as an argument to the command (lines 1 and 26-27): verifies that the file name for the incoming file doesn't duplicate an existing file name in the Uploads directory (lines 1-5); and collects a description of the file (appended to the text file /u/bbs/Uploads/Doc-.file), which the sysop can use to quickly figure out the nature of uploaded files.

The Public Message Subsystem

Three shell scripts support the exchange of public messages—scan displays the headers of all current messages, readmsg displays individual messages by number, and send sends a public message.

Scan is a short script that does nothing more than display the contents of a special text file, .index, to which an entry is added whenever a message is sent:

echo 'who am i ¦ cut -f1 -d" " ' 'date ¦ cut -c1-c16' "scan" >>/u/bbs/ log.file

echo echo more /u/bbs/msg/.index echo

Scan's purpose is to allow users to identify the numbers of messages that might be of interest.

Messages are displayed on the screen by the command *readmsg* (Listing Four, page 81). *Readmsg* uses two simple files—/u/bbs/.first, which contains the number of the first message available, and /u/bbs/

.last, which contains the number of the last message available—to remind users who might not have scanned the message index of valid message numbers (lines 1–9) to do so. The file names of messages stored in /u/bbs/msg are the same as their message numbers. For example, message number 64 has the file name /u/bbs/msg/64. Therefore, once the user enters a message number (line 12), readmsg can simply cat a file by that name (line 18). The script also traps nonexisting message numbers (lines 19–21).

The shell script send (Listing Five. page 82) stores messages for the public message system. The script increments the last message number (line 8) and then conducts a dialogue with the user to collect the message header (lines 11-25) and the body of the message (lines 27-34). In addition, the message header must be added to the top of the file /u/bbs/msg/.index (that is, scan displays message headers in descending order by message number). To maintain the descending order, the headers are first written to a temporary file (lines 35-33). Then, the existing .index file is concatenated onto the end of the temporary file (line 39). Finally, the temporary file is moved on top of the old index, effectively erasing the existing file (line 40). Because public messages can be addressed to groups of people (for example, ALL or All 3B2 Users), send does not bother to verify that the addressee is a valid system user ID.

The Private Mail Subsystem

Private mail is built around the Unix mail command. The shell script pmail (Listing Six, page 82) builds an interface to aid users in sending and reading mail. The script first displays the four available commands (lines 4–11): s to send mail, r to read mail, l to see system users, and x to exit from pmail.

The *l* command is not a command normally associated with Unix mail, but sending Unix mail requires an argument of either a valid user ID or the electronic-mail name of a remote system known to the system from which the mail is sent. Unfortunately, user IDs are often unrelated to users' real names. It is therefore important for the BBS to provide some way

for users to associate system user IDs with human names. Pmail displays the contents of the text file /u/bbs/ rbin/user.file, which contains a listing of user names and user IDs (lines 61-63). The same information is also available from the system file /etc/ passwd. In other words, it might be possible to avoid having to maintain user.file by using cut to obtain specific fields from /etc/passwd. The machine that supports Scholastech Telecommunications, however, also contains several accounts that are unrelated to the BBS. The code required to grep out the accounts that should not be displayed to BBS users far exceeds the effort to maintain user.file.

Sending Unix electronic mail is straightforward. The *mail* command is used with a single argument—the user ID of the recipient (although Unix mail can be directed to remote Unix systems, Scholastech Telecommunications does not make that capability available to BBS users). For example, *mail* sysop sends mail to the sysop account on the same system. To send private mail, the *pmail* shell script:

- displays a set of instructions that help ensure that the mail will be sent successfully (lines 19–29)
- collects the user ID of the account to which the mail is being sent (lines 31– 32)
- verifies that the user ID is valid (lines 34–39)
- sends the mail by issuing the Unix mail command (lines 40–45)

Users working at the Unix command level can read mail by issuing the command mail without any arguments (there is no reason why BBS users who are knowledgeable about Unix cannot do so from their restricted shell > prompt). Pmail reads mail in exactly that way (line 55). Several options are available once a piece of mail has been displayed on the screen, including forwarding the mail to other users and saving the mail under a different file name in the user's account. Pmail, however, only alerts the user of the options of either deleting the message or leaving it in the user's mailbox (lines 51-54). Though knowledgeable Unix users may attempt to save the mail, BBS

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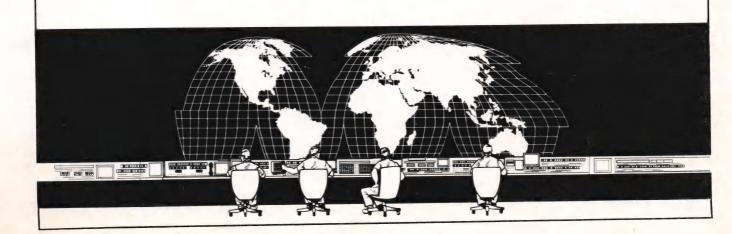
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(continued from page 58)

users do not have the right to write to their own accounts (this restriction was included to conserve disk space); any attempt to save mail in a user's account will therefore fail. Knowledgeable users may, however, succeed in forwarding mail to other BBS users.

Signing Up for a New Account

The special account new exists on Scholastech Telecommunications without a password. Its sole purpose is to permit prospective users to request accounts on-line. The code that manages the sign-up process is contained in the account's .profile (Listing Seven, page 86). Ideally, this .profile should automatically log users off the system once they have finished. The only way to do so that works on the 7300 (stty 0 does not work) is to make the .profile the account's default shell. There is, however, an important reason why the account wasn't set up in that way; any shells other than /bin/sh (the standard Bourne shell) and /bin/rsh (the standard restricted shell) are assumed to be extremely restricted. In particular, they do not permit any redirection of output. That means that data cannot be sent to a disk file, something that is essential if user account request data are to be stored. Instead, the new account is established as a standard restricted shell account. Its default path is set to an rbin directory that contains only the two harmless commands echo and cat (line 1); its level 1 prompt is set to the phrase "Type CNTRL-D now" (line 2). When the .profile ends or if a user breaks out of the .profile with the delete key, even those who know Unix will be powerless beause the commands in the rbin directory permit nothing but the display of text to the screen. The prompt continually reminds them of the key sequence to exit from the system.

Though somewhat lengthy for a profile, the account request shell script is quite simple. It first verifies that the user wants to sign up for an account (line 5). Assuming the user does want an account, the script then enters a loop to collect name and address data (lines 13–43). After data

have been entered, the user is given a chance to view the data (lines 33–38) and to correct it if necessary (lines 42–43). This process is repeated for the user ID and initial password (lines 47–60). Note that the system's password format is not enforced during the account request process. This is because initial passwords will be established by the superuser, bypassing the format restrictions. Finally, the script writes the data to a text file (lines 63–71).

The new account's profile does not actually establish new accounts.

I wrote two routines to keep records of system and command use.

This must be done manually by completing the following steps:

- 1. Make an entry in /etc/passwd for the new user.
- 2. Enter a password for the new user.3. Create a directory for the new user.
- 4. Copy the BBS .profile into the new user's directory.
- 5. Make an entry in /u/bbs/rbin/user.file for the new user.

Steps 1 and 2 must be performed by the superuser. The remaining steps are performed by the system operator (the account sysop). The sysop account owns all BBS accounts and their files. This prevents BBs users from viewing and/or modifying their own .profiles, an additional security measure often used along with the restricted shell.

Providing Organizational Information

The shell scripts that you have seen up to this point are generic; they might be used to support just about any BBs. The info account, however, provides a function needed by Scholastech that isn't typically found with most bulletin boards. Info's .profile (Listing Eight, page 88) works much like the new account's .profile to display information about upcoming Scholastech events and to allow users to register for workshops online.

Like new's .profile, the script first establishes a special rbin directory as the account's default path (line 1); info's .profile contains nothing but three text display commands—cat, echo, and more. It also changes the level 1 prompt to "Type CNTRL-D now . . ." (line 2). It then prints a welcome heading (lines 4–5) and displays the contents of a text file with the welcome message (line 7). The actual information about upcoming events is stored in the more extensive info .file, which is displayed page by page using the Unix command more (line 11).

Workshop sign-ups are straightforward. Users are prompted to enter registration data (lines 20–37). The profile then appends that data to a text file (lines 38–43). The registration file is later printed and turned over to the Scholastech personnel in charge of the workshops.

System Monitoring

It is important for any BBS sysop to be aware of everything that is happening on the system. While users are logged in, their activites can be monitored with the Unix ps command. which displays the status of all active processes. There is, however, a further need to keep records of system and command use. Unfortunately, the 7300's implementation of Unix V does not include the standard Unix accounting functions. I therefore wrote two routines that display and reformat data from a system log file and from the log file written by each BBS shell script in order to obtain archival information.

The two shell scripts are *logs*, which displays data to the screen, and *usage*, which creates a printed report. Both rely on process beginning and ending data kept by Unix in the system file /etc/wtmp. Normally, /etc/wtmp is cleaned out every time the system is rebooted. This can, however, result in a loss of data whenever an unexpected power

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add. list_connections(Place) (connect(Place,X), tab(2)

list_connections().

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Rules can refer to other rules, as well as to facts. Therefore, we can write "list_connections/1" just like "list_things/1".

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(continued from page 60)

failure occurs. I therefore made a modification to the system routine *.cleanup*, which is executed by the system boot routine *rc*, and turned the line > /etc/wtmp into a comment to prevent its execution (I could have just as easily deleted it).

Logs (Listing Nine, page 90) first creates a temporary file (/u/user.log) that contains the user's name and the time logged on for each remote user with the Unix command who. A pipe to the Unix command grep extracts all entries that contain ph1 (line 1). Because the device name for the internal modem being used by the BBS is ph1, the grep will include log-ins on that line and leave out all other devices (the console is usually either w1 or w2, for example). If the logs command is issued with an argument a (line 2), logs displays the entire contents of both the temporary file (line 8) and the BBS command log file (line 15). If no argument is included, logs displays only the last ten lines of the two files (lines 21 and 26). Finally, logs removes the temporary file it created (line 29).

Usage (Listing Ten, page 90) creates two printed reports. The first is a columnar listing of the user's name, the date, the time on, and the time off for each log-in contained in /etc/wtmp. The second is simply a hard copy of the BBS command use file. Usage is a "dangerous" shell script; one of its functions is to reinitialize /etc/wtmp (line 17). It is therefore extremely important that no user other than the system operator should have any rights to the program (that is, its permission line should appear as -rwx-----). Why shouldn't other users be allowed to read usage? Read rights allow another user to make of copy of a program to put in their own account. Once the copy is made, the user becomes the owner and can execute the script, wiping out /etc/ wtmp and destroying system use

The logic behind *usage* is tied to the way in which data are stored in /etc/ wtmp. There is one record for every process that starts and every process that ends (the two exceptions are the processes *LOGIN* and *rc*). The trick to capturing log-on and log-off times is

therefore to match up the entries for any given user. Because the user running the program will not have logged off, the program must be careful to exclude that user from the report. The particular version of usage that appears in Listing Ten handles only remote users on the device ph1. In this case, the code doesn't need to worry about excluding the system operator who is running the program from the console. If, however, users on other devices are to be included, two things must happen:

It is extremely important that no user besides the sysop has rights to usage.

- 1. The code for assembling the report must be duplicated for each device (records for concurrent users are interleaved in /etc/wtmp).
- 2. The name of the user running the program must be excluded from the report because there will be no matching process termination record for that user.

The usage script first creates the system use report. It writes report headings to a temporary output file (lines 1-3) and then begins to assemble the body. The body is assembled in the following manner:

- 1. A log-on record for each user entering the system on *ph1* is extracted and stored in the file temp1 (line 4).
- 2. The first field of the record, the user name, is *cut* from temp1 and stored in temp2 (line 5). This forms the leftmost column of the report.
- 3. The log-on time is *cut* from temp1 and stored in temp3 (line 6). This forms the middle column of the report.
- 4. A log-off record for each user entering the system on *ph1* is extracted. The log-off time is *cut* from the record and stored in the file temp4 (line

- 7). This forms the rightmost column of the report.
- 5. The three temporary files containing the contents of the columns of the report are put together side by side with *paste* (line 8).
- 6. To complete the process, the report is formatted for output with the Unix *pr* command (line 9). The temporary files are removed (lines 10–14).

Records in the BBS command use log file (/u/bbs/log.file) are already in columnar format (they were written in such a manner by each BBS shell script). Therefore, they are simply formatted for output with *pr* (line 15). The log file is also reinitialized (line 16). Both reports are then sent to the printer (lines 18 and 19). Note that *usage* does not delete the two final output files because doing so before the files have actually been printed can disrupt the printing process.

Acknowledgment

I received both the XMODEM file transfer program and a great deal of help from David Watson, who operates a Unix/Xenix BBS in my area. Though I thanked him profusely at the time, I think he deserves a public acknowledgment for his generosity and patience in putting up with some of the elementary (dumb) questions I asked when I was first working on my BBS software.

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(Listings begin on page 80.)

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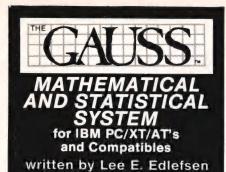


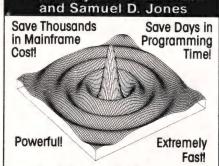
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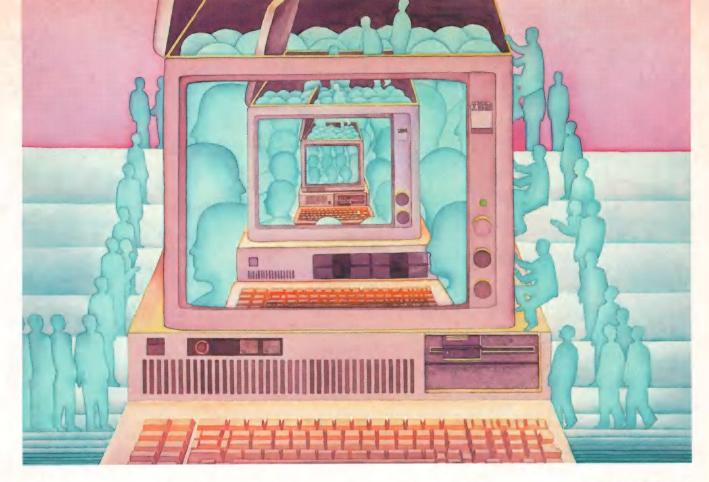
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COM PORT DRIVER

```
Listing One (Text begins on page 42.)
```

```
extended int 14 handler
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text
                  segment word public 'code'
                  assume cs:_text, ss:_text, ds:_text, es:_text
start:
                                                                                           ; must be first
                  Imp
                  misc constants
                                                                                             buffer size
input flow control turn off point
input flow control turn on point
select DTR input flow control
select RTS input flow control
select XON/XOFF input flow control
bit to set OUT2
CTS not required to transmit
DSR not required to transmit
set baud rate to 19200
set baud rate to 38400
com1 mask for 8259
com2 mask for 8259
BUFSZ
                  edn
                                    0100h
                                    00F0h
00E0h
XOFFSZ
                  equ
                  edn
edn
edn
DTR
                                    001h
RTS
                                    002h
XNXF
OUT2
                  equ
equ
                                    008h
NOCTS
                                    010h
                                    020h
                  equ
B19200
                  equ
                                    040h
B38400
                                    080h
                  equ
C2MASK
                  equ
                                    008h
                                                                                              com2 mask for 8259
                  structure for port control blocks
pcbs
                  struc
putptr
getptr
bufcnt
                                                                                              points to last char put
points to next char to get
number of characters in buffer
                                                                                              number of characters in burre
address past buffer end
port address
timeout outer loop
mask to enable port interrupt
set if recv has been stopped
bufend
                  dw
pbase
timeout
                  db
mask
                  db
rxoff
opts
                  db
                                                                                              holds extended options holds character attributes
lchar
                                                                                              old interrupt enable register
old line control register
old modem control register
old baud low divisor latch
old baud high divisor latch
oldler
                  db
oldmor
olddll
olddlm
                  db
buf
                                    BUFS2 dup (?)
                  allocate storage for port control blocks pcbs 2 dup (<>)
pcb
                                                                                           ; array of port control blocks
oldOB
                                                                                           ; old vector for int0B; old vector for int0C; old vector for int14
oldoc
                                    ?
oldmak
                  db
                                   table for UART initialization 0417h
bauds
                                                                                              divisor for 110 baud
divisor for 150 baud
divisor for 300 baud
                  dw
                                    0300h
0180h
00C0h
                                                                                              divisor for 300 baud
divisor for 600 baud
divisor for 1200 baud
divisor for 2400 baud
divisor for 4800 baud
divisor for 9600 baud
                  dw
                                    0060h
0030h
                                    0018h
                  dw
                                    000Ch
baudl 9
                                                                                              divisor for 19200 baud
                                                                                              divisor for 38400 band
......
      0Bh & 0Ch (com2 & com1) interrupt handler
put characters into the appropriate buffer
input flow control - stop input when buffer nears full
int OB:
                  : point to com2 port control block with ds:si
                  push
                  push
                  mov
                                    ds, si
                                    si, pcb + size pcbs
short al0
                  jmp
                 : point to com1 port control block with ds:si push ds push si
                                    ds
si
si, cs
                  mov
a10:
                  push
                                    di
                  push
                                    dx
                  push
                                    bx
                  push
                  ; read character, if this fills buffer to XOFFSZ, then send ; XOFF or drop DTR and/or RTS as indicated by extended options
```

(continued on page 66)



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COM PORT DRIVER

```
Listing One (Listing continued, text begins on page 42.)
```

```
; buffer put pointer
                                     di, [si].putptr
dx, [si].pbase
                                     al, dx
ah, al
[si].bufcnt, XOFFSZ
                                                                                            ; read character and clear interrupt
                  cmp
jl
cmp
jne
                                                                                            ; jump if buffer below turnoff point
                                                                                           ; jump if receive already disabled ; indicate receiver now disabled
                                     [si].rxoff, 1
[si].opts, XNXF
                  test
jz
mov
                                     al, 's' - 'e'
                  out
                                                                                           ; send ^S
a20:
                                    bl, [si].opts
bl, DTR or RTS
                  mov
                                    bl,
a30
                  and
                  jz
add
                                    dx, 4
al, dx
bl
                  in
                  not
                  and
                                     al. bl
                                                                                            ; drop DTR and/or RTS
a30:
                  ; if the buffer is full, set overflow status and exit cmp [si].bufcnt, BUFSZ
                  cmp
jl
or
                                    word ptr [di], 0200h
short a99
                  imp
a40 ·
                  ; read the port status
                                    dx, 5
al, dx
                  in
                  ; get new buffer pointer, adjust for wrap around if required
                  inc
                  ine lea
                                     di, [si] .bufend
                                     di, [si].buf
a50:
                  ; store new buffer pointer and put status:data into the buffer
mov [si].putptr, di
xchg ah, al
mov [di], ax
inc [si].bufcnt
a99:
                  ; send interrupt complete command to interrupt controller mov al. 020h
                                    al, 020h
020h, al
                  out
                  pop
pop
pop
                                    bx
                                    dx
di
 14h interrupt handler
                                     emulate pc bios functions plus extensions
                   dx is used to select the port for all calls
    dx = 0 for coml, dx = 1 for com2
                                    initialize the port, parameters in al al bits 7-5 set the baud rate 000 - 110, 001 - 150, 010 - 300, 011 - 600 100 - 1200, 101 - 2400, 110 - 4800, 111 - 9600 al bits 4-3 set the parity 00 - none, 01 - odd, 11 - even al bit 2 is number of stop bits 0 - 1 stop bit, 1 - 2 stop bits al bits 1-0 set the word length 00 - 5, 01 - 6, 10 - 7, 11 - 8 returns ah 1 al as per comm status (ah-3) below example: ah - 10101110 is 2400 baud, odd parity, 2 stop bits, 7 bit characters
                  ah = 1 transmit the character in al, al preserved
returns ah as per comm status (ah-3) below
                                     return receive character in al
                                     returns ah as per comm status (ah=3) below
only bits 1-2-3-4-7 can be set
ah having any bits set is a receive error or timeout
                  ah = 3
                                     return comm port status in ah & al
                                     ah bits are:

b0 - data ready b1 - overrun
b2 - parity error b3 - framing error
b4 - break detect b5 - xmit holding reg empty
b6 - xmit shift reg empty b7 - timeout
                                        b0 = delta CTS
b2 = trail edge ring
b4 = CTS
b6 = ring indicator
                                                                              bl - delta DSR
                                                                                                   b3 - delta recy detect
                                                                                b5 - DSR
                                                                                                   b7 - recv signal detect
                                     extended options returns 05A5A in ax, used to identify excom al contains options as follows:
```

(continued on page 70)

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CIRCLE 217 ON READER SERVICE CARD

COM PORT DRIVER

```
Listing One (Listing continued, text begins on page 42.)
```

```
bit 0 enables DTR input flow control
bit 1 enables RTS input flow control
bit 2 enables XON/XOFF input flow control
bit 4 don't require CTS to transmit
bit 5 don't require DSR to transmit
bit 5 don't require DSR to transmit
                                bit 6 sets baud rate to 19200
bit 7 sets baud rate to 38400
                             remove excom, restore old vectors, release memory
int14:
               push
               push
push
               push
               push
               ; point
                          to selected com port control block with ds:si
                             bx, cs
ds, bx
si, pcb
dx, dx
b00
               mov
               or
                              si, size pcbs
              ; get port address in dx, return if port not installed mov dx, [si].pbase
                             dx, [si].pbase
dx, dx
               or
jz
               ; ah has request type, jump to selected routine
               jz
dec
                                                                          : initialize
                              ah
               jz
dec
                              b100
                                                                          : transmit a char
               jnz
jmp
                             b10
                              b200
b10:
               dec
                              ah
               jnz
                                                                          ; get status
b20:
               jnz
               mp
                             b400
                                                                          ; extended initialize
b30:
                              ah
               jnz
jmp
                              b40
                             b500
                                                                          ; remove excom
b40:
                             bx
cx
              pop
pop
                             dx
                ah = 0 initialize the port
b000:
              push
                initialize the pcb, this empties the input buffers ea ax, [si].buf
                             ax, [si].Dur
[si].putptr, ax
              mov
              inc
              inc
                             [si].getptr, ax ax, BUFSZ * 2 [si].bufend, ax
              mov
                             ax, ax
[si].bufcnt, ax
[si].rxoff, al
              aub
              mov
              ; assert add
                           flow control signals and enable port interrupts
```

dx, 4 al, DTR or RTS or OUT2 dx, al mov out cli

in and

al, 021h al, [si].mask 021h, al out sub

; unmask com port on int controller dx, 3 al, 1 dx, al

; enable receive int on wart set up baud rates and character attributes

pop bl, al bl, al al, 01Fh [si].lchar, al al, [si].opts al, B19200 or B38400 and mov

and jz call b010 b440 short b020 jmp

cl, 4 bl, cl bx, 0Eh

; init only character attributes

; isolate baud rate selection

(continued on page 72)

b010:

rol

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Use C-scape's functions for Lotus-like, pull-down, or your own menu designs, automatic scrolling, pop-up windows (number limited only by RAM), logical colors, help, time and date, yes/no, tickertape fields, secure and protected fields, and many others, to turn your demo into a fully functioning and complete program in a fraction of the time spent coding screens from scratch.

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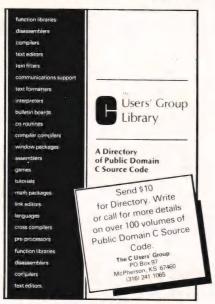
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CIRCLE 105 ON READER SERVICE CARD

COM PORT DRIVER

```
Listing One (Listing continued, text begins on page 42.)
                                                           ; init baud and character attributes
            sub
jmp
                                                            ; return status
            : ah - 1
            ; transmit the character in al
b100:
            push
add
mov
                        dx, 6
bh, [si].opts
            not
                        bh
            and bh, NODSR or NOCTS; optionally wait for DSR and/or CTS call bl30
            jnz
                        b110
                        dx
            mov
                        bh. 020h
            ; wait for call
                         the transmitter to be ready
                       b130
            jnz
sub
                        b110
                        dx, 5
            pop
                        ax
            push ax; actually send the character
                       dx, al
b310
            out
            call
                        short bl20
            dmc
b110:
            ; restore al and indicate a timeout (bit 7 of status, ah)
            call
                        b310
                        ah, 080h
            or
b120:
            pop
mov
dmt
                        cx
al, cl
b40
            ; wait for selected status, bh is status, dx is port address
b130:
                        bl, [si].timeout
                                                           ; delay counter
b140:
            sub
                        CX, CX
b150:
            in
                        al, h
al, h
b160
            and
             стр
            jz
loop
                                                            : status match
                        b150
bl
            dec
                        b140
                        bh, bh
                                                            ; timeout, set non-zero
b160:
            ret
            : ah - 2
            ; receive a character, put it in al
b200:
                        [si].bufcnt, 0
b230
                                                           ; chars in buffer, return one
            ; wait for a character to appear in the buffer, or return timeout
                       al, [si].timeout
al, al
            add
                                                           ; shorter timeout loop, so loop more
b210:
            aub
                       CX, CX
b220:
                        [si].bufcnt. 0
            amp
            loop
                        b220
            dec
            jnz
mov
                        b210
                       ax, 08000h
                                                           ; timeout return value
            jmp
b230:
            ; if receiver is disabled and room now exists, enable receiver
            je
cmp
                       [si].rxoff, 0
b250
                                                           ; receive is already enabled
                        [si].bufcnt, XONSZ
            jg
test
                        b250
                                                           ; buffer fuller that turn on point
                        (si).opts, XNXF
            jz
                        b240
                        [si].rxoff, 0
al, 'Q' - '@'
dx, al
            mov
                                                           ; indicate receiver enabled
            mov
            out
                                                           ; send ^O
b240:
            : it's
                    easier to just assert DTR & RTS rather than see if needed
                       dx, 4
al, dx
al, DTR or RTS
            add
            or
            out
                                                           ; assert DTR and RTS
            ; get the char from the buffer and update buffer get pointer call b310
                       b310
ah, 01Eh
            and
                                                           ; status reported with a receive
            inc
            inc
                            [si].bufend
            ine
                       b260
                        bx, [si].buf
b260:
```

```
[si].getptr, bx
[si].bufcnt
                                                                   ; updated pointer
             dec
             ; ah = 3
             ; get port status
b300:
             call
              inc
                                                                   ; modem status
              in
              imp
             ; get line status from input buffer and/or hardware into ah
b310:
                           dx, [si].pbase
                           dx, 5 al, dx
              add
             in
                                                                  : line status
                           ah, al
[si].bufcnt, 0
              cmp
je
             ; when chars in buffer, fix status as per status in buffer; return as above plus bx = getptr, al = character
mov cl, al
mov bx, [si].getptr
mov ax, [bx] ; status:data (ah
mov ch, 0lEh ; these status bi
                                                                   ; status:data (ah:al) ; these status bits from buffer
              and
                           ah, ch
              not
                           ch
                            cl, ch
              or
                            ah, cl
ah, 1
                                                                   ; char in buffer, show data ready
b320:
              ret
              ; ah = 4
; extended initialization
b400:
              ; save the options in the pcb
mov [si].opts, al
              ; if an extended baud rate, set up the uart add dx, 3
                            dx, 3
al, B19200
              test
                            b410
                            bx, baud19 - bauds
b430
              call
                            short b420
               jmp
 b410:
                            al, B38400
              test
                            bx, baud38 - bauds
b430
              call
 b420:
                                                                 ; magic value to identify excom
               ant
               ; init the baud rate and character attributes
 b430:
                            al, 080h
                            dx, al
ax, [bx+bauds]
dx, 3
dx, al
               out
                            ax,
                                                                    ; set low order divisor on uart
               out
               inc
                            dx
               mov
                                                                    ; set high order divisor on uart
               inc
                            dx
 b440:
                            al, [si].lchar
               mov
                                                                     ; set character attributes
               out
               : ah - 5
               ; remove excom
  b500:
               push
               ; restore status of interrupt controller and uarts cli
                             ah, oldmak
ah, ClMASK or C2MASK
al, 02lh
               and
               in
                             al, not (CIMASK or CZMASK)
                                  ah
               or
                             al, ah
021h, al
                out
                sti
                             si, pcb
b510
                                                                     ; com1 pcb ; restore com1 uart status
                lea
                             si, pcb + size pcbs ; com2 pcb
b510
                                                                     ; restore coml wart status
                             short b530
                dmp
                ; subroutine to restore uart status
  b510:
                             dx, [si].pbase
dx, dx
                                                                     ; no port
                             b520
                                                                            (continued on next page)
```

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Annual 68K Issue, 68K Mini Forth, OS-9 Operating System, Mac and Amiga Interface Programming.

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out

COM PORT DRIVER

```
Listing One (Listing continued, text begins on page 42.)
                            dx, 3
al, 080h
dx, al
                            dx, al
dx, 3
al, [si].olddll
              out
                                                                    ; set access to baud divisors
               aub
              out
                            dx, al
                                                                    : band divisor low
              inc
                           dx al, [si].olddlm
              mov
              out
                            dx.
                                                                    : baud divisor high
                            dx, 2
al, [si].oldlcr
               add
              out
                            dx. al
                                                                    ; line control register
                           dx, ai
dx
al, [si].oldmcr
dx, al
dx, 3
al, [si].oldier
              inc
              out
                                                                    ; modem control register
               sub
              out
                                                                    ; interrupt enable register
 b520:
 b530:
              ; restore old int 0Bh vector lds dx, old0B mov ah, 025h mov al, 0Bh int 021h
              ; restore old int OCh vector
              pop
                            da
                           dx, oldOC
ah, 025h
al, 0Ch
              mov
              int
                            021h
              : restore old int 14h vector
              pop
                           ah, 025h
al, 14h
              mov
               int
              ; free excom's memory, this memory
              push
                           es
ax, cs
                           es, ax
ah, 049h
021h
              mov
              mov
 installation entry point, must be at file end initialize and stay resident
 excom:
              ; initialize constant parts of each pcb
                           ax, 040h
es, ax
di, di
                                                                      bios data segment
                                                                      offset to com1 base address offset to com1 timeout
              sub
                           bx, 07Ch
cl, not ClMASK
              mov
                                                                      coml interrupt mask
                           si, pcb
c00
di
                                                                    ; com1 pcb address; init com1
              lea
              call
                                                                    ; offset to com2 base address; offset to com2 timeout
               inc
                            di
              inc
                            cl, not C2MASK
                                                                      com2 interrupt mask
              add
                                                                    ; com2 pcb address
; init com2
                            si, size pcbs
              call
              jmp
              ; subroutine to initialize parts of a pcb
 c00:
                           dx, es:[di]
[si].pbase, dx
al, es:[bx]
[si].timeout, al
[si].mask, cl
                                                                   ; port base from bios
              mov
              mov
                                                                   ; copy timout from bios ; save interrupt mask
              mov
 c10:
              ; save old status of interrupt controller and uarts in al, 021h mov oldmak, al lea si, pcb ; coml pcb call c20 ; save old
                           c20 ; com1 pcb
; save old com1 uart status
c20
c20 ; save old com1 uart status
              lea
call
              amt
                           short c40
              ; subroutine to save uart status
 c20:
                           dx, [si].pbase
              or
jz
                                                                   ; no port
                           c30
dx, 3
al, dx
al, 07Fh
[si].oldler, al
al, 080h
dx, al
dx, 3
              add
              and
mov
mov
                                                                   ; line control register
```

; set access to baud divisors

```
al, dx
[si].olddll, al
dx
                mov
inc
in
                                                                                 ; baud divisor low
                                al, dx
[si].olddlm, al
                mov
                                                                                 : baud divisor high
                                dx, 2
al, [si].oldler
dx, al
                mov
                                                                                 : restore register access
                out
                inc
                                                                                 ; modem control register
                                 [si].oldmcr, al
                mov
                 sub
                in
                                 [si].oldier, al
                                                                                 ; interrupt enable register
                mov
c30:
                ret
C40:
                ; release environment memory, not used
                mov
mov
                                bx, 02Ch
ax, [bx]
                                                                                 : address of env
                mov
                                es, ax
                                      049h
                                                                                 ; free memory
                int
                ; save old int OBh vector, install new handler push ds
                mov
mov
                                ax, cs
                                 ds, ax
ah, 035h
al, 0Bh
                 mov
                 int
                                 021h
                                 word ptr old0B, bx
                mov
                mov
                                 ax. es
                                 word ptr old0B + 2, ax dx, int0B ah, 025h al, 0Bh
                 mov
                mov
                 mov
                 ; save old int OCh vector, install new handler mov ah, 035h mov al, OCh
                 mov
                 mov
                                 021h
                                 word ptr oldoc, bx
                 mov
                 mov
                                 ax, es
                                 word ptr oldOC + 2, ax
dx, intOC
ah, 025h
al, 0Ch
                 lea
                 mov
                 mov
                 ; save old int 14h vector, install new handler
                                 ah, 035h
al, 014h
                 mov
                 mov
                                  word ptr old14, bx
                 mov
                                 word ptr old14, bx
ax, es
word ptr old14 + 2, ax
dx, int14
ah, 025h
al, 014h
021h
                  lea
                 mov
mov
int
                  pop
                  ; exit and keep everything above the entry point 'excom'
                                  dx, excom
cl, 4
dx, cl
                 lea
                  ahr
                                  dx
                                  ah, 031h
al, 0
                  mov
                  mov
                                                                                   ; terminate-and-stay-resident
                                  021h
  _text
                  ends
                  end
endm
                                   start
                                                                                                             End Listing One
   Listing Two
    #include
                   <atdio.h>
                   (C) 1987, Crystal Computer Consulting Inc.
This software may be used freely, at your own risk,
as long as this notice is not removed.
                                                                    /* ah value for extended init */
/* magic value to identify excom */
/* DTR input flow control */
/* RTS input flow control */
/* RTS input flow control */
/* EXD flow flow control */
/* CTS not required for transmit */
/* DSR not required for transmit */
/* set baud rate to 19200 */
/* port was selected */
    #define
#define
#define
                   EXINIT
                   EXCOM
DTR
RTS
                                    0x5A5A
                                    0x01
0x02
    *define
    *define
                    XNXFIN
NOCTS
                                    0×04
                                    0x10
                    NODSR
                                    0x20
    *define
                    B19200
                                     0×40
                    B38400
COMINIT
                                     0x100
    #define
    typedef struct {
                                                                     /* command string */
                    void
                                                                     /* command to execute */
                                     (*fnct)();
                    int
                                                                     /* argument to pass */
/* exmode command table */
                                    arg:
    ) CMDS;
    char
                    *insmsg = "excom installed",
*remmsg = "excom not installed",
                                                                                               (continued on next page)
                     *helpmsg;
```

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COM PORT DRIVER

```
Listing Two (Listing continued, text begins on page 42.)
                                                              /* extended init for com1 */
/* extended init for com2 */
/* port to initialize */
                init2 - 0,
portnum - -1;
void
                exit(), install(), remove(), exinit(), setport();
                "remove", remove, setport, "com1", setport, dtr", exinit, exinit,
                                                              DTR,
RTS,
XNXFIN,
                                "xnxfin",
                                "nocts",
"nodsr",
"19200",
"38400",
                                              exinit,
                                                              NOCTS,
NODSR,
                                              exinit,
                                                              B19200,
                                              exinit.
                               NULL
main(argc, argv)
register char
                **argv;
                register CMDS
                                *cmdp;
               /* is excom installed */
if (intl4(EXINIT, 0, 0) -- EXCOM && intl4(EXINIT, 0, 1) -- EXCOM)
    helpmag = insmag;
                               helpmsg - remmsg;
               if (argc == 1)
     help(helpmsg);
               while (*++argv != NULL) (

/* look up the argument in the command table */

for (cmdp = cmds: cmdp->str != NULL; ++cmdp)

if (strcmp(*argv, cmdp->str) == 0)

break;
                               if (cmdp->str -- NULL)
     help("bad command");
                               else (
                                              /* excom must be installed to set options */
if (helpmsg -- remmsg && cmdp->fnct !- install)
    help(helpmsg);
                                              /* execute the command */
(*cmdp->fnct) (cmdp->arg);
               /* actually send the init bits to excom */
if ((init1 & COMINIT) != 0)
    int14(EXINIT, init1, 0);
               if ((init2 & COMINIT) != 0)
    int14(EXINIT, init2, 1);
               exit(0):
/* install excom */
install()
               system("excom");
helpmsg = insmsg;
/* remove excom */
void
remove()
               int14(5, 0, 0);
helpmsg = remmsg;
  * collect up the init bits for each port */
void
exinit (thebit)
               thebit:
               if (portnum -- -1)
                              help("no port selected");
              else if (portnum -- 0)
init1 |- thebit;
```

```
setport (newport)
                     portnum - newport;
                     if (portnum - 0)
                                           init1 |- COMINIT:
                     else
                                           init2 |- COMINIT:
/* perform int 14h with ah, al & dx as passed, return ax value */
int14(cmd, val, port)
                      val.
                     union REGS
                                                                                         /* registers send to bios */
/* registers returned from bios */
                     ir.x.dx = port;
                     int86(0x14, &ir, &or);
                      return or.x.ax;
     provide a bit of assistance */
help(msg)
                       *msg:
                      printf("\n%s\n\n", msg);
                      fputs("install\t\tinstall excom\n", stdout);
                      fputs("install\t\tinstall excom\n", stdout);
fputs("remove\t\tremove xcom\n", stdout);
fputs("coml\t\tsubsequent commands for coml\n", stdout);
fputs("com2\t\tsubsequent commands for com2\n", stdout);
fputs("dtr\t\tuse DTR for input flow control\n", stdout);
fputs("rts\t\tuse XRS for input flow control\n", stdout);
fputs("xnxfin\t\tuse XON/XOFF (^S ^Q) input flow control\n", stdout);
fputs("nodsr\t\cdon't require DSR to transmit\n", stdout);
fputs("nodsr\t\cdon't require DSR to transmit\n", stdout);
fputs("19200\t\tatket baud rate\n", stdout);
                      fputs("19200\t\tset baud rate\n", stdout);
fputs("38400\t\tset baud rate\n", stdout);
                       exit(1);
                                                                                                                                                                                                                                                                      End Listings
```

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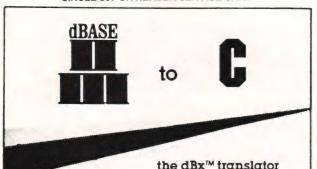
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TURBO PASCAL OVERLAYS

```
Listing One (Text begins on page 50.)
Listing 1: Memovrly.Inc
```

```
Turbo Pascal Memory Overlay Routines

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```

Limitations:

These routines have been tested only for Turbo 3.01A (both PC-DOS and generic MS-DOS). They may not work under 3.0 (the celebrated FileSize bug may cause trouble) and will certainly not work under 2.0 $\chi\chi$.

Memory overlay files must be < 64k in size!

NORMAL overlays nested inside memory overlays should work, but trying to nest memory overlays inside memory overlays would be disasterous!

OvrPath will not work in conjunction with memory overlays! (Writing a replacement routine would be simple if the code below makes sense to you.)

I/O testing in InitOverlay is just Turbo's Native. Anyone really needing memory overlays will probably wish to install their own I/O error checking.

CONST

RequiredHeap = \$1000; {Paragraphs of Heap Required by Program for other purposes than memory overlays. Change this to suit your needs for dynamic storage.}

TYPE

(Type used in both InitOverlay and DisposeOverlayStorage)
OverlayProcedure = RECORD
CASE Boolean OF

True:
(OldCall : ARRAY[1..3] OF Byte;
OldOffset : Integer;
FileName : ARRAY[1..13] OF Char;
);
False:
(NewCallInstruction : ARRAY[1..3]OF Byte;
NewCallAddress : Integer;
CurrentOffset : Integer;
OverlayCodeLoc : ^Byte;
NewRoutineLoc : Integer;
OverlaySize : Integer;

PROCEDURE NewOverlayHandler;

BEGIN INLINE (

(When this routine receives control, AX contains the number of bytes in the desired overlay & BX contains the offset (in pages) of the desired overlay within the overlay file (now on the heap).)

(First, check to see if the desired overlay is already in place by comparing DX with the offset recorded in memory immediately after the call instruction. If they match, no load is necessary)

\$46/ (INC SI) \$46/ (INC SI) \$2E/\$C5/\$34/ (LDS SI,CS:[SI])

(Multiply overlay page by 100H to get number of bytes code is displaced from start of overlay code area (on heap). Add to source offset in SI.} $88A/\$F2/ \qquad (MOV \qquad DH, DL \qquad) \\ \$32/\$D2/ \qquad (XOR \qquad DL, DL \qquad) \\ \$33/\$F2/ \qquad (ADD \qquad SI, DX \qquad)$

{Put number of bytes to move in CX} \$8B/\$C8/ (MOV CX,AX

{Copy CX bytes from DS:SI to ES:DI}

```
SEC /
                                                                (CLD
             $F3/$A4/
                                                                                    MOVSB
              (Recover mauled registers)
             $1F/
$5E/
                                                                                    DS
             {RUN_OVERLAY:}
$83/$C6/$0D/
                                                                                    SI, ODH
                                                                {ADD
             SFF/SE6
PROCEDURE InitOverlay(OverlayCallOffset : Integer);
         OverlayCallPtr : ^OverlayProcedure;
                                           : Integer;
: STRING[13];
: FILE;
         TestSize, i
     BEGIN
         OverlayCallPtr := Ptr(CSeg, OverlayCallOffset);
WITH OverlayCallPtr^ DO
             BEGIN
                   {Obtain overlay file name}
                   s := '';
WHILE FileName[i] <> #0 DO
                       BEGIN
                           s := s + FileName[i];
i := i + 1;
                       END:
                  Copen overlay file as untyped file}
Assign(f, s);
Reset(f);
                  Reset(f);

(determine file size in $80-byte sectors)

TestSize := FileSize(f);

(Check to see if there's enough space on the heap.)

(If there isn't, leave the overlay on disk)

IF (MemAvail > (RequiredHeap + TestSize * 8)) AND

(MaxAvail >= TestSize * 8) THEN (there's enough space)

BEGIN

OverlaySize := TestSize;

GetMem(OverlayCodeLoc, OverlaySize * $80);

BlockRead(f, OverlayCodeLoc^, OverlaySize, i);

NewCallInstruction[1] := $2E; (CS:)

NewCallInstruction[2] := $FF;

NewCallInstruction[3] := $16; {indirect near call}

NewCallAddress := Ofs(NewRoutineLoc);

NewRoutineLoc := Ofs(NewPowerlayHandler) + 7;

{extra 7 bytes skips turbo's procedure overhead}
                    {determine file size in $80-byte sectors}
                             (extra 7 bytes skips turbo's procedure overhead)
CurrentOffset := $FFFF; {force load on first call}
                       END;
                   Close(f);
              END;
     END:
 PROCEDURE DisposeOverlavStorage(OverlavCallOffset : Integer);
          OverlayCallPtr : ^OverlayProcedure;
     BEGIN
         SGIN
OverlayCallPtr := Ptr(CSeg, OverlayCallOffset);
WITH OverlayCallPtr^ DO
IF NewCallInstruction[3] = $16 THEN (Overlay is in memory)
FreeMem(OverlayCodeLoc, OverlaySize * $80);
```

End Listings

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CIRCLE 392 ON READER SERVICE CARD

UNIX BBS

```
Listing One (Text begins on page 54.)
1 TPATH=/u/bbs/rbin
   # check the command argument - $1
   if [ "$1" - MS] then
   # argument is for MS-DOS files
        echo 'who am i | cut -f1 -d" " ' 'date | cut -c1-16 ' "MS.list" >>/u/bbs/log.file
        more $TPATH/MS.list
   elif [ "$1" = Mac ]
    # argument is for Macintosh files
        echo 'who am i | cut -f1 -d" " ' 'date | cut -c1-16 ' "Mac.list" >>/u/bbs/log.file
        more $TPATH/Mac.list
12 elif [ "$1" - Unix ]
13 then
    # argument is for Unix files
        echo 'who am i | cut -fl -d" " ' 'date | cut -cl-cl6 ' "Unix.list" >>/u/bbs/log.
14
        more $TPATH/Unix.list
16 r
17 else
    # no argument was entered
        echo "List which directory? (MS, Mac, or Unix)"
                                                                        End Listing One
Listing Two
dwnld
    # identify which file directory (contained in the first argument - $1)
   if [ "$1" - MS ]
then
    # an MS-DOS file
    filedir=/u/bbs/MS-files
elif [ "$1" - Mac ]
    then
    # a Macintosh file
        filedir=/u/bbs/Mac-files
[ "$1" = Unix ]
    elif
    # a Unix file
        filedir=/u/bbs/Unix-files
    # no valid file directory was entered
        echo "Follow the dwnld command with a file directory - MS, Mac, or Unix"
    # verify that a file name has been entered (contained in the second argument - $2)
13 if [ -n "$2" ]
14 then
    # verify that the file name exists within the selected directory
        if [ -f "$filedir/$2" ]
            echo 'who am i | cut -f1 -d" " ' 'date | cut -c1-16 ' "dwnld" $1 $2
    # select a file transfer protocol
            echo "Transfer potocol (X = Xmodem; A = ASCII): \c"
            read method
if [ "$method" = X -o "$method" = x ]
    # send a binary file using the XModem protocol
```

echo "Prepare to transfer file. Press return to start."
read dummy

sleep 5
scho "Download complete. Press return to continue."

xmodem -sb \$filedir/\$2

cat \$filedir/\$2

send an ASCII file

22

```
read dummy
29
             fi
    # a valid file name wasn't entered
             echo "Sorry, that file name does not exist. Be sure to type the" echo "file name exactly as it appears in the directory listing." echo "Upper and lower case letters are different."
35
        fi
   else
    # no file name was entered
        echo "Please enter a file name after the directory on the command line." echo "Use the list command to see available file names."
                                                                              End Listing Two
Listing Three
    # verify that a file name has been entered (in argument 1 - $1)
    then
    # verify that a file with the same name isn't already stored in the Uploads directory
         if [ -f /u/bbs/Uploads/$1 ]
    else
    # collect the file transfer protocol
              echo 'who am i | cut -f1 -d" " ' 'date | cut -c1-16' "upld" $1 >>/u/bbs/log.
              echo "Transfer protocol (X = Xmodem; A = ASCII): \c"
             read method
if [ "$method" = X -o "$method" = X ]
10
     # receive a binary file using the XModem protocol
                  xmodem -rb /u/bbs/Uploads/$1
13
             6186
     # receive an ASCII file
                  echo "Begin sending ASCII file."
echo "Type a CNTRL-d when finished transmitting."
15
                  echo
cat > /u/bbs/Uploads/temp
cp /u/bbs/Uploads/temp /u/bbs/Uploads/$1
18
20
     # collect the file description
              echo "Please enter a one line description of your file."
              read desc
echo $1 $desc >>/u/bbs/Uploads/Doc.file
 23
     # no file name was entered
         echo "Please enter the name of the file you wish to upload on" echo "the command line."
                                                                              End Listing Three
 Listing Four
     FIRST=/u/bbs/.first
     LAST=/u/bbs/.last
MSG=/u/bbs/msg
echo 'who am i | cut -fl -d" " ' ' date | cut -cl-16' "readmsg" >>/u/bbs/log.file
     # display the number of the first and last messages available
     echo "Messages available:"
     cat $FIRST
echo "through"
cat $LAST
 10
     # accept the number of the first message to be displayed
 11 echo "Message number to read ('q' to quit): \c"
     # loop until the user wants to guit
      while [ "$message" !- q ]
     do
          echo
      # verify that the message number entered actually exists
                                                                     (continued on next page)
```

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CIRCLE 309 ON READER SERVICE CARD

UNIX BBS

```
Listing Four (Listing continued, text begins on page 54.)
```

```
if [ -f $MSG/$message ]
then

display the message

cat $MSG/$message
elif [ "$message" != q ]
then

the entered message number doesn't exist

echo "Sorry. That message doesn't exist."
fi
echo

accept the next message number to read

echo "Message number to read: ('q' to quit): \c"
read message
done
echo

End Listing Four
```

```
Listing Five
send
      echo 'who am i | cut -f1 -d" " ' 'date | cut -c1-16' "send" >>/u/bbs/log.file
SCAN=/u/bbs/msg/.index
     LAST=/11/bbs/ last
     MSG=/u/bbs/ms
L='cat $LAST'
D='date'
     # increment the number of the last message entered
    L='expr SL + 1'
     # save the new "last message" value
     echo $L >$LAST
     # collect message header information
     echo "To: \c"
     read to
echo "Subject: \c"
12
     read subject
echo "From: \c"
     read who
     trap 'rm -f $MSG/$L; continue' 2 3
     ech
     echo "Start typing your message. You have 20 lines available." echo "Type a period on a new line to end the message."
     * store the message header in the message file
     echo "# $L From: $who $D" >$MSG/$L
echo " To: $to" >>$MSG/$L
echo " Subject: $subject" >>$MSG/$L
echo >>$MSG/$L
     # collect the body of the message
25 NL-0
26
     read newline
     while [ "$newline" != "." -a "$NL" -1t 20 ]
           echo $newline >>$MSG/$L
NL='expr $NL + 1'
read newline
30
32 done
33
     # update the message index for use by the scan command
34 echo "Message being posted, please wait... \c"
35 echo "# $L From: $who $D" >>$MSG/temp
36 echo " To: $to" >>$MSG/temp
37 echo " Subject: $subject" >>$MSG/temp
38 echo " ">>$MSG/temp
39 cat $SCAN >>$MSG/temp
     mv $MSG/temp $SCAN echo
                                                                                                    End Listing Five
Listing Six
```

pmail

```
1 USERS=/u/bbs/rbin/user.file
2 echo 'who am i | cut -fi -d" " 'date | cut -cl-16' "pmail" >>/u/bbs/log.file

# display pmail instructions

3 echo
4 echo echo Commands:
6 echo S - send mail to another user

(continued on page 84)
```

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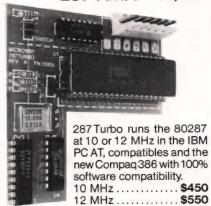
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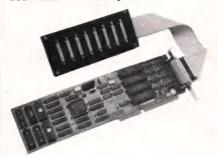
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UNIX BBS

Listing Six (Listing continued, text begins on page 54.)

```
r - read mail sent to you
l - list the users by user id and name
x - exit the private mail system
 11 echo
       # accept the first pmail command
 12 echo "Command: \c"
13 read choice
       # loop until the user wants to quit
     while [ "$choice" != x -o "$choice" != X ]
 16
            case $choice in [Ss])
      send private mail
print instructions
19
20
21
                        echo "WARNING: To successfully send mail observe these rules:"
                        echo
                        echo "
                        echo " 1. If you wish to erase a character, use the backspace key."
echo "DO NOT, DO NOT use the delete key."
echo " 2. Type a carriage return at the end of each line on the"
echo "screen. Lines should be no more than 80 characters."
echo " 3. Do not use the following characters in your message (they"
echo "have special meaning to the system): @ and #"
22
24
25
26
27
                        echo
                        echo "Press RETURN to begin sending your message: \c"
                        read dummy
      # collect user id of recipient
31
                        echo "To whom (user id): \c"
                        read to
      # verify that recipient entered is a valid user id
                        TPERS='grep $to $USERS | wc -c' if [ "$TPERS" -eq 0 ] then
33
      user id is invalid
                             echo "Sorry, that user id doesn't exist." echo "List them with the l command."
38
                        else
      # send mail
40
41
42
43
44
45
                             echo "Start typing your message. Type a period on a new" echo "line to send the message."
                             mail $to
                        Fi
      # collect next option
                        echo "Command (x to exit): \c"
48
                        read choice
                        continue;;
                 (Rrl)
      # read private mail
# print instructions
                        echo "After each message you will receive a '?' prompt." echo "Type 'd' to delete the message, <cr> to leave it in" echo "your mailbox, or 'q' to stop reading mail."
51
52
53
      # real the mail
                        echo
      # collect the next command
57
                       echo "Command (x to exit): \c" read choice
                continue;;
     # display a list of user names and associated user id's
61
                       echo
                       more $USERS
echo
62
     # collect the next command
                       echo "Command (x to exit): \c"
65
                       read choice
     # if an unrecognized command was entered, prompt for another one
                       echo "Command (x to exit): \c"
69
70
                       read choice
                       continue;;
                                                                                                        End Listing Six
                                                                         (Listing Seven begins on page 86.)
     echo
```

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write the new account data to a holding file

CIRCLE 336 ON READER SERVICE CARD

```
Listing Seven (Listing continued, text begins on page 54.)
 .profile from the Account new
     # set default path to an empty rbin directory
    PATH=/u/bbs/new/rbin
    # set level one prompt to remind users how to log off
2 PS1='Type CNTRL-D now... '
    # display the welcome message
    cat msgl
    # make sure the user really wants an account
    echo "Do you wish to request a login? (Y/N) \c"
    # continue only if the user answers yes
    if [ "$choice" - y -o "$choice" - Y ]
        echo
10
    # loop until user is satisfied with data
        while [ "$OK" !- Y -a "$OK" !- y ]
    # collect user name, address, etc.
13
            echo "Enter your real name: \c"
14
            echo
16
17
18
            echo "Enter the first line of your mailing address: (4 lines are available)"
            read address1
            echo
            echo "Enter the second line of your mailing address: <cr> if none" read address2
19
20
            echo "Enter the third line of your mailing address: <cr> if none" read address3
21
22
23
            echo
25
            echo "Enter the fourth line of your mailing address: <cr> if none"
            read address4
            echo
28
            echo "Enter your voice telephone number: \c"
            read phone
30
            echo
    display the data entered for the user
31
            echo "This is what you have entered:"
32
33
            echo $realname
34
            echo
            echo $address1
            echo $address2
37
            echo Saddress3
38
            echo $address4
            echo
40
            echo $phone
    # ask the user to verify the data
            echo "Is this correct? (Y/N) \c"
43
            read OK
       done
OK="N"
45
    # loop until user is happy
        while [ "$OK" !- Y -a "$OK" !- Y ]
    # collect the account data
            echo
49
50
            echo "Enter a one word login name: \c"
            read lognmae
51
            echo
52
53
54
            echo "Enter an initial password: \c"
            read initpasswd
    # display the entered data for the user
            echo "You have entered."
56
            echo
            echo $logname
echo $initpasswd
57
            echo
   # ask the user to verify the data
60
           echo "Is this correct? (Y/N) \c"
           read OK
61
       done
```

(continued on page 88)

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Listing Seven (Listing continued, text begins on page 54.)

```
echo $realname >>signups
echo $address1 >>signups
echo $address2 >>signups
echo $address3 >>signups
64
65
66
67
68
69
70
71
              echo $address4 >>signups
echo $phone >>signups
echo $logname >>signups
              echo $initpasswd >>signups
echo >>signups
       # tell the user that he or she is finished
              echo "This completes the login request process. You'll receive an account" echo "confirmation through the mail within a week."
74
       # tell the user to log off
77 echo
78 echo "Enter CTNRL-D to log off the system."
79 echo
```

Listing Eight

```
.profile for the info Account
```

- # set the default path to an empty rbin directory
- PATH=/u/bbs/info/rbin
 - # set the level one prompt to the logoff message
- PS1='Type CNTRL-D now...'
 - # display the welcome message

```
echo
    echo "
                           Welcome to Scholastech Telecommunications"
Scholastech Info"
    echo "
    echo
    cat msg
    echo
    echo "Press the carriage return to begin: \c"
10 read dummy
```

- # display the contents of the information file
- more info.file echo "Press the carriage return when done: \c" read dummy
 - # give the user the choice of whether or not to sign up for workshops
- 16 echo "Do you wish to sign up for a workshop? (Y/N) \c" 17 read dummy
 - # loop until user doesn't want to register for any more workshops

```
while [ "$dummy" - y -o "$dummy" - Y ]
          echo
20
          echo "Enter your name: \c"
          read regname
23
          echo "For which workshop are you registering? \c"
          read workshop
26
27
28
          echo
          echo
          echo "Enter your school or company name: \c" read school
          echo "For how many people are you registering? \c" read people
32
          echo "Enter a voice phone number where you can be reached in case" echo "there are any questions about your registration: \c" read phone
34
35
```

write the registration data to disk

```
echo $regname >>signups
              echo $workshop >>signups
echo $school >>signups
echo $people >>signups
39
40
              echo $phone >>signups
echo >>signups
```

allow the user to register for more than one workshop

44 45 echo "Do you wish to register for another workshop? (Y/N) \c" read dummy 46 done

tell the user how to log off

echo "Type CNTRL-D to log off the system."

End Listing Eight

End Listing Seven

(Listing Nine begins on page 90.)

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```
Listing Nine (Listing continued, text begins on page 54.)
     # create a temporary file with a log on record for each remote user
     who /etc/wtmp | grep ph >/u/user.log
     # if an "a" was entered as an argument to the command, display the entire log files
     if [ "$1" - a ]
         echo
         echo "System Logins"
     # display all system use since /etc/wtmp was last purged
         more /u/user.log
          echo "Next?"
 10
         read dummy
     # display all BBS use since /u/bbs/log.file was last purged
         echo "BBS User Log"
         more /u/bbs/log.file
     * otherwise, show only the last ten entries in each log file
         echo "Last System Logins"
        tail /u/user.log
 22
         echo "Last BBS Activity"
         tail /u/bbs/log.file
        echo
     # remove the temporary file
                                                                      End Listing Nine
 30 rm /u/user.log
 Listing Ten
 usage
     # print headings to report file
                                         Logout" >usage.temp
    echo "User
                    Date
     echo "----echo " " >>usage.temp
     # get a log in record from /etc/wtmp for remote users on device phl
    who /etc/wtmp | grep phl >templ
     # get the user name
    cut -f1 -d" " temp1 >temp2
    get the log in date and time
   cut -c25-36 temp1 >temp3
    # get log out record from /etc/wtmp and extract the log out time
    who -d /etc/wtmp | grep phl | grep -v LOGIN | grep -v rc | cut -c32-36 >temp4
    # paste together for the report
    paste temp2 temp3 temp4 >>usage.temp
    # print the report
9 pr usage.temp >usage.summary
    # remove temporary files
10 rm templ
11 rm temp2
12 rm temp3
    # format the BBS command use log for output and purge the log file
15 pr /u/bbs/log.file >log.summary
16 > /u/bbs/log.file
    d clean out /etc/wtmp
17 > /etc/wtmp
    # send both reports to the printer
18 lp usage.summary
19 lp log.summary
```

End Listings



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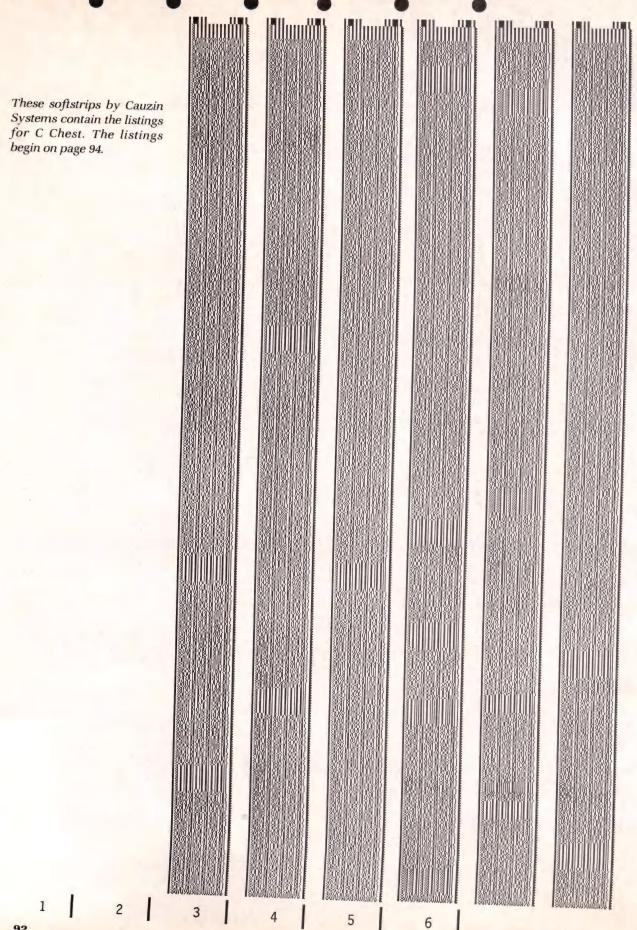
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C CHEST

Listing One (Text begins on page 102.)

```
11 finclude <stdio.h>
 31 /* PQ.C
                      General-purpose priority-queue routines.
                      (C) 1987 Allen I. Holub, All Rights Reserved.
 41
 51
        typedef char *PO:
                               Dummy typedef for priority queue.
     * PQ *pq_create( numele, elesize, cmp, swap, initheap)
* int numele; Max # of elements in the queue
 81
 91
       int elesize;
int (*cmp)()
101
                               Size of one element in bytes
111
              (* cmp) ();
                               Pointer to comparison function
       int (*swap) () ;
                               Pointer to swap function
121
        char *initheap;
                               Inital heap or NULL to allocate
141
151
      * pq_ins(p, item)
                               Insert Item into queue
     * PQ *p;
* char *item;
                               Pointer to priority queue
Pointer to item to insert
171
181
                               Return number of empty slots
                               before insertion (0 if none).
201
     • int pq_del(p, target) Delete item from queue
211
221
        PO
              *p:
                                   Pointer to priority queue
     * char *target;
231
                                    Pointer to place to put deleted
241
                                    item.
251
                                    Return # of items in queue before
261
                                   delete (0 if nothing deleted) .
271
     * char *pq_look( queue )
* PQ *queue;
281
                                   Look at (don't delete) top element
                                   Pointer to queue.
301
311
     " int pq_numele( queue ) Return # of elements in queue.
" PQ "queue; Pointer to queue.
331
341
361
371
    typedef struct
381 (
391
                               /* compares two objects
                (*cmo) ():
401
                               /* swap two objects
         int
                (*swap) ();
              itemsize;
411
         int
                               /* size of one element in heap
421
         int
                               /* Number of items in the heap
               nitems:
431
               maxitem;
                               /* Maximum number of items in heap
441
         char *bottom;
                               /* Ptr. to most-recently added item
         char *heap;
451
                               /* pointer to the heap itself
461
471
    PO:
481
501
51| static void reheap_down(p)
    PO
             *p;
531 (
541
         /* Reheap the Heap, starting at the top and working
          a down;
561
         +/
571
         int
                               /* index of parent
/* index of child
                parent;
                child:
591
         int
601
         char
                *poarent:
                               /* pointer to parent
61 1
         char
                               /* pointer to child
                *pchild;
                               /* pointer to child's sibling
621
         char
                *psibling:
631
        char *heap;
                               /* pointer to heap
641
651
        heap = p->heap;
661
671
        for ( parent = 0, child = 1; child < p->nitems ;)
691
             pparent = heap + (parent * p->itemsize);
701
             pchild = heap + (child * p->itemsize);
721
             if ( child+1 < p->nitems )
731
                 psibling = pchild + p->itemsize :
751
                 if ( (*p->cmp) (pchild, psibling) < 0 )
771
781
                      pchild = psibling;
791
                      ++child:
801
821
831
             if( (*p->cmp) ( pparent, pchild ) >= 0)
851
861
             (*p->swap) ( pparent, pchild );
             parent = child;
child = (parent = 2) + 1;
881
911
```

```
931 /=
 941
 95| static void
                       reheap up( p )
 961 PQ
 971 (
          /* Reheap the Heap, starting at the bottom and working up.
 981
           " Note that we must use a divide-by-2 rather than a
 991
           * shift-right in the while loop because -1/2 -- 0 but
1001
101
             -1 >> 1 == -1.
1021
1031
1041
                                /* index of parent
         1nt
                 parent:
105
                  child;
                                /* index of child
                                /* pointer to parent
/* pointer to child
1061
          char
                 *pparent;
1071
                  *pchild:
          char
108
                 *heap;
                                /* pointer to heap
1091
1101
          child = p->nitems - 1:
         heap = p->heap;
1111
1121
1131
          while ( (parent = (child-1) / 2) >= 0 )
1141
              pchild = heap + (child * p->itemsize);
1151
              pparent = heap + (parent * p->itemsize);
1161
1171
1181
              if ( (*p->cmp) ( pparent, pchild ) >= 0)
119
1201
1211
               (*p->swap) ( poarent, pchild ):
1221
               child - parent;
1231
124| }
1251
1261 /
1281 PQ
              *pq_create( numele, elesize, cmp, swap, initheap )
1291
1301 Int
              numele;
                                /* max # of elements in the queue */
                                /* size of one element in byte */
/* pointer to comparison function */
131| int
               elesize;
1321 int
               (*cmp) ();
                                /* pointer to swap function */
/* inital heap, NULL to allocate */
               (*swap) ();
1331 int
     char
1351 (
          /* Create a priority queue that can hold at most * "numele" elements, each of size "elesize". The
1361
1381
           * cmp function is passed two pointers to queue
            elements and it should behave as follows:
1391
140
1411
                   (*cmp) (pl. p2)
1421
           * For descending priority queues (pq_get() returns the
143
1441
                                                       largest iteml.
1451
                   *p1 < *p2
*p1 -- *p2
                                     return < 0
1461
                                     return -- 0
1471
                                     return > 0
1481
           1491
150
151
                                     return > 0
                   *p1 -- *p2
1521
                                     return -- 0
                   *pl > *p2
                                     return < 0
1531
154
           " If the initheap argument is NULL, an empty heap is
155

    created automatically, otherwise initheap must point
    at an initialized numele-element-long heap.

1561
1571
1581
1591
1601
          PO
                   *p, *malloc() ;
161 |
                heapsize;
1621
1631
          heapsize = numele * elesize ; /* heap size in bytes */
164
1651
          if ( initheap )
1661
167
               if( !(p = malloc(sizeof(PQ))) )
1681
                   return 0;
1691
1701
              p->heap = initheap;
1711
              p->bottom = initheap + ((numele - 1) * elesize);
1721
              p->nitems = numele;
1731
1741
1751
1761
              if( !(p = malloc(sizeof(PQ) + heapsize)) )
1771
                   return 0:
1781
1791
              p->heap = (char *) (p + 1);
p->nitems = 0;
1801
1811
              p->bottom = p->heap - elesize :
1821
1831
1841
         p->cmp
                       - cmp;
```

```
p->swap
                       - swap;
          p->itemsize = elesize;
186
1871
          p->maxitem = numele ;
1881
          return n:
1891 }
1901
1911 /
1921
193| pq ins(p, item)
194| PQ *p;
                                  /* Pointer to priority queue
                                 /* Pointer to item to insert
1951
      char
1961
      1
              Insert a new item into the priority queue (provided
              that space is available.
1981
1991
2001
              Return the number of empty slots in the queue before
              the insertion. This number is 0 if the queue is full and nothing is inserted. Algorithm is:
201
2021
 2031
204
               if ( space is available in the queue )
                        increase queue size copy new item into bottom of queue
2051
 206
 207
                         reheap from the bottom up.
 2081
 2091
 2101
           int space avail = p->maxitem - p->nitems;
 2111
           if ( space_avail > 0 )
 2121
 2131
 2141
                ++ ( p->nitems );
                   cpy( p->bottom += p->itemsize, &item, p->itemsize );
 2161
                reheap up ( p );
 2171
 2181
 2191
           return space_avail ;
 2201 }
 2211
 2221 /
 2231
 2241 int
                pq_del(p, target)
 2251
 2261 PQ
                                /* pointer to priority queue
                                /* place to copy current largest item *
 2271
      char
                *target;
 2281 (
                Copy the largest item in the priority queue to the address held in target, then delete the item.
 2291
 230
 2311
                Return the number of items in the queue before the
 2321
                delete. If this number is 0, then nothing was in the queue and *target will not have been
 2331
 2341
                modified. Algorithm is:
 2351
 2361
                                                                     [1]
                 if ( there's something in the queue )
                      remember pointer to former first item
 238
                      replace the first item with the last one
 2391
                      shrink the heap by one element
 2401
                                                                     [5]
                      reheap from the top down
 2411
 2421
              ./
 2431
             int
                      slots_in_use;
 245
                                                                    /* 1 */
             if ( slots_in_use = p->nitems )
 2461
 247
                 memcpy( target, p->heap, p->itemsize );
memcpy( p->heap, p->bottom, p->itemsize );
 248
                                                                    1 3 1/
 2491
 2501
                  -- (p->nitems) ;
                  p->bottom -= p->itemsize;
 2521
 2531
                  reheap down ( p );
  2541
  255
  2561
  2571
             return slots in use ;
  2581 1
  2591
  2601 /
  2611
        char
                 *pq_look ( queue )
  2621
  2631 PQ
                 *queue;
  2641
        -
                 /* Return a pointer to the largest/smallest element
  2651
                  * in the priority queue but don't dequeue it.
  2661
  2671
  268
                 return queue->heap;
   2691
  2701
   2711
   2721
   2731
                  pq_numele ( queue )
   2741 int
   2751
        PQ
                  /* Return number of items in queue
   2761
        1
   2771
   278
   2791
                  return queue->nitems;
   2811 1
                                            (continued on next page)
   2821
```

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C CHEST

Listing One (Listing continued, text begins on page 102.)

```
2841 #ifdef MAIN
2851
286| int Descending = 1;
                                /* Change these in Codeview
287| int Makequeue = 1;
                               /* to change the tests.
288
289| cmp(s1, s2)
290| char **s1, **s2;
2911 (
2921
              int rval:
293
              rval = strcmp( *s1, *s2 );
2941
2951
              if ( Descending )
2961
                  return rval:
2971
              else
                  return (rval < 0) ? 1:
(rval > 0) ? -1:
/* rval == 0 */ 0;
2981
2991
3001
301| }
3021
3041
305| swap( s1, s2 )
              **s1, **s2;
3061 char
3071 1
              char
3081
                      *tmp;
3091
              tmp = *s1;
*s1 = *s2;
3101
3111
              *s2 = tmp;
3121
313| |
3141
3151 /* - - - - - - - - - - - -
3161
3171 printq(p)
318| PQ
               *p;
3191 (
3201
          int 1:
3211
3221
          printf("Queue is:\n"):
3231
         3241
3251
3261
                                                    p->itemsize );
          printf( "\tnitems..... %d\n",
3271
                                                    p->nitems
         printf( "\tmaxitem...... %d\t", p->maxitem
printf( "\tbottom...... 0x%04x\n", p->bottom
3281
                                                    p->maxitem );
3291
         printf("\tbottom - heap... %d\n\n", (char **)p->bottom -
330
3311
3321
                                                    (char **)p->heap );
3331
          if ( p->nitems <= 0 )
3341
              printf("\tqueue is empty\n");
3361
          else for ( i = 0 ; i < p->nitems ; i++ )
3371
338
            printf("\t%-2d: %10.10s (0x%04x) (children: %2d, %2d)\n",
3391
                   i, ((char **)p->heap)[i], ((char **)p->heap)[i], (2*i)+1, (2*i)+2
3401
3411
3421 1
3441
3451
3461
    main()
3481
               *queue;
3491
          char buf[80];
3501
                12
351
3521
353
          static char *testq[] =
354
3551
              "0", "1", "2", "3", "4", "5", "6", "7", "8", "9"
3561
3571
358
3591
              queue = pq_create( 10, sizeof(char*), cmp, swap, 0 );
360
361
              queue - pq_create( 10, sizeof(char*), cmp, swap, testq);
3621
3631
          if ( Iqueue )
3641
              printf("pq_create failed\n");
3661
3671
368
3691
                                                                  +\n*);
3701
         printf("| Enter i<string><CR> to add string to
printf("| queue, d<CR> to dequeue top element, q to
printf("| exit the program.
                                                                  (\n");
371
3721
                                                                  1\n*):
         printf("+--
                                                                  +\n");
374
         while( 1 )
376
```

```
printq( queue
printf( "\n[i|d|q][string) :"
3781
              gets ( buf
3791
3801
381
              if( *buf == 'q' )
3821
                     exit(1);
3831
3841
              else if( *buf -- '1' )
3851
3861
                  i = pq_ins( queue, strsave(buf + 1) );
3871
3881
3891
                    printf("%d slots avail. before insert\n", i);
3901
3911
                   printf("Queue was full, did nothing\n\n");
3921
3931
              else
3941
3951
                  i = pq del ( queue, (char*) &p );
3961
                 printf("%d slots used before delete, got <%s>\n",
3971
3981
3991
4001
                      free ( p );
                      p = "nothing";
4011
4021
4031
4041
        1
4051 )
4071 #endif
                                                    End Listing One
```

Listing Two

Listing 2 -- strsave.c

```
11 char
             *strsave( str )
 21 char
         /* Save the indicated string in a malloc()ed section

    of static memory. Return a pointer to the copy or
    0 if malloc failed.

71
81
        register char *rptr;
        extern char *malloc();
101
111
121
        if ( rotr = malloc( strlen(str) +1 ))
141
                 stropy( rptr, str );
                 return rptr;
161
171
181
        return (char *)0;
191 1
                                                    End Listing Two
```

Listing Three

Listing 3 - freq.c

```
#include <stdio.h>
 31
              FREO.C
                                  Print a list of the frequency
                                  of occurance of all bytes
               in a list of files given on the command line
 61
              Frequencies are printed as a probability x 100.
 71
              For example, if we read a total of 20 characters,
              5 of which are 'e', the probability of an 'e' occuring in the input is 5/20 (.25) and freq will
 91
101
              output (.25 * 100) or 25.
111
121
13| typedef struct
141
151
              int
                       val:
              double count;
17| | ITEM;
191
    #define TABSIZE 256
                                    /* I'm counting on this being
20| ITEM Tab[ TABSIZE ];
21 |
                                    /* initialized to zero.
221
241
25| cmp(iteml, item2)
              *iteml, *item2;
271
         /* Comparison function used by ssort(), below.

- Count is the primary sort field and val is

- the secondary field.
281
291
301
31 |
321
         int rval:
```

```
351
         return
                   ( item1->count < item2->count )
                    ( item1->count > item2->count )
361
                  /* iteml->count == item2->count */
371
                                            item1->val - item2->val ;
381
40
411 /
421
43| main( argc, argv )
              **argv;
451 1
                         *bin to ascii();
                                                   /* in pchar.c */
461
         char
471
         FILE
                         *fp;
         int
                         1;
                         smallest:
49
         double
         double
                          largest;
501
51|
         double
                          numchars
                                       - 0.0 ;
52
         double
                          sum
                                       - 0.0 2
                         probability = 0.0 ;
531
         double
541
                                          /* Needed only for
55
         reargy ( &argc, &argv );
                                          /* On Command! shell
561
571
         for ( -- argc, ++ argv ; -- argc >= 0 ; argv++ )
591
              if( |(fp = fopen( *argv, "rb")) )
601
                  perror( *argv );
611
 621
 631
                   fprintf( stderr, "%s\n", *argv );
 641
 651
                   while ( (i - getc(fp)) !- EOF )
 661
 68
                        ++ numchars:
                        ++ Tab[ i & Oxff ].count ;
 691
 701
 711
                   fclose (fp);
 721
 731
 741
          /* Find largest and smallest elements and at the mame
 76
           * time initialize the val fields of the Tab entries.
 77
 781
 791
 801
          largest = 0.0;
          smallest = numchars :
 82
          for ( 1 = 0 ; 1 <- Oxff ; 1++ )
 831
              Tabiil.val = is
 85
 86
               if( Tab[i].count > 0.00000001
              44 Tab(i).count < smallest )
smallest = Tab(i).count;</pre>
 88
 891
 90
 911
               if ( Tab[i].count > largest )
 92
                        largest - Tab(i).count;
 931
 941
          /* Sort the list by probability. A shell sort is used.
    You can replace the ssort call below with a call to
 961
             the quort () subroutine, available with many compilers.
 971
           * Qsort() takes the same arguments as ssort().
 981
           * A version of quort appeared in the in the C Chest,
 991
           DDJ #102 (April, 1985; also Bound Volume 10, p.316).
The ssort() subroutine appeared in DDJ #113, C Chest
1001
101
102
            * (March, 1986) p. 70.
1031
104
           ssort ( Tab, TABSIZE, sizeof (ITEM), cmp );
105
1061
1071
108
            * Print the list. Each element is printed as three
            numbers: the value of the charcter (in hex), the
probability of that character apearing in the input,
109|
1101
              and the probabability normalized so that the least-
1111
1121
            * frequently occuring probability has the value 1.
1131
1141
           for ( i = 0 ; i < TABSIZE ; i++ )
115
 116
                    probability = Tab[i].count / numchars;
sum += probability;
1171
1181
 1191
                    printf( "0x402x\t47.6f\t41.0f\n",
 1201
 1211
                                  Tabfil.val.
 1221
                                  Tab[i].count / smallest );
 123
 1241
 125
           fprintf( stderr, "Total = %8.6f (N = %8.2f)\n",
 126
                                                     sum, numchars );
 1271
 1281
```

End Listings

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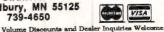
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ARTIFICIAL INTELLIGENCE

Listing One (Text begins on page 116.) Listing 1. Inheritance in SCOOPS ; (C) Copyright 1987 Ernest R. Tello (define-class artifact (instvars material weight purpose cost) (options (gettable-variables material weight purpose cost) settable-variables inittable-variables)) (define-class transport-means (instvars medium time-range power-source) (mixins artifact) (options (gettable-variables medium time-range power-source) settable-variables inittable-variables)) (define-class transport-vehicle (instvars load-capacity length max-speed) (mixins artifact transport-means) (options (gettable-variables load-capacity length max-speed) settable-variables inittable-variables)) (define-class passenger-vehicle (instvars capacity safety dining facilities) (mixins artifact transport-means transport-vehicle) (gettable-variables capacity safety dining facilities) settable-variables inittable-variables)) (define-class water-transport-vehicle (classvars (body-name 'hull) (dof 2) (dangers 'sink) (advantages 'relaxing)) (mixins artifact transport-means transport-vehicle passenger-vehicle) (gettable-variables dof dangers) settable-variables inittable-variables)) (define-class surface-vessel (instvars #-decks #-masts #-engines) (mixins artifact transport-means transport-vehicle passenger-vehicle water-transportvehicle) (gettable-variables #-decks #-masts #-engines) settable-variables inittable-variables)) (define-class ship (instvars x-position y-position x-velocity y-velocity mass) (mixins surface-vessel) (options (gettable-variables x-position y-position x-velocity y-velocity mass) settable-variables inittable-variables)) (define-method (ship speed) () (sqrt (+ (expt x-velocity 2) (expt y-velocity 2)))) (define-method (ship direction) () (atan y-velocity x-velocity)) (define-class ocean-liner (instvars company launched homeport tons) (mixins ship) (options (gettable-variables company launched homeport tons) settable-variables inittable-variables)) (define ship) (make-instance ship 'x-position 100 'y-position 150 'x-velocity 30 'y-velocity 40 'mass 100)) (compile-class artifact) (compile-class transport-means) (compile-class transport-vehicle) (compile-class passenger-vehicle) (compile-class water-transport-vehicle) (compile-class surface-vessel)

(compile-class ship)

(compile-class ocean-liner)

End Listing One

Listing Two

Listing 2. Multiple Inheritance in SCOOPS

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(define-class business

(instvars name location industry business-type size year-founded ownership-type gross-sales costs

market-share)

(options

(gettable-variables name location industry business-type size year-founded ownership-type

gross-sales costs market-share) settable-variables

(define-method (business calc-net-gain) (gross-sales costs) (- gross-sales costs))

(define-class adversary

(instvars aggressiveness allies goals common-goals strengths weaknesses)

(options

(gettable-variables aggressiveness allies goals common-goals strengths weaknesses)

settable-variables inittable-variables))

(define-class competitor

(mixins business adversary))

(compile-class business) (compile-class adversary) (compile-class competitor)

(define your-business (make-instance business))

(define competitor-1 (make-instance competitor)) End Listings

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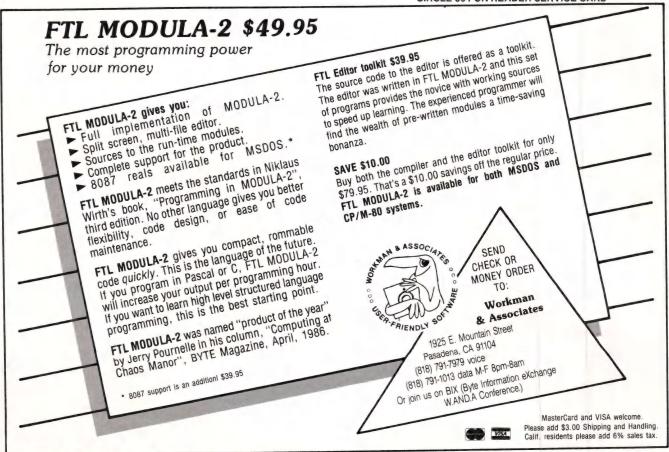


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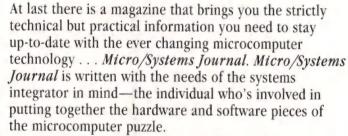
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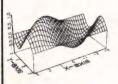
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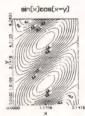
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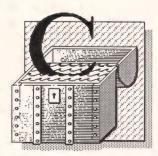
his month is the first of what has unintentionally become a two-part article. I had intended to implement an interesting variant of the Huffman encoding data-compression algorithm that is used by the Unix COMPACT program and is described in Robert Gallager's article "Variations on a Theme of Huffman" (IEEE Transactions on Information Theory, vol. IT-24, no. 6 [November 1978]: 668-674). Gallager describes an adaptive one-pass Huffman code in which the code changes as the input file is processed. This way the code tree stays optimal over the life of the file. As usual, the problem proved more intractable than I had at first anticipated. So, this month I'll describe some of the stuff I developed along the way to a solution: I'll discuss the actual compression algorithm in a future column.

Priority Queues

A queue, in the normal data-structure sense of the word, works like a line in a bank does. New entries are added to the back of the queue, and items are removed from the front. That is, the items in the queue are ordered by insertion time—those items

by Allen Holub

that were inserted first are removed first. Another kind of queue is a priority queue or heap, a queuelike data structure in which something other than time is used to order the elements. For example, the largest or smallest element—rather than the least-recently inserted element—could be the first to be dequeued. Heaps have several uses (other than Huffman codes). For example, a heap is used to do the merge phase of the



external-sorting program described in the June 1986 C Chest. A clever sorting algorithm (heapsort) uses a priority queue to do the sorting. Elements are put into the queue in random order, and the smallest element is extracted repetitively until the array is sorted.

The file PQ.C (Listing One, page 94) contains a set of general-purpose priority queue routines. Queues can be constructed of any sort of object (numbers, pointers, structures, and so forth) and can be ordered in any sort of way. You could, for example. keep a queue of structures, adding structures to the queue at random but always extracting the one with the smallest key. You could also create a queue of pointers to structures. You could even use these routines to maintain a normal (but inefficient) queue by adding a time-entered field to the structure and then extracting by smallest time entered. Similarly, a stack can be represented by a priority queue in which the item with the largest time entered is removed first.

PQ.C contains five externally accessible routines. The first of these creates a new queue:

typedef char QUEUE;
QUEUE *pq_create(numele, elesize, cmp, swap, initheap)

int numele; int elesize; int (*cmp)();
int (*swap)();
void *initheap;

The QUEUE type is a dummy pointersize type that is used in the same way as is the FILE pointer returned from fopen(). Pq_create() is passed five arguments. Numele is the maximum number of elements in the queue. Elesize is the size of one element. Cmp is a pointer to a comparison function. It is passed pointers to two enqueued objects and should return a value as indicated in Table 1, page 105. Swap is a pointer to a swap function. It is passed pointers to two objects, and it should swap the objects (not the pointers). Finally, initheap can be used in two ways: if it is NULL, an empty priority queue is created; otherwise, initheap is assumed to be a pointer to an already-initialized numele-long array that will be used as the heap. I'll discuss the utility of this in a moment. I've used the ANSI syntax here to declare initheap. A void pointer is one that doesn't point at any explicit object; you have to cast it to a pointer to a real type to use it. If your compiler doesn't support void pointers, use a pointer to char.

A queue that's created by pq_create() can be deleted by free(); just pass it the pointer returned from pq_create(). Note, however, that this will only free memory allocated by pq_create() itself. If you create your own initial queue and pass it to pq_create() via the initqueue argument, you'll have to free up the memory that you allocated. By the same token, if you create a queue of pointers to strings, free() will free the memory used by the queue but not by the strings.

I'll illustrate how to set up a queue

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C CHEST

(continued from page 102)

with a moderately complex example-say you want to keep a queue of pointers to the following structures:

```
typedef struct
  int weight;
  char *stuff:
  long more_stuff;
ITEM;
```

The ITEMs can be inserted in any order, but they will be extracted in order of decreasing weight. The comparison function used for this purpose looks like this:

```
cmp(p1, p2)
ITEM **p1, **p2;
  return (*p1)->weight
    - (*p2)->weight;
```

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and the swap function looks like this:

```
swap(p1, p2)
ITEM **p1, **p2;
 ITEM *tmp;
 tmp = *p1;
  *p1 = *p2;
  p2 = tmp;
```

If you had wanted to order the queue by increasing-rather than decreasing-weight, you would have reversed p1 and p2 in the comparison function's return statement.

An empty ten-element queue can now be created with the following:

```
QUEUE *qp;
qp = pq_create( 10, sizeof(ITEM*),
                    cmp, swap, 0);
```

Items are added to the queue using:

```
int pq_ins(qp, item)
QUEUE *qp;
char *item;
```

where qp is a pointer returned from a previous pq_create() call and item is a pointer to the object to insert. Pq_ins() returns the number of empty slots that were in the queue

before the insert. If the return value is 0, the queue was full and $pq_ins()$ will have done nothing. A new item is inserted in the queue that I created earlier with the following code:

ITEM *p;
p = (ITEM *) malloc (sizeof(ITEM));
p->stuff = "A string";
p->weight = 5;
if(!pq_ins(qp, &p))
 printf("queue is full");

Note that you have to pass in the address of the object to enqueue—in this case the address of the pointer *p*. Items are extracted from the queue with:

int pq_del(qp, item) QUEUE *qp; char *item;

where, again, qp is a <code>QUEUE</code> pointer returned from $pq_create()$ and <code>item</code> is a pointer to a place into which the dequeued object will be copied. The number of items in the queue before the delete is returned. If this number is 0, the queue was empty and the contents of *item* are undefined. For example, the largest element of the queue can be dequeued with the following code:

ITEM *p;
if(!pq_del(qp, &p))
printf("queue is empty");

Two additional support routines are provided:

char *pq_look(p) char *pq_numele(p) QUEUE *p;

Again, *p* is a pointer to a *QUEUE* returned from a previous *pq_create()* call. *Pq_look()* returns a pointer to the object at the head of the queue. The object is not actually dequeued, however. *Pq_numele()* returns the number of elements currently in the queue.

Implementation

A priority queue can be represented by a binary tree that has the following properties:

1. All children in the tree have a value less than their parent.

- 2. The tree is as perfectly balanced as possible—that is, the difference in height of all leaves is at most 1.
- 3. Leaves are inserted into the tree from left to right until an entire rank is full, then the next rank is started.

An example tree is shown in Figure 1, below. Note that this tree is not strictly ordered in the normal way. That is, rule 1 doesn't require that the tree be sorted, only that the children are smaller than the parent. You could exchange the subtrees rooted at 1 and 2 without violating rule 1. This ordering guarantees that the root node of the entire tree always holds the largest element, however. Notice that an insert operation is harder than normal because of rules 2 and 3. An 11th node must be inserted as the right child of node 4 and a 12th node as the left child of node 5. If you do this, however, rule 1 may be violated by the newly inserted node. To avoid this last problem, you have to adjust the contents of the nodes at every insertion (a process called reheaping). For example, if you insert a new node having the value *k* as the right child of node 4, you'll have to shuffle things around because of rule 1.

The reheaping process is illustrated in Figure 2, page 106. You reheap from the bottom up with an insertion. Because k is greater than h, the contents of nodes 10 and 4 must be swapped. You then go up one level to node 4 and compare its key to its parent's. Here k is still greater than i, so you swap again, moving the k to node 1. Finally, nodes 0 and 1 must be swapped as well, moving k to the root position.

Because of rules 2 and 3, it's convenient to represent the tree as an array in which the element at node N has children at nodes 2N+1 and 2N+2. For example, the node at array/0/ has children at array/1/ and array/2/, the node at array/2/ has children at array/2/ has children at array/3/ and array/3/ and so forth. This array representation, usually called a heap, has both advantages and disadvantages. The main

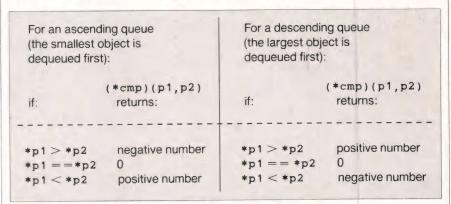


Table 1: Values returned from the comparison function

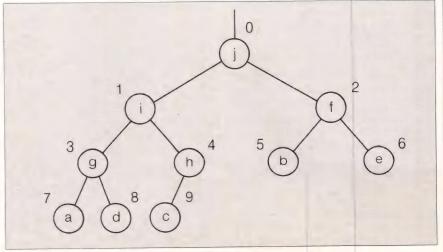


Figure 1: A priority queue represented as a binary tree

C CHEST

(continued from page 105)

problems are that the maximum number of elements must be known at create time, and it's difficult to delete an interior node at random or to merge two queues. Nonetheless, the heap representation is the most efficient for the overwhelming majority of applications—in which the only operations are insert and delete.

The tree from Figure 1 is shown as an array in Figure 3, right. To insert a node in the tree, you make the array one element larger, put the new node in the rightmost element, and then reheap from the bottom up (right to left). The largest node is always at array[0], just as it was always at the root position in the tree. Delete the largest element by copying the rightmost node in the tree (array[9] in Figure 3) into array[0], making the array one element smaller, and then reheaping from the top down (left to right).

The heap representation is a reasonably efficient one. Unlike a normal binary tree, you never have to search for the insertion point in the tree (it's always at the far right). Similarly, the largest element is always in a fixed place (at the far left). Though there's a certain amount of copying that has to be done during a reheap, no more than log₂N copies need ever be done. Nonetheless, it's worthwhile to make the queue elements themselves small. Use an array of pointers to structures rather than an array of structures. This way you have to swap only two pointersrather than two complete structures-when you reheap.

Note that a sorted array is a valid heap (though a heap is not necessarily a sorted array). This property explains the *initheap* argument to $pq_create($). If your input is already a sorted list, there's no point in doing a series of insert operations to create the heap—you can just pass the array directly to $pq_create($).

The priority queue is represented internally by the following structure:

typedef struct
{
 int (*cmp)();
 int (*swap)();

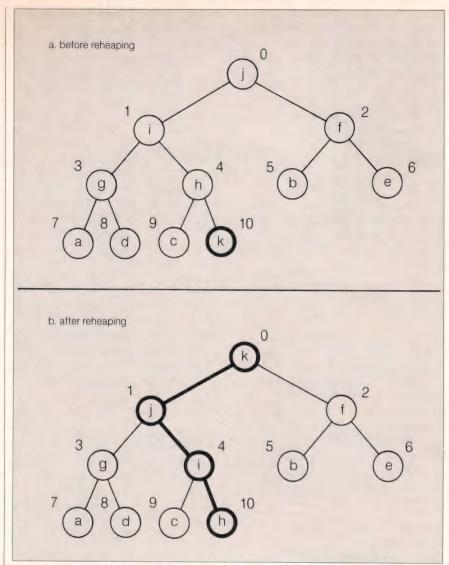


Figure 2: Inserting a new node

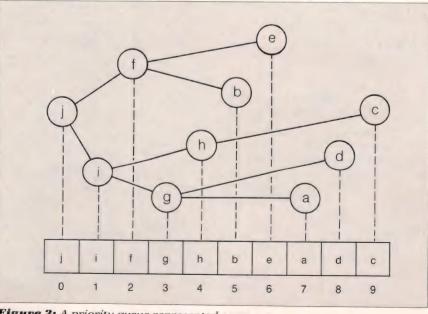


Figure 3: A priority queue represented as an array

C CODE FOR THE PC

source code, of course

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U. S. Atlas (29,000 cities with retrieval program)
KST Fonts (13,200 characters in 139 mixed fonts: specify TEX or bitmap format)
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U. S. Map (15,701 points)

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C CHEST

(continued from page 106)

int itemsize; int nitems; int maxitem; char *bottom; char *heap;

PQ;

The comparison and swap function pointers are remembered in *cmp* and *swap. Itemsize* is the size in bytes of one element of the heap, *nitems* is the number of elements currently in the heap, and *maxitem* is the maximum number of items that the heap can hold. The *heap* field points at the array itself, and *bottom* points at the

most recently added element. It will point at heap[-1] if the heap is empty. Note that you could derive nitems by subtracting bottom from heap but it's more convenient to keep it as a separate number.

The pq_create() function on lines 128–188 of Listing One creates a PQ structure and initializes the various fields. If *initheap* is not NULL, space for the PQ structure alone is allocated (on line 167) and heap is made to point to the specified array. The various other fields are adjusted so that the heap is full. If *initheap* is NULL, space for both the PQ structure and the heap itself is allocated (on line 176), and an empty heap is initialized. Finally, a pointer to the PQ structure is returned on line 188.

The insert and delete operations are performed in pq_ins()(lines 193-219) and pq_del() (lines 224-258) using the method described earlier. Because the compiler doesn't know the size of a heap element at compile time, items must be copied with memcpy() calls at run time. The reheaping is deferred to two static workhorse functions-reheap_ down() and reheap_up()—declared at the top of the file. Reheap_down() (lines 51-89) starts at the root node (heap[0]) and works down the tree. It selects the larger of the two children on lines 72-81 and then swaps the root element and the larger child (and moves down the tree) if necessary. The reheap process stops as soon as the root is larger than both

Flotsam and Jetsam

Standard #include Files At present I'm using the Microsoft C compiler, Version 4.0, for two reasons-the first is the CodeView debugger, and the second is the degree of Unix compatibility. (The MS-DOS compiler is the Xenix compiler: there's literally no difference.) I spend a lot of time bouncing around between the Unix system at Berkelev and my own PC, so having compatible compilers is essential. Using the Microsoft compiler, I've never had to modify a program written under one system and ported to the other unless that program did some sort of verylow-level communication with the operating system (such as talk directly to the BIOS).

There's another advantage to Unix compatibility, and that's standardization. Although the emerging ANSI standard finally incorporates the I/O library into the language—at least in terms of regularizing the function names and argument-passing conventions-there aren't nearly as many ANSI functions as there are Unix functions, so Unix must continue to provide a de facto standard. You don't really know C unless you know the Unix library. Your code just won't be portable because you don't know what's standard and what isn't. Learning your own compiler's library is a necessary but not

sufficient condition for knowing C. As a consequence, it's very much to your advantage to pick up a copy of the Unix System V documentation and read through it. Hitherto, these manuals have been hard to come by. All that you could get were the nowout-of-date Version 7 manuals (the big blue and green paperbacks). AT&T, however, has just published the manuals for Release 2.0 of Unix System V. The complete five-volume set is overkill unless you're really using Unix. Nonetheless, Volume 2-System Calls and Library Routinesshould be in every C programmer's library. It's available by mail order (for \$29.95) from the Computer Literacy bookshop in San Jose, California, and is well worth the money. (Computer Literacy's phone number is [408] 435-1118. The five-volume set is Steven V. Earhart's Unix Programmer's Manual [New York: Holt, Rinehart & Winston, 1986].)

Because I develop all the programs that appear in C Chest in the Unix environment (the Microsoft environment is the Unix environment for all intents and purposes), it seems worthwhile to list the various #include files that I use in C Chest. These are usually presented without comment because they're both standard and well documented elsewhere. If your compiler doesn't have one of

these files, then you don't need to *include it in your program because none of the library functions will require any of the things included within the file. By the same token, if your compiler doesn't have one of these files, then it's not Unix compatible, in spite of what the advertisements may say.

It's impossible for me to describe how to port programs to every nonstandard compiler on the market. You'll have to learn how to do this yourself. You'll need to know your own compiler's library pretty well and have at least a working knowledge of the Unix equivalents to various library functions. Do your homework. I will say, while on the subject. that the Lattice compiler is one of those that falsely claims Unix compatibility-for example, it doesn't support a Unix-compatible stat() function, it uses the name fork() in incorrect ways, and it doesn't have a Unix-compatible #include file system. These inconsistencies are surprising because Lattice seems to have gone to a lot of trouble to add Unix-compatible functions to its library in the last release. Unfortunately, the degree of Unix compatibility that is indeed present lures you into a false sense of security.

The various #include files that you're likely to see in a C Chest pro-

children. Note that an ascending queue (one where the root holds the smallest rather than the largest element) can be created by modifying the comparison function, without touching the reheap code at all.

Reheap_up() (lines 95–124) reheaps in the other direction. It starts with the most recently entered item (which is at heap[nitems-1]) and works up the tree to the root. The routine is simpler than reheap_down() because you don't have to find the larger sibling—a child has only one parent. Again, the reheaping processes stops as soon as a parent node that is larger than the current child node is detected.

The remainder of the file is compiled only if *MAIN* is *defined. In this

case, a stand-alone test program is compiled. This program creates a ten-element-long heap of character pointers and then adds or deletes strings from the heap according to commands entered from the keyboard. Type i < string > to insert <string> into the heap, d to delete an item, and q to exit from the program. The priority queue is created either on lines 359 or 361, depending on the value of Makequeue. If Makequeue is set, an empty queue is created; otherwise, a preinitialized queue lusing the array declared on lines 353-356) is created. If Ascending is true, the queue is an ascending priority queue (items will be removed in ascending order); otherwise, it's descending. There's no explicit code to modify Makequeue and Ascending (I just modified them with CodeView, the Microsoft debugger, as I was debugging).

New items are inserted into the queue on line 386 and deleted on line 395. The comparison and swap functions are declared on lines 289–313. Finally, the *strsave()* function, used on line 389, is shown in Listing Two, page 96. Because this is just a small test program, I'm ignoring the error return from *strsave()*—you shouldn't do this in a real application, of course.

FREQ.C

The priority queue routines are of general utility. You need a few special-purpose utilities to make Huff-

gram are listed below. Most of these are both Unix and Microsoft compatible, but some are used just in the DOS compiler. Note that this isn't a list of all the Unix/Microsoft #include files—just the ones I'm likely to use.

ctype.h—contains various text processing macros: isalpha, isupper, islower, isdigit, isxdigit, isspace, ispunct, isalnum, isprint, isgraph, iscntrl, isascii, toupper, tolower, and toascii.

dos.h—not Unix compatible. Contains #defines for the MS-DOS interface functions: bdos, dosexterr, int86, int86x, intdos, and segread.

errno.h—#defines for the various error condition codes returned from the I/O library: EZERO, EPERM, ENOENT, ESRCH, EINTR, EIO, ENXIO, E2BIG, ENOEXEC, EBADF, ECHILD, EAGAIN, ENOMEM, EACCES, EFAULT, ENOTBLK, EBUSY, EEXIST, EXDEV, ENODEV, ENOTDIR, EISDIR, EINVAL, ENFILE, EMFILE, ENOTTY, ETXTBSY, EFBIG, ENOSPC, ESPIPE, EROFS, EMLINK, EPIPE, EDOM, ERANGE, EUCLEAN, and EDEADLOCK.

fcntl.h—contains definitions needed to use the unbuffered I/O function open(): O_RDONLY, O_WRONLY, O_RDWR, O_APPEND, O_CREAT,

O_TRUNC, O_EXCL, O_TEXT, and O_BINARY.

io.h—contains function declarations for access, chmod, chsize, close, creat, dup, dup2, eof, filelength, isatty, locking, lseek, mktemp, open, read, rename, setmode, sopen, tell, umask, unlink, and write.

math.h—contains definitions for abs, acos, asin, atan, atan2, atof, bessel, cabs, ceil, cos, cosh, exp, fabs, floor, fmod, frexp, hypot, labs, ldexp, log, log10, matherr, modf, pow, sin, sinh, sqrt, tan, and tanh.

process.h—contains definitions for various process-control functions: abort, execl, execle, execlp, execlpe, execv, execve, execvp, execvpe, exit, _exit, getpid, spawnl, spawnle, spawnlp, spawnve, spawnve, spawnvp, spawnvpe, and system.

signal.h—contains definitions needed by the *signal()* subroutine. Some common definitions are: *SIGINT*, *SIGFPE*, *SIG_DFL*, and *SIG_IGN*.

stdarg.h—contains definitions needed to write an ANSI-compatible subroutine with a variable number of arguments (see varargs.h). The following are defined: va_list, va_start,

va_arg, and va_end.

stdio.h—contains #defines and so on for all the buffered I/O functions (fopen, fprintf, and so forth). Stdin, stdout, stderr, FILE, EOF, and NULL are all defined here. Note that several psuedofunctions that you're used to thinking of as subroutines (getchar, putchar, getc, putc, feof, ferror, and fileno) are actually macros that are #defined in stdio.h. If you forget to #include this file, you'll get error messages from the linker (such as "Unresolved external:_putchar"). Putchar is a macro that ultimately evaluates to a call to a system-level subroutine usually called either _flsbuf or _flushbuf.

sys/stat.h—contains definitions for the system-status subroutines *stat()* and *fstat()*.

sys/types.h—also required by *stat()* and *fstat()*. Various time functions (such as *utime()*) use the information declared here as well.

varargs.h—Definitions for the variable-number-of-argument mechanism used by Unix System V (see stdarg.h). The following are defined here: va_list, va_dcl, va_list, va_alist, va_arg, and va_end. C

(continued from page 109)

man trees, however. One of these is the FREQ.C program shown in Listing Three, page 96. FREQ.EXE is a standalone program that takes as input a list of files and outputs a table showing the frequency of occurrence of every 8-bit pattern in these files (I'll call these patterns "characters" from here on). The output table is sorted in ascending order of frequency (least frequently occurring characters toward the top). The output format uses one character per line with three numbers on each line-the leftmost number is the value of the character itself (in hex); the next two numbers are character frequencies. output in two forms. The first form is the probability of occurrence of every character-the number of times that that character occurred in the input divided by the total number of input characters. The second form is a normalized character count in which the least frequently occurring nonzero pattern has the value of 1it's the character count divided by the count associated with the least frequently occurring character.

The table itself is declared on lines 13–21. It is a 256-element array of *ITEMs*, indexed by character value. One field of the *ITEM* structure holds a count that is incremented every time the associated character is encountered in the input. The other

field holds the character value itself. You can derive this from the index, of course, but putting it into an *ITEM* lets you sort the array by frequency of occurrence without loosing the value.

The array is loaded with the for loop on lines 58-74. The count associated with every byte is incremented on line 69. The counts are all initialized to 0 by default because globallevel objects are always initialized to 0 unless there's an explicit initializer present as part of the declaration. The next for loop, on lines 83-93, initializes the value fields of the array and at the same time finds the element having the smallest nonzero value. (I'm getting the largest value, too, but am not using it for anything at present.) The array is sorted with the ssort() call on line 105. This subroutine works just like the Unix-compatible gsort() does, but it does a shell sort. See the comment on lines 95-102 for more information. The comparison function used by ssort() is declared on lines 25-39. Finally, the table is printed (and the probabilities and so forth are computed) in the loop on lines 115-126. The sum of the probabilities is printed to stderr on line 126 just to make sure that everything worked correctly. It should always be 1.00.

Availability

All the source code for articles in this issue is available on a single disk. To

order, send \$14.95 to *Dr. Dobb's Journal*, 501 Galveston Dr., Redwood City, CA 94063 or call (415) 366-3600 ext. 216. Please specify the issue number and format (MS-DOS, Macintosh, Kaypro).

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DD.

(Softstrips begin on page 92.)

(Listings begin on page 94.)

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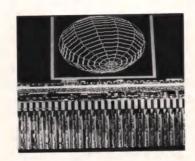
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Resources

Ward, Robert. *Debugging C.* Indianapolis, Ind.: Que Corp., 1986. 350 pages with index. ISBN 0-88022-261-1.

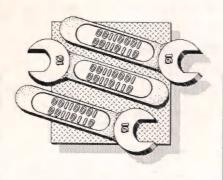
Finding this book among the flotsam and jetsam of computer publishing is like finding a \$100 bill in the street. The art, craft, and/or skill of effective debugging is a topic that has rarely been discussed in any useful way in the computer literature. The debugging chapters in manuals or textbooks are usually focused on defensive strategies and the design of test data, rather than on the stabilization, isolation, localization, and correction of bugs once detected. To new programmers, who have not yet developed a methodical approach to debugging, the process seems a little magical and their own approach is frequently haphazard and based largely on luck. Although native talent is an element of effective debugging, as it is of effective programming, a far larger component is simply the accumulated experience and the strategies learned by finding and fixing many different types of bugs in the past.

In his introduction, Mr. Ward writes: "Until I started teaching, I assumed that good debugging skills were a natural outgrowth of good design skills. Not so. A bright student may intuitively decompose a problem into beautifully coherent, cohesive, functional modules. That same student may not be able to find the

by Ray Duncan

most trivial syntax errors, let alone discover subtle runtime bugs. Equally bright students turn in working designs that literally defy analysis. While I don't believe that we learn debugging by studying design, I do believe that we can learn efficient debugging.

"We can develop a methodological



model that directs our efforts toward more productive searches. We can acquire heuristic knowledge (a kind of folk wisdom) about where to look first. We can be deliberately sensitive to the different variables and observable phenomena in different environments. We can become expert at selecting and using appropriate tools. And, through critical analysis of our attempts to find 'worthy' bugs, we can learn from our own mistakes."

The first chapters of Debugging C concentrate on the debugging process itself, with emphasis on recognition of bugs (lexical, syntactic, execution, and intent errors), their localization (using the principles of lexical, temporal, or referential proximity), a methodical approach, and good record keeping. Later chapters discuss the localization of compiletime errors and tracing methods. The last chapters of the book are particularly C-oriented, with excellent discussions of the special problems of C programmers: bugs due to data type mismatches, operator precedence or misuse, and uninitialized or out-of-range pointers. Here is where the author gently presents many useful tips that are common sense to veteran C coders, usually acquired through painful experience.

Example: "C, unlike Pascal and Basic, doesn't give an initial value to local variables when it creates them.... Most stack frames are relatively small (less than 20 bytes) but every stack frame has a frame pointer (an address higher in the stack) and a return address (an address in the code area). Because uninitialized

pointers will use these obsolete addresses which occur frequently in reused areas of the stack, the programmer can expect 10 to 30 percent of the uninitialized pointers to reference code or areas of the stack." Similarly, Mr. Ward explains, with stack traces and memory dumps, why common errors such as initializing too many bytes of an automatic character array result in bizarre transfers of control or crashes "between lines" of source code.

The last few chapters describe the use of some common machine-level and symbolic debuggers, interpreters, and integrated development environments. The book's appendices include the C source code for a debugging subsystem that can be linked to a C application. The debugger provides tracing at several layers of granularity, stack frame displays, memory dumps, and watch points (periodic checks to see if a certain variable has been altered or has taken on a specified value).

The author couples a direct, lucid, and informal style of writing with a depth of understanding and highly structured presentation that are unusually effective. In spite of the title, most of the book has broad applicability and would be perfectly comprehensible to anyone with even a passing acquaintance with the C language.

Hyman, Michael I. Memory Resident Utilities, Interrupts, and Disk Management with Ms and PC DOS. Portland, Oreg.: Management Information Source, 1986. 373 pages with index. ISBN 0-943518-73-3.

This book, with the short catchy title, is anything but self-effacing. The introduction reads: "You will find this book to be the ultimate reference guide to getting the most out of your machine. As you read it, you'll learn how to use and enhance DOS to ex-

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plore your computer and make mighty programs. Powerful examples will lead you along the way and provide models for later reference "Such hype is commonplace from publishers' PR departments, but it's a little more unusual coming straight from the author's pen. Several paragraphs farther on, you find the admonition "The letter 'I' and the number '1' are represented in the program code by the same character. In no case is the letter 'I' used as a variable name. This should eliminate any confusion." If there's anything I'm already confused about at this point, it's why the publisher of an "ultimate" reference book would allow it out of the door containing program listings wherein the numeral 1 and the letter l cannot be distinguished. But no matter, let's see what the rest of the book has to offer.

Chapters 1 through 8 discuss the boot sector, file allocation table, directory structure, and file area of MS-DOS disks. Incredibly tortuous Turbo Pascal source code for a disk peeker/ patcher named Explorer is developed as part of the exposition. Each chapter ends with a summary of sorts entitled "Key Programming Points." Here are the Key Programming Points listed at the end of Chapter 4: "Sectors are the smallest organizational unit. They contain 512 bytes. You can find interesting messages and modify programs by editing sectors. Explorer uses arrow keys to move the cursor."

The following few chapters discuss the partition table, "un-erasing" files, and a program that patches COMMAND.COM to change the names of MS-DOS internal commands (real useful!). Next the author covers input and output with the mouse, keyboard, and screen using IBM PC ROM BIOS drivers; an overview of file and record I/O; directory searching; and memory management. Nothing new here.

Finally, in Chapters 28 through 38, you get to the apparent reason for the book's existence: the problems and pitfalls of programming Terminate and Stay Resident (TSR) utilities. This part of the book contains much useful information, poorly organized

and presented, about chaining onto interrupts, putting up and taking down pop-up displays, monitoring the keyboard for hot keys, and the like. If this part of the book were properly structured and edited, it would make a decent magazine article, but the book as a whole is a poor investment.

Norton, Peter; and Socha, John. Peter Norton's Assembly Language Book for the IBM PC. New York: Brady/ Prentice-Hall, 1986. 413 pages with index. ISBN 0-13-661901-0.

This book is Yet Another Assembly-Language Tutorial of average quality. It is built around the design and stepwise enhancement of a simple disk sector modification utility called DSKPATCH, discussing in passing some issues of processing keyboard input and updating screen displays using the IBM PC ROM BIOS video driver.

This book is mainly notable as a demonstration of the recent trend toward commodity marketing of computers and related products. As the competition in the wonderland of silicon has become more intense, we have seen the adoption of marketing tactics that were previously the province of car, appliance, and apparel manufacturers, such as scratch-off sweepstake tickets (Borland); cash rebates (Apple); "free gifts with purchase" (with every dBASE III Plus, get a Cross writing instrument free!); mystical mumbo jumbo such as Robert Carr, Wayne Ratliff, and Jonathan Sachs being touted as "Chief Scientists" of their respective companies; and last but not least, celebrity endorsements.

Back in 1983-1984, while John Socha was a contributing editor for Softalk/PC, he wrote a book called Assembly Language Safari for the IBM PC: First Explorations (Bowie, Md.: Brady, 1984. ISBN 0-89303-321-9). For various reasons, including rather poor production values in the book itself and massive financial problems at the Brady corporate level, the book received little attention and went out of print shortly thereafter. Nowadays, John Socha works for Peter Norton Inc. in Santa Monica, California, and is the author of the program sold as the Norton Commander. When Peter Norton's Assembly Language Book appeared, with Peter Norton and John Socha listed as coauthors, I thought it would be instructive to make a page-by-page comparison of John Socha's old book with the new one.

A fairly generous assessment of the two books is that there are approximately 17 pages of new material in the "new" Norton/Socha version, scattered among the following topics: the Proceed command of DEBUG, a MASM program skeleton, MAKE files, SYMDEB, linker maps, .COM vs. .EXE programs, the ASSUME directive, segment overrides, and phase errors. The remainders of the two books are identical, except for some redrawn figures, the addition of some titles and divider pages, and some minor changes in wording that would have been introduced by any competent editor. There are also three new appendices that are mostly filler from other sources, such as MASM error messages and character tables. In other words, the substance of Mr. Norton's contribution to this book seems to be his picture on the cover, his billing as a coauthor, and the running head "Peter Norton's Assembly Language Book" on the top of each right-hand page.

Some might argue that the kind of misrepresentation involved here hurts no one and is therefore of no consequence. To be sure, the book's purchasers, although they might have been misled about the creative origins of the book, have still bought a reasonably useful introduction to assembly language. The three principals—John Socha, Peter Norton, and Brady-are undoubtedly happy because their ploy has caused the book to vault onto the computer best-seller lists-which means everyone involved has made some money. And last but not least, Peter Norton's recognition as an expert on all matters concerning the IBM PC has been magnified.

A Poor Man's MAKE

J. F. Philippe Marchand, of Webster, New York, sent in the programming goody of the month—a short program called CHKDATE.C. This program compares the date of two files and returns an "error level" code that can be tested within a batch file. For those people who are not fortunate enough to have one of the commercial MAKE utilities (which are bundled with many of the compilers and assemblers being sold today), CHKDATE can help to automate the process of compiling and linking the various modules of an application program. Example 1, below, contains the C source code for the CHKDATE program, and Example 2, below, contains an example batch file that demonstrates the use of CHKDATE.

MS-DOS Programming Tips

George Smith, of Lilburn, Georgia, writes: "I have run across one puzzling problem [in MS-DOS 3.0]. DOS function *0eh*, which makes a specified drive the current drive, has al-

ways reported a reliable count of the number of drives on the system in register al. Under DOS 3.0 and later, though, the value returned in al is always at least 5, even if you happen to be running short-handed with fewer drives.

"I have no explanation, though I do have a solution in the form of function *DriveCnt* [Example 3, below]. The logic behind *DriveCnt* is simple. It takes the value *al* reports from function *0eh* and passes it as a drive code to function *47h* (Get Current Directory), which will set the carry flag if the drive code is invalid. It repeats calls to function *47h* with successive-

ly lower drive codes until DOS

finds one it likes."

Although George has been kind enough to send in the routine *Drive-Cnt* for the edification of *DDJ* readers, he sells a package called Boosters for Turbo Pascal programmers that includes this subroutine and 77 others, a screen generator, some 40 example programs, and a 93-page manual. The Boosters package costs \$40 and can be ordered direct from George Smith at 609 Candlewick Lane, Lilburn, GA 30247; (404) 923-6879.

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```
#include <stdio.h>
#include <sys types.h>
#include < sys stat.h>
main(argc,argv)
int argc;
char *argv[]:
{ struct stat buf 1, buf 2;
  int result, k;
  if ( argc < 3 ) {
   printf("chkdate: usage chkdate f1 f2 .. fn\n");
   printf("
              will return errorlevel 1\n");
   printf("
               if f1 older than f2 .. fn\n");
   printf("
               or if f1 .. fn don't exist. \n");
   exit(1):
  if ( stat( argv[1], &buf1) != 0 )exit(1);
  for (k=2; k < argc; k++)
    if ( stat( argv[k], &buf2 ) != 0 ) exit(1);
    if ( buf1.st_atime < buf2.st_atime ) exit(1);</pre>
  exit(0):
```

Example 1: Phil Marchand's CHKDATE.C program

```
chkdate main.obi main.c
IF ERRORLEVEL 1 goto : compile1
goto:next1
:compile1
msc main, main;
:next1
chkdate func.obj func.c
IF ERRORLEVEL 1 goto : compile2
goto:next2
:compile2
msc func, func;
:next2
chkdate main.exe main.obj func.obj
IF ERRORLEVEL 1 goto : link
goto:exit
link main.obj + func.obj, main.exe;
:exit
```

Example 2: TEST.BAT file, a demonstration of the use of the CHKDATE program in a batch file to automate the compilation and linking of an application

```
FUNCTION DriveCnt : Integer;
         (C) 1986 George F. Smith & Company
         Purpose:
            Returns number of logical drives on host
            computer - DOS 2.0 and above.
         Sample Usage (Turbo Pascal):
            Writeln('Number of drives is ', DriveCnt);
         Suggested processing sequence:
            MASM DriveCnt,,,
            Link DriveCnt
            Exe2Bin DriveCnt DriveCnt.com
            C2I DriveCnt.com >DriveCnt.inl
             (C2I utility converts .com files to Turbo Pascal
            inline code. See DDJ, 10/86, pg. 90 )p
         segment
         assume cs:code
DriveCnt:
                        ; standard
         push bp
         mov
               bp, sp
                              ; subroutine
         push ds
                              ; overhead
               sp, 64
                              ; create scratch area,
         sub
                             ; save in si for 47h call
               si, sp
                ah, 19h
                              ; get default drive
               21h
dl,al
         int
                               ; returns drive code in al
                               ; move code to dl
         mov
               ah, Oeh
                              ; set default drive to itself
         mov
         int
               21h
                               ; On return, al=# drives
      If DOS 3.0 or above, number of drives in al will
be minimum of 5. Will repeat calls to function 47h
until a valid drive code is obtained.
:
                              ; number of drives in dl
         mov
               dl.al
                              ; segment address
         push ss
                                   of scratch buffer
               ds
         pop
search:
                ah, 47h
                               : is drive code okay?
         mov
                21h
                               ; carry flag set if
          int
                               ; drive code invalid
                               ; jump if code valid
          inc
                okay
                               ; drive code too big;
          dec
                               ; decrement it and
          qmt
                search
                                   try again
                               ; clear dh
okav:
          xor
                dh, dh
                [bp+4], dx
                              ; give results to caller
          mov
                              ; adjust stack ptr
          add
                sp, 64
                               ; restore caller's regs
          pop
                ds
                sp, bp
          pop
                bp
          ret
code
          ends
          end
               DriveCnt
```

Example 3: George Smith's DriveCnt routine to determine the number of disk drives present in an MS-DOS

Object-Oriented Programming in SCOOPS

This month, I conclude my review of PC Scheme (which I consider to be the Turbo Pascal of the PC LISP family) with some examples of object-oriented programming using SCOOPS, the object-oriented extension of the language. To do this, I'll have to make some extensions to PC Scheme itself—but first, some discussion about LISP programming in general.

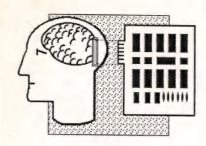
LISP is the most organic and lifelike of all programming languages. Most of its dynamic character comes from the combination of complex, nested structures with dynamic reassignment of structures of pointers and a simple syntax that uses the same representation for data and programming code.

The functional programming aspect of LISP involves a special implementation of argument passing so that usually few variables need to be stored permanently. The main thing in pure functional programming is not modifying objects in permanent storage but passing symbols as if they were values being passed between mathematical functions. The main result of such a program is the structure it returns rather than the state it creates in the permanent storage of the machine.

But LISP is not simply a single-paradigm programming language. It is a language that so far has been able to absorb each new programming model as it appears and to incorporate these models in a functioning whole. Over the years it has absorbed other programming concepts and is con-

by Ernest R. Tello

tinuing to do so—for example, as I mentioned last month, SCOOPS has assimilated some features of Smalltalk. The original model on which the LISP language was based was that of functional programming, using the lamb-



da calculus.

In pure functional LISP programming, anything that does not just return a structure but modifies the machine is generally considered a side effect. But it is often important in an object-oriented environment to modify the object hierarchy dynamically in complex and carefully controlled ways. Take, for example, the case of performing simple list processing functions, such as updating and modifying list structures. Here it is often the case that either or both of the functions of returning the necessary structure and producing the necessary structure in permanent storage are important parts of the required

In the following example, I will show various versions of a function add-to-end that are implemented so as to return different values and produce different side effects. This function extends the list processing functions of LISP to include the ability to add an element to the end of an already existing list structure.

The function give-n-take was written to demonstrate the side effects of the add-to-end function. First, two lists are created—nums, which contains the list of number words (one two three four five); and morenums, which is composed of the complementary number word list (six seven eight nine ten). Here is what give-n-take looks like:

(define nums '(one two three four five))
(define morenums '(six seven eight nine ten))

(define (give-n-take)

(add-to-end (car morenums) nums)

(set! morenums (cdr morenums)))

As you can see from the short session that follows, what give-n-take returns is different from the side effects it has on these lists. It simply returns the morenums list that was passed to it. When you ask LISP for the contents of these lists by typing their names at the interpreter prompt, however, you see the effects that give-n-take has each time it is called. It takes numbers successively from the beginning of the morenums list and adds them to the end of the nums list:

[2] nums (ONE TWO THREE FOUR FIVE) [3] morenums (SIX SEVEN EIGHT NINE TEN) [4] (give-n-take) (SEVEN EIGHT NINE TEN) [5] nums (ONE TWO THREE FOUR FIVE SIX) [6] morenums (SEVEN EIGHT NINE TEN) [7] (give-n-take) (EIGHT NINE TEN) [8] nums (ONE TWO THREE FOUR FIVE SIX SEVEN) [9] morenums (EIGHT NINE TEN)

Now compare this to another version of the function called add-to-end-2. In the next session I create the list of integers from 1 to 4. My original add-to-end returns something unusable, but the side effects are the correct result. Add-to-end-2 does just the opposite—it returns the list with the number added to the end, but when you examine the list of integers, you see that nothing has changed.

[6] (define integers '(1 2 3 4)) INTEGERS [7] (add-to-end 5 integers) (4 5) [8] integers (1 2 3 4 5) [9] (add-to-end-2 6 integers) (1 2 3 4 5 6) [10] integers (1 2 3 4 5)

This hasn't been just an academic exercise. The side-effects version of add-to-end is useful for doing necessary housekeeping in an object-oriented LISP environment. It is valuable to be able to maintain lists of all the current instances of various classes that are alive in a system that is changing dynamically. Without the version of add-to-end that can actually modify such lists, you would not be able to update them continually.

Another function that could be useful is *delete-last!*. It performs the opposite service—that of destructively removing the final element in a list. Its definition is:

(define (delete-last! lst) (delete! (car (last-pair lst)) lst))

The *last-pair* function in PC Scheme returns the last pair in a list. The function works for returning the last element as a single element list because the *car* reduces the pair to a simple list structure.

The following quick session shows the behavior of *delete-last!*. As you can see, in this case both what the function returns and its side effects are identical.

[2] (define numbers '(one two three four five))

NUMBERS
[3] (delete-last! numbers)
(ONE TWO THREE FOUR)
[4] numbers
(ONE TWO THREE FOUR)

Programming in SCOOPS

Sending messages is a rather simple matter of using the *send* function with the *receiver* object and the message to be sent plus its arguments. So, to send a message to the *my-body* object (introduced in the May column), giving it a new body part called *toes*, you would say:

[1] (send my-body put-cpart-name 'toes)

[2] body-parts (HEAD NECK ARMS HANDS TRUNK LEGS FEET TOES) If you then changed your mind and decided to remove this new body part, you could do so by globally accessing the body-parts list using the newly defined function deletelast! as follows:

[3] (delete-last! body-parts)
[4] body-parts
(HEAD NECK ARMS HANDS TRUNK LEGS FEET)

The code in Listing One, page 98, demonstrates simple inheritance in PC Scheme through several levels of a fairly linear hieararchy. First, the root class artifact is defined with the instance variables material, weight, purpose, and cost. Then transportmeans is defined as a subclass of artifact with the additional instance variables medium, time-range, and power-source. Naturally, transportmeans inherits all the variables and methods of the artifact class. Then transport-vehicle is defined as the next subclass and passenger-vehicle as a subclass of it. Descending further in the same linear manner of adding more and more specific classes that inherit everything from the previous class, the classes water-transport-vehicle, surface-vessel, ship, and oceanliner are defined. The instance object ship1 is created as an instance of the class ship. The two methods speed and direction are also provided for the ship class.

Listing Two, page 99, provides an example of multiple inheritance in SCOOPS. Here I have implemented the example used in April's column in PC Scheme. First, the classes business and adversary are defined. Then the class competitor is defined and uses the multiple inheritance feature to inherit everything from both of these two classes.

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All the source code for articles in this issue is available on a single disk. To order, send \$14.95 to *Dr. Dobb's Journal*, 501 Galveston Dr., Redwood City, CA 94063 or call (415) 366-3600 ext. 216. Please specify the issue number and format (MS-DOS, Macintosh, Kaypro).

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(Listings begin on page 98.)

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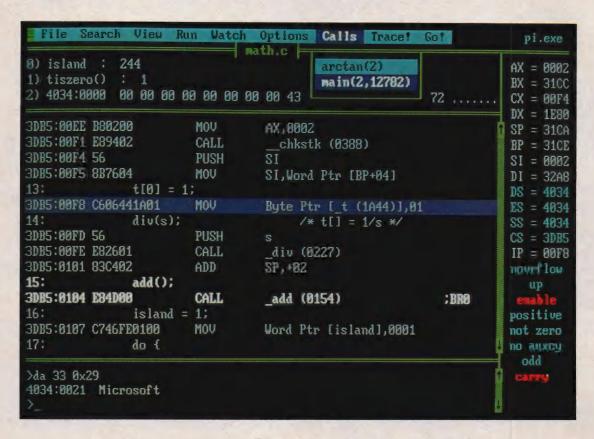
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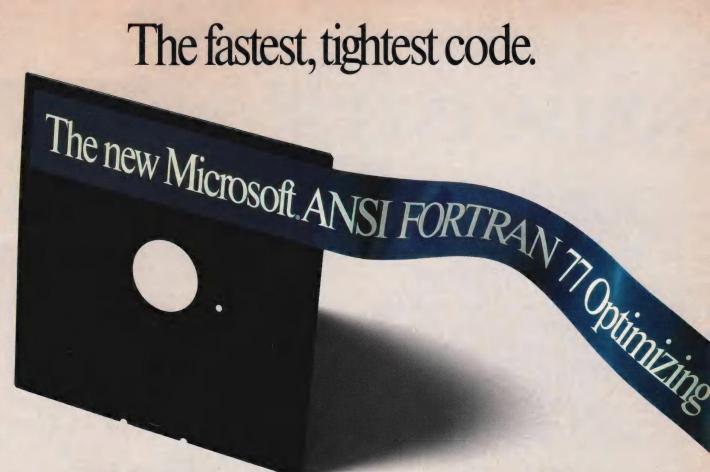
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fastest and just about the handiest editor under ten fingers (if it isn't, just redefine your macros the instant you think of something better). Just take the compiler and all-purpose resident macro program as bonuses.

Now, the importance of speed is simply that we want to work at programming, not at running an editor. Like a fine sound system, a good editor should be transparent. When I write while $v \le maxv \ do$, I don't think of the keys I press, only of the statement I'm producing. It should be the same when I want to move that statement, to indent it and what follows, or simply to erase the line—or to find it in a 2,000-line file. What no one wants is to sit and look at the editor editing.

Which brings us to WordStar imprinting. I got imprinted with Word-Star because my mother imprinted me with ten fingers. For the first two days, it may be easier to use an editor that has F3 for "delete word" and shift-F3 for "delete line," but after those two days, only a WordStarstyle editor gives you a chance to get to the point at which you only have to think "delete line" and not notice the actual fingerwork needed to do it—because it doesn't demand that your hands leave the typing position.

Hunt-and-peck chickens may not understand this, but there is a world beyond the barnyard, you know. The WordStar commands aren't meant for the user's manual but for the user's hands. They're the natural extension of sheer speed because they remove another nontransparent interface between your mind and the program text.

So, my ultimate editor is simply the Turbo Pascal editor with a good mac-

ro program plus several-document capacity, windowing, wordwrap for comments, block-limited search-and-replace with more complex specification capacity, and files greater than 64K. But never at the price of speed or the WordStar keyboard code. Don't mistake us far-voyaging mallards for clay pigeons.

Philippe Ranger 6120 Hutchison Montreal, Canada H2V 4C2

6502 Hacks

Dear DDJ,

When I read Mark Ackerman's "6502 Hacks" (February 1987), I was both surprised and delighted that, in this world of 68000s and hypercubes, there is still anyone left who would spend the time and effort to write a good article about 6502 programming.

But, before anyone's programs start crashing, let me correct Mr. Ackerman on two related points. First, the software interrupt instruction (*BRK*) uses the *IRQ* (maskable interrupt) vector, not the *NMI* (nonmaskable interrupt) vector. But to make matters even worse, the *BRK* instruction pushes its address+2 on the return stack. In order to return to the opcode immediately following the *BRK*, the return address on the stack must be dug up and decremented before returning (very messy).

Second, and more important, is that the *RTI* (return from interrupt) instruction is not functionally compatible with the sequence *PLP* (pull processor), *RTS* (return from subroutine). The *RTI* instruction pulls the processor status byte and then continues execution at the address pulled from the stack, increments the address by 1, and then continues execution at the adjusted address. This compensates for the fact that the *JSR* pushes the address of the next opcode–1.

The reason for this lies in a quirk of the 6502 opcode processing. To execute the *JSR* instruction, the processor first fetches the opcode and the low byte of the address. It then pushes the current program counter (which is now pointing to the third byte of the instruction) and, finally,

```
ENTRY STX SAVEX ;Save [X]
... ;(Body of subroutine)
LDX #$FF ;Restore [X]
SAVEX EQU *-1 ;Data field of preceding instruction
RTS
```

Example 1: A fast method for saving and restoring registers

```
TXBLOCK
            STX
                  _TXLENGTH
                               :Set block length (0=256)
            STA
                 TXADDR
                               ;Set block address
                 _TXADDR+1
                               ; Init timeout
            LDX
                 # - 1
                               ;Init block index
            LDY
                 # 0
            STY
                  _TXSUM
                               :Reset TXSUM accum.
TXBLOOP 1
            LDA
                 $FFFF, Y
                               ; First, get byte to send
_TXADDR
             EOU
                 *-2
                               ; Address portion of LDA
                                instruction
TXBLOOP2
                  HWREADY
             BIT
                               ; Check if Tx port is ready
                  TXBLOOP3
            BPL
                               ; (N) Check timeout
                               ; (Y) Tx the byte
                  HWDATA
            STA
             EOR
                  #$FF
                               : Accum TXSUM
_TXSUM
             EOU
                  *-1
                               ;Data portion of EOR instruction
             STA
                  _TXSUM
                               :Update running XOR sum
             INY
            CPY #$FF
                               . End block?
TXLENGTH
             EOU
                               ;Data portion of CPY instruction
                 TXBLOOP 1
             BNE
                               : (N)Continue sending
                  _TXSUM
                               ; Send the checksum
             LDA
TXBLOOP3
             DEX
                               ; Timeout expired?
            BNE
                  TXBLOOP2
                               · (N)Continue sending
                               ; (Y) Return timeout error
```

Example 2: An actual 6502 code fragment

fetches the high byte of the address. So, the address that gets pushed on the stack is always one byte shy of the next instruction. In contrast, the interrupt acknowledge sequence will only happen between instructions, so the progam counter is always pointing to an opcode when the interrupt return address is pushed.

In his coverage of self-modifying code, I think Mr. Ackerman missed one very useful trick. I often save registers (because they are so very scarce) by storing their contents in the data portion of a load instruction located at the end of the routine [see Example 1, page 122]. It takes 5 bytes of code but is absolutely the fastest method of saving and restoring a register (six cycles total) and both the instruction and data storage are localized in the subroutine.

To demonstrate variations of the trick, I have included a programming fragment of an actual routine [see Example 2, page 122]. The routine transmits, over an extremely fast synchronous communication port, a block of data (from 1 to 256 bytes) followed by the exclusive-ORed sum of the block. It must also make sure that the port is not "hung" by keeping a decaying timer in the X register. The routine is passed the address of the block in registers A and Y and the length of the block in X. When the loop is cooking, it can transmit a byte every 27 cycles.

Of course, I won't tell you what product this code is running in (the labels have been changed to protect the innocent), for fear that someone wouldn't buy our system if they knew I programmed this way on a regular basis.

James Bucanek C-Si Systems 572 W. Pima Coolidge, AZ 85228

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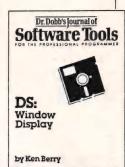
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VIEWPOINT

(continued from page 14)

ple) to use simple trigonometry in a high-level language than it is to use it in assembly language.

High-level languages can enhance the value of a programmer's work by allowing code to run with very little alteration on many different computers-all that is required is a compatible compiler for each. My 68000 cross assembler is now operational on computers that use Z80, 8088, and 68000 processors. I would like to challenge Suman-or any assemblylanguage programmer-to convert an assembly-language program of similar functionality and complexity to run in two environments. I suspect that it would take them longer than the few days each that it took to port the 68000 cross assembler I presented in the April and May 1986 issues of

With reference to Suman's specific criticisms, the use of named constants (for any language, including assembly language) is almost universally considered to improve maintainability by localizing changes. Although I used the constants FIRST and LAST only twice (not once, as Suman stated), I feel their use was justified.

As to the lack of initialized variables in Modula-2, I agree that this is a defect of the language. That should not, however, be an indictment of all high-level languages—C does allow initialized variables, including initialized structures and arrrays. If Modula-2 allowed initialized variables, there would have been no need for InitOperationCodes (the module that Suman disliked so much). The mnemonic lookup table could have been created at compile time—in the module that needed the table—and with much less overhead.

Certainly, the mnemonic lookup table of X68000 could have been written more efficiently using assembly language. It would have been harder to write or to read, however, and it would not have been portable (some processors invert the order of bytes within words). The sets that were used for the allowable addressing modes (ModeA and ModeB) repre-

sented the 68000 addressing scheme in a recognizable way. For example, one of the set members was called *Size67* to indicate that the sixth and seventh bits of the 68000 operation code were used to indicate the size of the operation. If that member were to be converted to a computer word (1s and 0s), a programmer reading the listing might be at a loss to figure out what the significance of a particular bit (or combination of bits) was.

Finally, Suman seems to contradict his own point about turgidity and redundancy when he suggests that I should have used 118 individual write statements instead of the simple loop:

FOR i := FIRST TO LAST DO
WriteRec (f, Table68K[i]);
FND:

Perhaps an ideal solution to the dilemma of choosing between the expressiveness of high-level languages and the efficiency of assembly language is to mix the two. Most decent compilers allow the integration of assembly-language modules. A program can be written and debugged in Modula-2 or C, then profiled to identify the bottlenecks. Finally, these sections can be rewritten in assembly language. The original highlevel language code can be left behind as a comment block to the assembly language. Although this would hamper portability somewhat, the sections of a program that benefit from recasting in assembly language are usually a small portion of the overall code. Projects that would be Herculean tasks in assembly language become quite comfortable using modern high-level languages. Get your project working, then if you need more speed or better memory utilization, tune it up using a small amount of hand-coded assembly language.

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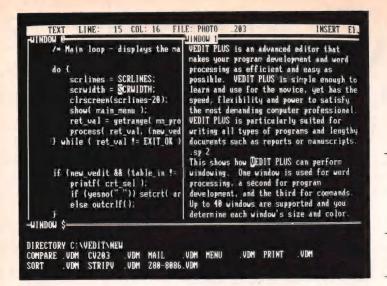
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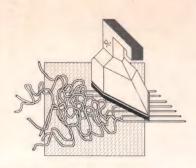
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THE STATE OF BASIC



BASIC Functions

In this issue we'll look at user-defined functions as implemented by QuickBASIC, Turbo BASIC, True BASIC, and Better BASIC.

QuickBASIC 2.0 permits user-defined nonrecursive functions to extend over multiple lines. All function names must start with the letters FN. A data type symbol may be required, depending on the data type returned by the function and any global default name declarations used. A function has an optional list of scalar arguments (no arrays are allowed) with all the arguments passed by value. Multiline functions end with an END DEF statement and can use EXIT DEF to exit from the function.

Most of the new BASIC dialects allow passing of all arguments by value. A function with no arguments that returns a simple value is one way of implementing Pascal-like constants; you cannot accidentally alter the function's value, as may be the case with a variable.

Nonparameter variables that are used within the functions are global. To avoid undesirable side effects and to localize such variables, include them in a *STATIC* declaration. This technique ensures that new addresses are assigned to these variables every time a *STATIC* declaration is encountered.

Turbo BASIC functions are similar to those of QuickBASIC, with the following differences:

- Turbo BASIC functions can be recursive.
- Local variables are declared inside functions using the LOCAL keyword.
- The SHARED keyword is used to explicitly declare global variables.

Turbo BASIC offers *LOCAL* and *STATIC* declarations, which give programmers the ability to explicitly declare local nonstatic variables. In Turbo BASIC, local arrays are first declared in the *LOCAL* list and then dynamically dimensioned using *DIM DYNAMIC*.

True BASIC implements user-defined functions in a slightly different way from that of the previous two BASIC dialects. First, function names do not have to start with the letters *FN*. The price to pay, however, is the use of *DECLARE DEF* declarations to inform True BASIC of the function

```
DEF FNFACT(N%)
'QuickBASIC nonrecursive factorial function
STATIC I%, F' static local variables
F = 1' Initialize
'loop to get factorial
FOR I% = 2 TO N%
F = F * I%
NEXT I%
FNFACT = F
END DEF
```

Example 1: QuickBASIC listing for a nonrecursive factorial function

```
DEF FNFACT(N%)
' Turbo BASIC recursive factorial function
IF N% > 1 THEN
FNFACT = N% * FNFACT(N%-1)
ELSE
FNFACT = 1
END IF
END DEF
```

Example 2: Turbo BASIC listing for a recursive factorial function

```
DEF Fact(N)
! True BASIC recursive factorial function
IF N > 1 THEN
   LET Fact = N * Fact(N-1)
ELSE
   LET Fact = 1
END IF
END DEF
```

Example 3: True BASIC listing for a recursive factorial function

```
REAL FUNCTION: Fact
INTEGER ARG: N
EXTERNAL: Fact
10 IF N > 1 THEN RESULT = N * Fact(N-1) ELSE Fact = 1
END FUNCTION
```

Example 4: Better BASIC listing for a recursive factorial function

```
REAL FUNCTION: LOGN

REAL ARG: X

REAL ARG: BAISE/OPT=10

10 RESULT = LOG(X) / LOG(BAISE)

END FUNCTION
```

Example 5: Simple power function in Better BASIC that demonstrates the default parameter feature

names imported from external libraries and modules. True BASIC supports recursive multiline functions that take scalar- and array-type parameters. All the function parameters are passed by value.

In essence, True BASIC supports two levels of functions: internal and external. Internal functions are those located in the main program, before the unique END statement. All the variables in the internal function and not in the argument list are global, which enables internal functions to create and manipulate global variables. External functions are located either after the END statement of the main program or in an external library or module file. External functions defined in modules can access information through the argument list, SHARED variables, and PUBLIC variables. External functions in libraries have a strict data interface because they rely mainly on the argument lists. Argument lists in True BASIC can contain a file I/O channel number that allows file I/O, which provides another method of accessing large data that are stored in intermediate files.

Better BASIC approaches the implementation of functions in a Pascallike fashion: the type of the function or its arguments is explicitly declared using data-type keywords and not symbols. Much of the discussion about Better BASIC procedures that appeared in the May column applies to functions. In addition:

- Function names do not need to start with the letters FN.
- The function type and name are declared on a separate line. This causes Better BASIC to respond interactively by creating a new workspace for the function and display and by displaying the memory available for it. Consequently, functions can use any range of valid line numbers without conflicting with other functions, procedures, and the main program.
- Parameters are declared by first listing their type, the keyword ARG:, and the parameter name. Like Better BASIC procedures, parameters can be assigned default values.
- The result of the function is returned using the standard identifier

RESULT.

 Recursive functions need to declare the function and any local variables as external. This enables Better BASIC to allocate new addresses instead of using the same ones as in the calling function.

Examples 1-4, page 128, show versions of the factorial function written in each of the BASIC implementations discussed here. Recursive versions are used with all implementations

except QuickBASIC. Example 5, page 128, shows a simple Better BASIC function that returns the logarithm to any base. The default is base 10. That is, using the function *LOGN(X)* returns the base-10 logarithm, and *LOGN(X,BAISE)* returns the logarithm of *X* to base *BAISE*.

DD.

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Lattice Works

LATTICE ANNOUNCES MICROSOFT WINDOWS SUPPORT IN VERSION 3.2

Version 3.2 of the Lattice MS-DOS C Compiler features full support for Microsoft Windows—including the "far," "near," and "pascal" keywords.

In addition, version 3.2 includes the ability to generate more than 64K bytes of static data and to declare objects larger than 64K bytes. It also includes improved support for ROM-based applications via the "const" data type. Version 3.2 is a significant release because it eliminates Microsoft's claimed monopoly on future MS-DOS C development tools. Now that the Lattice MS-DOS C Compiler supports a window interface, programmers using Lattice C can avoid the problems caused by switching to a different compiler. \$500.00

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Amiga's FFP format floating point library, and multi-tasking support.

With Version 3.1, Lattice has broken free of the reliance on the Amiga standard linker and object file format. This new release includes completely new expanded documentation, and a Lattice assembler and linker which remain compatible with previous software but allows professional programmers to take advantage of both the Amiga's speed and the industry's standardization.

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LATTICE RELEASES NEW VERSIONS OF C CROSS COMPILER AND LINKER

Version 3.1 of the Lattice C Cross Compiler to MS-DOS and version 2.12 of the Plink86Plus Overlay Linker are now available for Sun and Apollo workstations as well as the DEC VAX Family of processors running VMS, UNIX or Berkeley UNIX.

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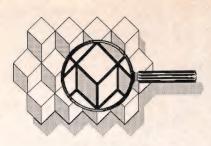
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OF INTEREST



Languages

Computer Crossware Labs has introduced Real BASIC, a BASIC interpreter for the Atari 520 and 1040 ST that executes BASIC code 20 to 100 times faster than ST BASIC while maintaining full compatibility. It contains an in-line Motorola-compatible assembler that allows you to switch to assembly language without leaving the interpreted BASIC environment and to have full access to BASIC variables while in assembly mode. Real BASIC's features include an integrated, full-screen editor that supports ST BASIC commands and some Micro Emacs commands; extended graphics instructions; and function calls that deal with the mouse and joystick ports. Real BASIC sells for \$69.95. Reader Service No. 16. Computer Crossware Labs Inc.

516 Fifth Ave., Ste. 507 New York, NY 10036 (212) 677-3686

Digitalk has released two optional extension kits to accompany Release 1.2 of Smalltalk/V for PC-DOS machines. The first kit integrates full EGA color capabilities into the Smalltalk/V environment, greatly enhancing the system's bit-mapped graphics. The second kit, called Goodies, offers several new programming kernels and lets you extend Smalltalk to handle applications that require discrete event simulation, forwardchaining interface operations, and connection to external sensors and instrumentation. Release 1.2 of Smalltalk/V is priced at \$99, and updates to 1.1 can be downloaded from CompuServe and BIX or obtained on disk from the company for \$10. The optional extension kits cost \$49 each. Reader Service No. 17.

Digitalk Inc. 5200 W. Century Blvd. Los Angeles, CA 90045 (213) 645-1082

Rational Visions is shipping a fullfeatured PROLOG programming system for the Atari ST. The system uses the Edinburgh standard syntax, making it compatible with most popular tutorials on PROLOG. Features include an Emacs-style text editor within the system's interpreter that allows you to interact with external disk files, the internal knowledge base, or invoke PROLOG goals from within the editor; a built-in grammar-rule translator that allows the development of natural-language interfaces; support for floating-point values; and extended math functions. Also, all PROLOG's metaprogramming primitives have been implemented, as have Atarispecific features such as GEMDOS primitives and interfaces to VDI and AES. The system is not copy-protected and supports user-written applications free of licensing restrictions. The package sells for \$39.95. Reader Service No. 18.

Rational Visions 7111 W. Indian School Rd., Ste. 131 Phoenix, AZ 85031 (602) 846-0371

Datalight has introduced Optimum-C, a full-featured global optimizing C compiler for IBM PCs and compatibles that supports the Unix System V C language along with several proposed ANSI extensions. Features include 8087 and software floatingpoint support, strong type checking, ROMable code generation, a MAKE program, MS-DOS object files format, Lattice C compatibility, and full library/start-up source code. Support is provided for compact-, small-, medium-, and large-memory models. Source licenses and support contracts are available for the compiler. Optimum-C sells for \$139. Reader Service No. 19.

Datalight P.O. Box 82441 Kenmore, WA 98028 (206) 367-1803

Tools

Devpac Amiga, from HiSoft, is a

68000 program development system for the Amiga computer that features a combined editor and assembler and a symbolic disassembler/debugger. The full-screen editor can be used from either the mouse or from the keyboard. The assembler is a complete, fast, macro assembler that supports include files read from disk, conditional assembly, and Motorolastandard macro handling and that can produce executable or linkable code. The debugger includes all the expected commands, such as breakpoints and single-stepping, and also allows you to use your original symbols when debugging programs. It uses its own screen for its display so as not to disturb that of other programs. Devpac Amiga is available in the United States from Apex Resources. It runs on any Amiga with at least 512K RAM and costs \$99.95. Reader Service No. 20.

Apex Resources 129 Sherman St. Cambridge, MA 02140 (800) 343-7535 In MA (617) 876-2505

Numerical Algorithms Group (NAG) now offers 172 selected routines from its extensive NAG Fortran Library for use on workstations and personal computers. The NAG Fortran Workstation Library provides the routines most frequently used by engineers and PC programmers. Each routine includes an example program, the source, and data results to illustrate the routine's usage. The NAG Fortran Workstation Library is available for the DEC MicroVAX, IBM PC line, and Sun workstations. License fees range from \$1,296 for a single workstation to \$384 each for 11 or more. Reader Service No. 21. Numerical Algorithms Group Inc.

Numerical Algorithms Group Inc. 1101 31st St., Ste. 100 Downers Grove, IL 60515 (312) 971-2337

P-tral, from **Woodchuck Industries**, is a native code BASIC-to-Pascal translation package that translates IBM BASICA/MS BASIC to Turbo Pascal. The program is interactive and lets you pick out or name subroutines as well as rename variables not fitting Pascal criteria. The program works

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(continued from page 132)

best with a hard disk and requires MS-DOS or PC-DOS 2.1 or later with AN-SI.SYS. P-tral sells for \$179. Reader Service No. 22.

Woodchuck Industries Inc. 340 W. 17th St., #2B New York, NY 10011 (212) 924-0576

RTC Plus, a FORTRAN- and RATFOR-to-C translator package from **Cobalt Blue**, allows you to tap old FORTRAN code for new C development. The translator is designed for translating non-I/O FORTRAN libraries and code in which I/O is concentrated in a few routines. Also, more than 95 percent of STUG'S RATFOR statements are supported by RTC Plus for complete translations of clean RATFOR code. RTC Plus runs under MS-DOS, Version 2.2 or later and costs \$325. Reader Service No. 23.

Colbalt Blue 1683 Milroy, Ste. 101 San Jose, CA 95124 (408) 723-0474

TurboPower Software has released an upgrade to T-DebugPLUS, a symbolic run-time debugger for Turbo Pascal. The new version (1.04) allows Turbo Pascal programmers to debug code in overlays and access CPU registers and memory. Other improvements include easier customization of the debugger and a new MAKELST utility program that creates a commented disassembly listing of Turbo Pascal programs. T-DebugPLUS requires Turbo Pascal 3.0 and a PC-DOS machine with 256K RAM. It sells for \$60 (\$10 for upgrade). Reader Service No. 24.

TurboPower Software 3109 Scotts Valley Dr., Ste. 122 Scotts Valley, CA 95066 (408) 438-8608

Laney Systems has released StruBAS 2.0, a structured BASIC toolkit for application development with Quick-BASIC and IBM BASIC 2.0 compilers in network and single-user environments. Integrated tools include enhanced structured programming facilities, screen handling, a Btrieve interface, StruBAS/ISAM for those not wishing to buy Btrieve, an object library of utility subroutines, and utility programs. A preprocessor ex-

pands BASIC with new commands to produce a more readable and structured language while making all BASIC's features and constructs available as well. StruBAS programs are translated by the preprocessor to compilable BASIC, which is then compiled and linked to produce executable programs. StruBAS is not copyprotected, and no royalties are charged for programs written using the package. It requires a PC-DOS machine with 320K RAM and a hard disk and sells for \$495. Reader Service No. 25.

Laney Systems Inc. 3 Office Park Dr., Ste. 105 Little Rock, AR 72211 (501) 225-7755

BASIC Program Analyzer, from the Expert Systems Division of Expert Information Systems, is a rulebased system that reads Microsoft's GW-BASIC program files in ASCII format and performs analysis, listing, and cross-reference functions upon them. The analyzer prints individually tabulated cross-references of line numbers, variables, keywords, quoted strings, and numerical constants found in theograms. Error, warning, and comment messages are inserted where applicable into tabulated cross-references, and an ASCII file is produced of all portions of the printout, including the cross-references. Optional parameters can be used to modify some of the analyzer's operations, such as testing for compiler syntax compatibility and for redirecting output. BASIC Program Analyzer is available for most MS-DOS computers and costs \$99.95. Reader Service No. 26.

Expert Information Systems Inc. Expert Systems Division P.O. Box 1310 El Campo, TX 77437-1310 (409) 543-9222

Libraries

United States Software Corp. has released a 68HC11 floating-point library. This complete math package operates on the 68HC11, is designed in a modular format, and is delivered in source assembly language for maximum flexibility. The library conforms to the IEEE 754 Floating Point Standard. The 68HC11 FPAC/

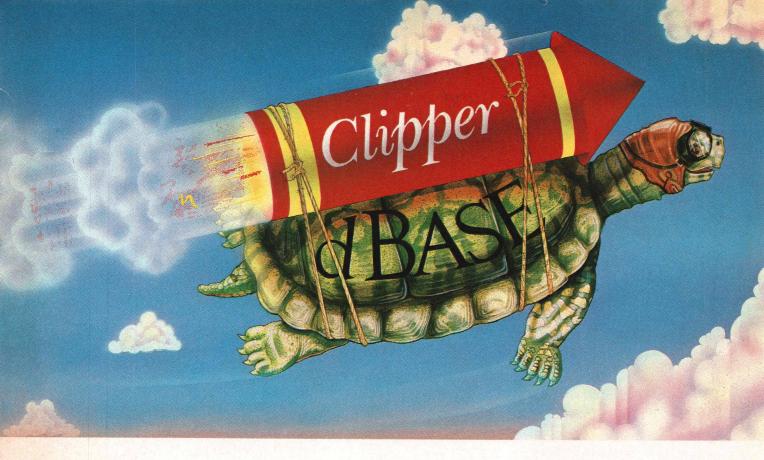
DPAC is available in either single-precision or double-precision formats. The library includes trigonometric, logarithmic, and exponentiation functions, and data-conversion and floating-point utility procedures. It is available on IBM PC-compatible media for a one-time fee of \$950. Reader Service No. 27.

United States Software Corp. 14215 N.W. Science Park Dr. Portland, OR 97229 (503) 641-8446

LINLIB, from Information and Graphic Systems, is a library of C routines for developers of scientific and statistical software. The package performs all basic operations on vectors and matrices, with no limitation on size of vector or matrix. Of major importance are the routines that factor matrices (LU, QR, and Cholseky factors). These factorizations enable you to solve any kind of system equations-underdetermined, determined, and overdetermined. Also included in LINLIB are routines based on the B-spline for manipulating splines. LINLIB sells for \$150. Reader Service No. 28.

Information and Graphic Systems 15 Normandy Ct. Atlanta, GA 30324 (404) 231-9582

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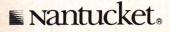
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SWAINE'S FLAMES

think I may be the first person who has ever actually seen a computer bug.

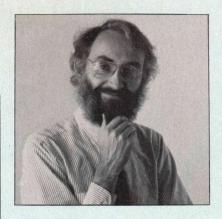
It came to me at the end of a long and disillusioning night when my body, no doubt in sympathy with the new punitive tariffs on Japanese goods, began rejecting a dinner of overripe sushi. En route from bathroom to loft. I heard an odd tapping sound from the office and stopped to investigate.

In my collywobbled state I had neglected to turn off the Macintosh, and what I saw in the glow from the screen was a bug-it looked a lot like a cockroach—repeatedly climbing to the top of the Mac and leaping to the keyboard to enter a character. The exertion must have been tremendous, and he looked pretty beat when he finally finished and crawled into the disk drive. I'm not surprised he didn't bother with capitalization or punctuation (the bottom row of keys was almost too long a jump for him).

He left this behind.

boss i m afraid that the computer curmudgeon john dvorak has been talking about our beloved dr dobb s journal of whatever again if dr. dobb s became outspoken it would be dangerous dvorak said in one of those tasty magazines you leave lying around this office i wish you would leave a crust of bread i m getting tired of paper and glue

i wonder does he mean like that 1940s science fiction writer was dangerous when he explained how an a bomb would work or like the white house press corps would be dangerous if it decided it wasn t the white house s press corps



or like sesame street would be dangerous if it taught children to question their teachers and their parents or maybe like infoworld was dangerous when it was publishing dvorak

sometimes i think that john dvorak is a modeless dialog box the max headroom of the personal computer industry

how could dr dobb s become dangerous is what i d like to know maybe by examining legislator s voting records on software issues or by doing a study of compiler pricing vs value or by publishing detailed specs on data formats of commercial products so that data held hostage by obsolete applications could be liberated or by showing how to develop software for the 386 without waiting for you know who or by questioning the need for an operating system on a personal computer at all

i used your modem to call dvorak s number and talked with his dog sparky maybe dr dobb s should list the companies that preannounce products to outflank their competitors sparky suggested or do their beta testing in

release version one point oh or cripple their products with copy protection tricks or refuse to eliminate known bugs what s this about bugs i asked him but he just barked

one thing i know is that dr dobb s could stand to be a little more rough and tumble a bit more capricious and corybantic there s a dance in the old dame vet and if that s dangerous it s fine by me please save this file as i can t work the mouse nor would i want to i haven t always had the best of relations with mice yours archy

Scarab journalism was invented by Don Marquis, a New York Sun columnist, when he found a cockroach composing in his office late one night some seventy years ago, learned that the bug was a reincarnated vers libre poet, and cannily conned the poet roach into meeting copy deadlines for him for the next ten years.

The saga of archie the cockroach itself allegedly owes its existence to the newspaper account of a typewriting rat who plied his or her art by night circa 1916 in a Dobb's Ferry, New York, garage. By some odd symmetry in the transmigration of souls, a Dobb has once again played host to verse vermin.

Fine by me.

Michael Swades Michael Swaine editor-in-chief



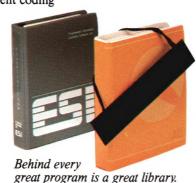
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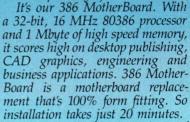
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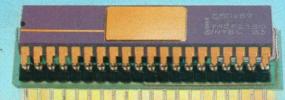
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